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```

```
1 0001 0 MODULE RM3SPLUDR (LANGUAGE (BLISS32) ,
2 0002 0 IDENT = 'V04-000'
3 0003 0 ) =
4 0004 1 BEGIN
5 0005 1
6 0006 1 *****
7 0007 1 *
8 0008 1 * COPYRIGHT (c) 1978, 1980, 1982, 1984 BY
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25 0025 1 *
26 0026 1 *
27 0027 1 *****
28 0028 1
29 0029 1 ++
30 0030 1
31 0031 1 FACILITY: RMS32 INDEX SEQUENTIAL FILE ORGANIZATION
32 0032 1
33 0033 1 ABSTRACT:
34 0034 1 split user data record buckets
35 0035 1
36 0036 1
37 0037 1 ENVIRONMENT:
38 0038 1
39 0039 1 VAX/VMS OPERATING SYSTEM
40 0040 1
41 0041 1 --
42 0042 1
43 0043 1
44 0044 1 AUTHOR: Wendy Koenig CREATION DATE: 5-JUL-78 14:46
45 0045 1
46 0046 1
47 0047 1 MODIFIED BY:
48 0048 1
49 0049 1 V03-013 JWT0157 Jim Teague 23-Feb-1984
50 0050 1 When RMS attempted to calculate whether a series of
51 0051 1 duplicate records (including the new record) would
52 0052 1 fit within a single bucket, it neglected to account
53 0053 1 for the fact that the first record in the chain will
54 0054 1 undergo full expansion when it ends up as the first
55 0055 1 record in the new bucket. If it is currently partially
56 0056 1 compressed based on the previous key, then that could
57 0057 1 (and sometimes DID) cause bucket overflow when the
```

58 0058 1
59 0059 1
60 0060 1
61 0061 1
62 0062 1
63 0063 1
64 0064 1
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66 0066 1
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69 0069 1
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74 0074 1
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76 0076 1
77 0077 1
78 0078 1
79 0079 1
80 0080 1
81 0081 1
82 0082 1
83 0083 1
84 0084 1
85 0085 1
86 0086 1
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99 0099 1
100 0100 1
101 0101 1
102 0102 1
103 0103 1
104 0104 1
105 0105 1
106 0106 1
107 0107 1
108 0108 1
109 0109 1
110 0110 1
111 0111 1
112 0112 1
113 0113 1
114 0114 1

duplicate chain is moved merrily into the new bucket.
Keep track of the compression count for the first
record in the dup chain, and add it to the total
size of the chain before comparing to bucket size.

V03-012 JWT0142 Jim Teague 16-Dec-1983
Correct incorrect bucket VBN comparison.

V03-011 MCN0008 Maria del C. Nasr 22-Mar-1983
More changes in the linkages

V03-010 MCN0007 Maria del C. Nasr 28-Feb-1983
Reorganize linkages

V03-009 TMK0004 Todd M. Katz 10-Nov-1982
At the present time, under certain circumstances, the number
of RRVs which will be required to be created when a simple
two-bucket split is done is being incorrectly calculated. This
will happen only during SUPDATES when the record being updated
is to go into the old (left) bucket and prior to the split is
in its original bucket. Even then it does not happen under all
possible circumstances, but only when duplicate records are
involved. It is possible that the number of RRVs calculated to
be required will be several less than the actual number which
will be needed. Under certain circumstances, the number of RRVs
needed may actually be calculated as a negative number - an
impossibility. Much depends upon the bucket composition. While
this does not influence the actual creation of RRVs, what it
does affect is where the bucket split point is calculated to be
since RRVs to be created do take up space in the old (left)
bucket. In fact, this problem came to my attention because of
the occurrence of a bucket split which resulted in the right
bucket, the new bucket, being empty, and the old (left) bucket
containing all the records even though there was no room for
them (or the bucket split would not have been required in the
first place). This split was caused by the number of RRVs
required being calculated as -1 instead of 0 such that, instead
of having the RRV spacial requirements added to the left bucket
size requirements, they were subtracted.

To fix this problem I have adjusted how the number of needed
RRVs are to be calculated. To start, the number of needed RRVs
is calculated to be the number of records (including the record
being updated which is not currently in the bucket) whose
original bucket is the bucket splitting. Then, as the split
point of the bucket is adjusted from left to right, this number
is decremented as records (which are in their original bucket)
are designated to stay in the left or old bucket. This is where
my change comes in. Previously, that the updated record was to
stay in the old bucket was determined at several different
points, and each time the count of the number of needed RRVs
was decremented, as long as the other conditions were met.
Unfortunately, this allowed for this determination to take place
more than once, and for the RRV count to be decremented multiple
times for the same record. My fix prevents this from occurring.
While it is still determined in several places that the updated
record is to go in the old bucket, I have made sure that those

115 0115 1
116 0116 1
117 0117 1
118 0118 1
119 0119 1
120 0120 1
121 0121 1
122 0122 1
123 0123 1
124 0124 1
125 0125 1
126 0126 1
127 0127 1
128 0128 1
129 0129 1
130 0130 1
131 0131 1
132 0132 1
133 0133 1
134 0134 1
135 0135 1
136 0136 1
137 0137 1
138 0138 1
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158 0158 1
159 0159 1
160 0160 1
161 0161 1
162 0162 1
163 0163 1
164 0164 1
165 0165 1
166 0166 1
167 0167 1
168 0168 1
169 0169 1
170 0170 1
171 0171 1

places are orthogonal to one another, so that the RRV count is not decremented more than once for the same record, the record whose update is causing the split.

V03-008 KBT0234 Keith B. Thompson 23-Aug-1982
Reorganize psects

V03-007 TMK0003 Todd M. Katz 02-Jul-1982
Implement RMS cluster solution for next record positioning. The next record positioning context is now kept in the IRAB, where it may be retrieved from, instead of in the NRP list which has been eliminated. When referring to the RFA address of the new/changed primary data record use the subfields IRBSL_PUTUP_VBN and IRBSW_PUTUPD_ID.

V03-006 MCN0006 Maria del C. Nasr 29-Jun-1982
Allow keys of different data types other than string in prologue 3 files.
Change all CH\$COMPARE calls to RMS\$COMPARE_KEY to compare keys taking into consideration the different data types.

V03-005 MCN0005 Maria del C. Nasr 11-Jun-1982
Eliminate overhead at end of data bucket that was to be used for duplicate continuation bucket processing.

V03-004 TMK0002 Todd M. Katz 31-May-1982
Performance enhancements. I have made four changes to the routine RM\$SPLIT_UDR_3 which should cut down on the length of the bucket scans required at various times to re-expand keys. The enhancements involve setting the IRAB field IRBSL_LST_NCMP to the current record if key compression is enabled and the key of the current record is zero front compressed during various bucket scans required in the determination of the split point(s). The first of these scans is made to position to and extract the key of the last record in the bucket. The second and third scans are when the split code has decided that the best split point is one record previous to the current position, and must scan the bucket to obtain the record previous to that position so its key can be extracted and used during the index update. The fourth scan is the right-to-left record-by-record scan made to decide whether a two-bucket split is possible, and if so, where is the best place to split. In all four cases, I have added code to set the last noncompressed record pointer before continuing the scan with the next record, if the current record was zero front compressed.

V03-003 TMK0001 Todd M. Katz 10-May-1982
The algorithm for determining the split point of a prologue three data bucket with compressed keys first determines whether a two-bucket split can be done by scanning the old bucket from left-to-right record-by-record determining whether the lefthand sides and righthand sides of each possible split point will fit into a bucket. This size determination must take into account the position of insertion of the new (or updated) record, and the size determination of the righthand side must take into account the number of characters currently front

compressed of what will become its low-order (and thus non-compressed key) record. What was missing, and what this change rectifies, is that what may become the low-order record of the righthand bucket is in fact the new (updated) record whose insertion is forcing this split to take place. In this case, the number of front compressed characters to be added to the righthand side total must come from the compressed key in keybuffer 5, if this is an \$UPDATE, or from the compressed key in the record buffer whose address is stored in IRB\$L_RECBUF, if this is a \$PUT. This change will be included as a patch on the V3.1 update floppy.

V03-002 MCN0004 Maria del C. Nasr 31-Mar-1982
Do not count records that will not need rrv's when moved out of the bucket. Their id's cannot be recycled in plg 3 files.

V03-001 MCN0003 Maria del C. Nasr 25-Mar-1982
Use macro to calculate keybuffer address.

V02-016 DJD0001 Darrell Duffy 1-March-1982
Fix references to RBF for better probing

V02-015 MCN0002 Maria del C. Nasr 09-Jul-1981
Fix a problem with update of the first record in a duplicate chain, in both old code, and new code. Also fix problem in new code with non-compressed keys.

V02-014 MCN0001 Maria del C. Nasr 02-Jun-1981
Add the routine to split prologue 3 data buckets.

V02-013 REFORMAT Ron Schaefer 23-Jul-1980 14:10
Reformat the source

V02-012 CDS0000 Christian Saether, 01-Jan-1980 15:00
FIX PROBLEM WHEN SPLITTING BECAUSE OUT OF ID'S.

REVISION HISTORY:

Wendy Koenig, 18-SEP-78 16:53
X0002 - FIX BUG IN BACKING UP PAST NEW RECORD

Wendy Koenig, 19-SEP-78 10:52
X0003 - DO SPLIT AT POINT OF INSERT IF ASCENDING ORDER DETECTED

Wendy Koenig, 12-OCT-78 13:21
X0004 - CHANGES FOR UPDATE

Wendy Koenig, 18-OCT-78 14:03
X0005 - IF WE PASS BY POS_INSERT WHILE SKIPPING OVER DUPS, NOTE IT

Wendy Koenig, 18-OCT-78 14:37
X0006 - FIX SOME PROBLEMS W/ 4-BKT SPLIT (\$UPDATE ONLY)

Wendy Koenig, 24-OCT-78 14:03
X0007 - MAKE CHANGES CAUSED BY SHARING CONVENTIONS

Wendy Koenig, 7-NOV-78 8:58

```

229 0229 1  X0008 - FIX EMPTY_BKT BUG, NOT BEING SET WHEN SHOULD BE
230 0230 1
231 0231 1  Wendy Koenig, 22-JAN-79 17:03
232 0232 1  X0009 - IF LOA TRIES TO FORCE US TO SPLIT ALL DUFS, SPLIT AT POS_INS
233 0233 1
234 0234 1  Wendy Koenig, 24-JAN-79 9:51
235 0235 1  X0010 - CONDITION HOLDS EVEN IF LOA NOT SET
236 0236 1
237 0237 1  Wendy Koenig, 29-JAN-79 15:58
238 0238 1  X0011 - FIX PROBLEM W/ DUPLICATE ENTRIES IN INDEX
239 0239 1
240 0240 1  *****
241 0241 1
242 0242 1  LIBRARY 'RMSLIB:RMS';
243 0243 1
244 0244 1  REQUIRE 'RMSSRC:RMSIDXDEF';
245 0309 1
246 0310 1  ! define default psects for code
247 0311 1
248 0312 1  PSECT
249 0313 1  CODE = RMSRMS3(PSECT_ATTR),
250 0314 1  PLIT = RMSRMS3(PSECT_ATTR);
251 0315 1
252 0316 1  ! Linkages
253 0317 1
254 0318 1  LINKAGE
255 0319 1  L_COMPARE_KEY,
256 0320 1  L_PRESERVE1,
257 0321 1  L_RABREG_4567,
258 0322 1  L_RABREG_67,
259 0323 1  L_REC_OVRD,
260 0324 1
261 0325 1  ! Local linkages
262 0326 1
263 0327 1  RL$BUILD_KEY = JSB ( ) :
264 0328 1  GLOBAL (R_IDX_DFN) PRESERVE(1,2,3,4,5),
265 0329 1  RL$MOVE_KEY = JSB (REGISTER = 0, REGISTER = 6) :
266 0330 1  GLOBAL (R_RAB, R_IRAB, R_IFAB, R_IDX_DFN, R_BKT_ADDR);
267 0331 1
268 0332 1  ! Forward Routine
269 0333 1
270 0334 1
271 0335 1  FORWARD ROUTINE
272 0336 1  RMSBUILD_KEY : RL$BUILD_KEY NOVALUE,
273 0337 1  RMSMOVE_KEY : RL$MOVE_KEY NOVALUE;
274 0338 1
275 0339 1  ! External Routines
276 0340 1
277 0341 1
278 0342 1  EXTERNAL ROUTINE
279 0343 1  RMSMOVE : RL$PRESERVE1,
280 0344 1  RMSRECORD_VBN : RL$PRESERVE1,
281 0345 1  RMSRECORD_KEY : RL$PRESERVE1,
282 0346 1  RMSREC_OVRD : RL$REC_OVRD,
283 0347 1  RMSVBN_SIZE : RL$PRESERVE1,
284 0348 1  RMSCOMPARE_KEY : RL$COMPARE_KEY,
285 0349 1  RMSCOMPARE_REC : RL$RABREG_67,

```

RM3SPLUDR
V04-000

M 15
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

VAX-11 Bliss-32 V4.0-742
[RMS.SRC]RM3SPLUDR.B32;1

Page 6
(1)

: 286

0350 1

RM\$GETNEXT_REC

: RL\$RABREG_67;


```
288 0351 1 ++
289 0352 1
290 0353 1
291 0354 1
292 0355 1
293 0356 1
294 0357 1
295 0358 1
296 0359 1
297 0360 1
298 0361 1
299 0362 1
300 0363 1
301 0364 1
302 0365 1
303 0366 1
304 0367 1
305 0368 1
306 0369 1
307 0370 1
308 0371 1
309 0372 1
310 0373 1
311 0374 1
312 0375 1
313 0376 1
314 0377 1
315 0378 1
316 0379 1
317 0380 1
318 0381 1
319 0382 1
320 0383 1
321 0384 1
322 0385 1
323 0386 1
324 0387 1
325 0388 1
326 0389 1
327 0390 1
328 0391 1
329 0392 1
330 0393 1
331 0394 1
332 0395 1
333 0396 1
334 0397 1
335 0398 1
336 0399 1
337 0400 1
338 0401 1
339 0402 1
340 0403 1
341 0404 1
342 0405 1
343 0406 1
344 0407 1
```

ALGORITHM FOR A TWO-BUCKET 50/50 SPLIT

GIVEN: that the record will not fit in the bucket.
i.e., we must split the bucket in some form.

INPUTS: the bucket, the record size and the position
to insert the record in the bucket

GOALS: to make the split as efficient as possible:

- 1) to create the fewest number of new buckets possible
- 2) to use the space in the available buckets efficiently --
i.e., the bucket with the most available space should contain
the most data after the split.

ALGORITHM IN A NUTSHELL:

- 1) A two-bucket split will occur IF AND ONLY IF there is a point in
the bucket at which all records to the left of the point and
necessary rrv's fit in a single bucket and all records to the right
of the point fit in a single bucket. This point must be on a
record boundary and must not be in the middle of a chain of
duplicates.
- 2) Given that such a point exists, the most optimal point for a
2-bucket split is the point at which the actual data records
are divided evenly between the available space in the original
bucket and the available space in the new (previously empty) bucket.

In theory, therefore, the idea is to find a point in the bucket such that
the point is on a boundary between duplicate records and that

- 1) records in the left hand side / space in the left hand bucket
=
- 2) records in the right hand side / space in the right hand bucket.

In practice, the idea is to minimize the absolute difference between
ratio 1) and ratio 2). Just to make it clearer, "records in the left
hand side" means the total size of the data records left of this point
(not including rrv's of any kind) and "space in the left hand bucket"
means the bucket size of the data bucket minus the total size of existing
rrv's and the total size of rrv's which would have to be generated.

IMPLEMENTATION:

This algorithm needs two scans of the bucket. The first scan is very
quick and determines the total size of the existing rrv's. It also
counts the number of rrv's that would have to be generated in a worst
case situation (i.e., all records would be moved out). Thus, as the
second scan proceeds, all information needed to calculate the above
ratios EXACTLY is available.

In order for there to be a 2-bucket split, there must be a point
in the bucket such that the right hand side fits in a single bucket.
Scanning from the left (beginning) of the bucket, we can find the
first point at which the right hand side will fit. Since as we
continue scanning to the right we are decreasing the right hand side,
the righthand side will continue to fit as we scan rightward.

If at this point, the left hand side will not fit, we can not possibly

345 0408 1
346 0409 1
347 0410 1
348 0411 1
349 0412 1
350 0413 1
351 0414 1
352 0415 1
353 0416 1
354 0417 1
355 0418 1
356 0419 1
357 0420 1
358 0421 1
359 0422 1
360 0423 1
361 0424 1
362 0425 1
363 0426 1
364 0427 1
365 0428 1
366 0429 1
367 0430 1
368 0431 1
369 0432 1
370 0433 1
371 0434 1
372 0435 1
373 0436 1
374 0437 1
375 0438 1
376 0439 1
377 0440 1
378 0441 1
379 0442 1
380 0443 1
381 0444 1
382 0445 1
383 0446 1
384 0447 1
385 0448 1
386 0449 1
387 0450 1
388 0451 1
389 0452 1
390 0453 1
391 0454 1
392 0455 1
393 0456 1
394 0457 1
395 0458 1
396 0459 1
397 0460 1
398 0461 1
399 0462 1
400 0463 1
401 0464 1

have a 2-bucket split, since continuing our scan would only make the left hand side larger (or it may stay the same size). Once we have found a point at which we can do a 2-bucket split we can always return to it, if in our search for a more optimal split point we leave the range in which the left hand side will fit. This can occur if the records in the bucket are of minimal size, that is to say that the records are the same size as rrv's and therefore no additional space for data is gained by scanning to the right.

At this point (the first point at which the right hand side will fit), ratio 1 is less than ratio 2. As we proceed to the right, ratio 1 will increase and ratio 2 will decrease. This is due to the fact that the size of the right hand side (the numerator of ratio 2) decreases as we move rightward and the available space in the right bucket is a constant (the denominator of ratio 2). In ratio 1, both the numerator and denominator are increasing, but the numerator is increasing at a faster rate. As soon as we reach a point where ratio 1 is greater than or equal to ratio 2, we can stop the scan. Now we have a choice of split points available. We can use this point or the one immediately before it (if such a point exists). The decision is made by minimizing the absolute difference between the ratios and we have an optimal split point.

Things become complicated by the presence of duplicate records. When duplicate records occupy more than one bucket, the subsequent buckets are termed continuation buckets. In prologue version 1 and 2 files, there is a pointer from the index to the first bucket only, and the continuation buckets are found only from the horizontal links in the buckets. At one point, it was thought that disaster would ensue if the continuation buckets ever had a record with a key value other than that of the duplicates. Normally, this will not happen because the key value of the index pointer to the first bucket will be the same as that of the duplicate records in the chain and a record with a higher key value will follow the next index pointer down when positioning for insert. This will place it in the next bucket beyond the chain of continuation buckets. However, a bucket in which the record with the highest value has been deleted that subsequently receives a series of duplicates creating a continuation chain will generate a situation where a record with a key value between that of the duplicate chain and the original high key value of the bucket will be inserted at the end of the duplicate chain. A far more common situation is created by RMS-11 (at least thru v1.5) when loading a file in ascending primary key sequence will pack the buckets 100% (or the load factor) full, including records of non-dupe key values at the end of continuation buckets.

At any rate, the fact that the situation exists notwithstanding, much of the code that follows is there to keep duplicates together when splitting, and to put only records with duplicate key values in continuation buckets. It appears to be a good thing to do from an overall space efficiency standpoint over a period of time, but the code could probably be considerably simplified if it wasn't necessary. With all that in mind, the split situation with all possible record 'partitions' within the bucket prior to splitting is as follows:

402 0465 1
403 0466 1
404 0467 1
405 0468 1
406 0469 1
407 0470 1
408 0471 1
409 0472 1
410 0473 1
411 0474 1
412 0475 1
413 0476 1
414 0477 1
415 0478 1
416 0479 1
417 0480 1
418 0481 1
419 0482 1
420 0483 1
421 0484 1
422 0485 1
423 0486 1
424 0487 1
425 0488 1
426 0489 1
427 0490 1
428 0491 1
429 0492 1
430 0493 1
431 0494 1
432 0495 1
433 0496 1
434 0497 1
435 0498 1
436 0499 1
437 0500 1
438 0501 1
439 0502 1
440 0503 1
441 0504 1
442 0505 1
443 0506 1
444 0507 1
445 0508 1
446 0509 1
447 0510 1
448 0511 1
449 0512 1
450 0513 1
451 0514 1
452 0515 1
453 0516 1
454 0517 1
455 0518 1
456 0519 1
457 0520 1
458 0521 1

! low set ! low dupes !! high dupes ! high set !

^
point of insert (new record)

From the point of view of the split code, an update operation in which the record is growing and causes a split is identical (almost) to a new record being inserted. The original record is removed from the bucket after determining that the updated record will cause a split and the updated record is more or less treated as a new record. One of the most important differences is that in an update situation, the 'new' record gets the id of the old record, rather than a new id. Another is that because duplicate records are always inserted at the end of a chain of duplicates, some split cases can only occur on an update operation.

In fact, the situation postulated above can happen only in an update situation, and may cause 3 new buckets to be generated on the split operation. This will occur when the updated record is in the middle of a group of duplicate records and grows to the extent that no other records will fit in the bucket with it anymore. Using 1 byte key values to make this easier to visualize, the bucket above prior to the update may look like this (the artificial partitioning of the bucket corresponds to the breakdown above):

! A B C ! D D D ! D ! D D ! E F G !

^\
this record gets updated

The record being updated changes size and grows such that it needs an entire bucket for itself. To keep all the duplicates together, the situation after the split looks like this:

----- ----- ----- -----
! A B C D D D ! -> ! D ! -> ! D D ! -> ! E F G !

this is the these two are
original bucket continuation buckets

The original bucket probably had an index pointer with the value 'G' pointing to it (or some previous bucket if there was a previous index update failure). After the split, the key value for that pointer will be updated to have the key value 'D', and the key value that used to point to it (probably 'G'), will now point to the right hand bucket (with 'E', 'F', and 'G' in it). The continuation buckets never have an index pointer to them.

All other split situations are a variation of this one, with one or more of the 'partitions' not present, dependent on the key value and position of insert within the bucket of the record being inserted or updated. For example, if there are no duplicates, there are no 'low dupes' or 'high dupes'. Or if the position of insert is at the end of the bucket, there is no 'high set'.

Now that I've started on it, may as well try to document some other

interesting split situations. Note that a '2 bucket split' means that there are 2 buckets after the split, i.e., 1 new bucket is added. The situation described above is a 4 bucket split.

The most interesting split from an index updating point of view is the 3 bucket split where a record is being inserted in the middle of the bucket and doesn't fit in a bucket with either the low set or the high set. Again with 1 byte key values to illustrate:

G (this is supposed to represent an index
pointer to this bucket with key value 'g')

↓
v

! A B C ! ! E F G !

/\

new record with key value 'D' inserted, but is so large
that it has to have bucket of its own.

After split (with new index pointers):

C D G
↓ ↓ ↓
v v v

! A B C ! -> ! D ! -> ! E F G !

The new pointer 'C' is the bucket pointer from the original index record 'G' with the new key value 'C'. The 'D' pointer is an entirely new record (i.e., key value 'D' and bucket pointer). The pointer 'G' is the key value from the original record 'G' with a new bucket pointer. The bucket pointer for the 'D' bucket comes from `irb$1_vbn_mid` and the bucket pointer for the 'G' bucket comes from `irb$1_vbn_right`. Remember that all of this stuff works correctly if the index update failed and we got to the bucket that's splitting by following the horizontal bucket links at the data level. For example, consider the following case where prior index corruption exists:

G (index update failed when right hand bucket split off
during a previous insert operation)

↓
v

! A B C ! -> ! D ! ! F G !

/\

new record 'E' will be inserted here and cause split

After split:

E G
↓ ↓
v v

-----V-----V
! A B C ! -> ! D E ! -> . F G !

The reason for the index updating behavior becomes more obvious. The key value of the original down pointer 'G' has been changed to the new value 'E', but retaining the original bucket pointer. Note that we split the bucket with 'D' in it, yet there is no bucket pointer to it now (there wasn't before). The key value 'G' of the original bucket pointer 'G' has been used with a new bucket pointer for the new bucket created by the split (this is `irb$l_vbn_right`). Sometimes there will be a bucket split and no records will be in the left hand bucket after the split. This may happen if the record being inserted belongs at the beginning of the bucket, but there are enough rrv's present so that it doesn't physically fit. In that case, all of the existing records will have to be moved out also. This may also occur if there are no id values left in the bucket (typically caused by deleted rrv's). In this case, we would like to swing the index pointer away from the 'empty' bucket to keep random access times from deteriorating. As of prologue versions 1 and 2, however, it will remain in the horizontal link of data buckets. However, we can only change the down pointer if it already points to that bucket or we can potentially create crossed down pointers. The situation is illustrated below:

-----C-----I (index update failed on a previous split)
-----V-----V
! A B C ! -> ! D E F ! -> ! H I !
-----/\-----
record with key value 'G' will be
inserted here

Also presume that the bucket 'G' is being inserted in has so many rrv's in it that it won't fit into the existing bucket, even though it will fit into a bucket without any rrv's in it.

After split:

-----C-----I
-----V-----V
! A B C ! -> ! D E F ! -> ! rrv's only ! -> ! G H I !

Note that the index pointer 'I' was not moved to point to the new bucket. If it had been, the bucket containing 'D E F' would have been 'lost' by random access from the index. This condition is detected by setting `irb$l_vbn_left` to the vbn of the rrv only bucket. During the index update procedure, the pointer will be moved to point to the new bucket only if the existing down pointer points to the bucket that was split, i.e., `irb$l_vbn_left` (this

```
573 0636 1 is normally the case as index corruption is not normal). Note that
574 0637 1 an empty left hand bucket may also be present in a 3 bucket split
575 0638 1 situation.
576 0639 1
577 0640 1 Following is a list of the specific split cases handled in the
578 0641 1 code. They are basically variations of the above cases.
579 0642 1
580 0643 1 these are all the cases of 3 and 4 bucket splits that i can think of
581 0644 1 any or all of these cases can have the empty left-hand bucket
582 0645 1 -- this would occur if the first split point is at the beginning
583 0646 1 -- of the bucket and all data records got moved out
584 0647 1
585 0648 1 low dups exist -- no high dups
586 0649 1 low dups fit w/ rec
587 0650 1 3 bkt split low, low dups w/ rec, hi set -- rec goes w/ lo
588 0651 1 ( SPLIT TYPE 1 )
589 0652 1
590 0653 1 low dups don't fit w/ rec
591 0654 1 3 bkt split w/ rec in its own continuation bucket
592 0655 1 ( SPLIT TYPE 2, W/ DUPS SEEN )
593 0656 1
594 0657 1 hi dups exist -- no low dups
595 0658 1 hi dups fit w/ rec
596 0659 1 3 bkt split low, hi dups w/ rec, hi set
597 0660 1 ( SPLIT TYPE 1 )
598 0661 1
599 0662 1 hi dups don't fit w/ rec
600 0663 1 if no more hi, 3 bkt split low, rec, hi = hi dups is a cont. bkt
601 0664 1 ( SPLIT TYPE 2 )
602 0665 1
603 0666 1 if there is more hi, 4 bkt split low, rec, hi dups, hi
604 0667 1 ( SPLIT TYPE 2B )
605 0668 1
606 0669 1 no dups at all
607 0670 1 record goes in its own bucket, 3 bkt split
608 0671 1 (SPLIT TYPE 3 )
609 0672 1
610 0673 1 low dups and hi dups
611 0674 1 all dups fit together
612 0675 1 3 bkt split w/ dups in middle bkt
613 0676 1 ( SPLIT TYPE 1 )
614 0677 1
615 0678 1 no dups fit w/ record
616 0679 1 if no more hi, 3 bkt split low, rec = cont. bkt, hi = hi dups = cont. bkt
617 0680 1
618 0681 1
619 0682 1 ( SPLIT TYPE 2B )
620 0683 1
621 0684 1 if there is more hi, 4 bkt split low, rec = cont. bkt, hi dups, hi
622 0685 1 ( SPLIT TYPE 2B )
623 0686 1
624 0687 1 hi dups fit w/ record
625 0688 1 if no more hi, 2 bkt split low, rec w/ hi dups = cont. bkt
626 0689 1 ( this is a 2 bkt split case that the previous alg. wouldn't handle)
627 0690 1 ( SPLIT TYPE 4B )
628 0691 1
629 0692 1 if there is more hi, 3 bkt split low, rec w/ hi dups = cont. bkt, hi
```

:	630	0693	1	:	(SPLIT TYPE 4)
:	631	0694	1	:	
:	632	0695	1	:	low dups fit w/ record
:	633	0696	1	:	if lo and hi, 4 bkt split low, low dups w/ rec, hi dups, hi
:	634	0697	1	:	(SPLIT TYPE 5)
:	635	0698	1	:	
:	636	0699	1	:	if no lo and no hi, 3 bkt split (rrv's), low dups w/ rec, hi dups = cont.
:	637	0700	1	:	
:	638	0701	1	:	bkt
:	639	0702	1	:	(SPLIT TYPE 5, w/ empty original bkt and no high)
:	640	0703	1	:	
:	641	0704	1	:	if lo but no hi, 3 bkt split lo, low dups w/ rec, hi dups = cont. bkt
:	642	0705	1	:	(SPLIT TYPE 5, w/ no high)
:	643	0706	1	:	
:	644	0707	1	:	if hi but no low, 4 bkt split (rrv's), low dups w/ rec, hi dups, hi
:	645	0708	1	:	(SPLIT TYPE 5, w/ empty bkt)
:	646	0709	1	:	--

RMSMOVE_KEY

```
648 0710 1 %SBTTL 'RMSMOVE_KEY'
649 0711 1 ROUTINE RMSMOVE_KEY (ADDRESS, CUR_REC_ADDR) : RL$MOVE_KEY NOVALUE =
650 0712 1
651 0713 1 ++
652 0714 1
653 0715 1 FUNCTIONAL DESCRIPTION:
654 0716 1
655 0717 1 Routine to move the key from wherever it is desired into
656 0718 1 key buffer 2.
657 0719 1
658 0720 1 CALLING SEQUENCE:
659 0721 1 bsbw rmsmove_key (address,cur_rec_addr)
660 0722 1
661 0723 1 INPUT PARAMETERS:
662 0724 1 address from which to get the key from
663 0725 1 the current value of rec_addr
664 0726 1
665 0727 1 IMPLICIT INPUTS:
666 0728 1 BKT_ADDR,
667 0729 1 RAB -- user's buffer address
668 0730 1 IRAB -- pos_ins, rec_w_lo, keybuf
669 0731 1 IFAB -- kbufsz, prologue version
670 0732 1 IDX_DFN -- for call to record_key, and compression flags
671 0733 1
672 0734 1 OUTPUT PARAMETERS:
673 0735 1 none
674 0736 1
675 0737 1 IMPLICIT OUTPUTS:
676 0738 1 key is moved into key buffer 2
677 0739 1
678 0740 1 ROUTINE VALUE:
679 0741 1 none
680 0742 1
681 0743 1 SIDE EFFECTS:
682 0744 1 key is move' into key buffer 2
683 0745 1 AP is clobbered
684 0746 1
685 0747 1 --
686 0748 1
687 0749 2 BEGIN
688 0750 2
689 0751 2 BUILTIN
690 0752 2 AP;
691 0753 2
692 0754 2 GLOBAL REGISTER
693 0755 2 R_BDB,
694 0756 2 R_IMPURE,
695 0757 2 R_REC_ADDR_STR;
696 0758 2
697 0759 2 EXTERNAL REGISTER
698 0760 2 R_IFAB_STR,
699 0761 2 R_RAB,
700 0762 2 R_IRAB_STR,
701 0763 2 R_IDX_DFN_STR,
702 0764 2 R_BKT_ADDR_STR;
703 0765 2
704 0766 2 IF .CUR_REC_ADDR - .BKT_ADDR EQ'U .IRAB[IRBSW_POS_INS]
```



```

: 705      0767 2      AND
: 706      0768 2      .IRAB[IRB$V_REC_W_LO]
: 707      0769 2      . THEN
: 708      0770 2      BEGIN
: 709      0771 3      AP = 3; ! no overhead, not compressed
: 710      0772 3      REC_ADDR = .IRAB[IRB$L_RBF];
: 711      0773 3      END
: 712      0774 2      ELSE
: 713      0775 3      BEGIN
: 714      0776 3      AP = 0;
: 715      0777 3      REC_ADDR = .ADDRESS;
: 716      0778 3
: 717      0779 3      ! In prologue 3 version files, if the key is compressed, it must be
: 718      0780 3      ! rebuilt. Make sure that the last non-compressed pointer, is before
: 719      0781 3      ! the record we are looking at.
: 720      0782 3
: 721      0783 3
: 722      0784 3      IF .IDX_DFN[IDX$V_KEY_COMP]
: 723      0785 3      THEN
: 724      0786 4      BEGIN
: 725      0787 4
: 726      0788 4      IF .IRAB[IRB$L_LST_NCMP] GTRU .ADDRESS
: 727      0789 4      THEN
: 728      0790 5      BEGIN
: 729      0791 5
: 730      0792 5      IF .(.ADDRESS + RM$REC_OVHD() + 1) < 0,8> EQLU 0
: 731      0793 5      THEN
: 732      0794 5      IRAB[IRB$L_LST_NCMP] = .ADDRESS
: 733      0795 5      ELSE
: 734      0796 5      IRAB[IRB$L_LST_NCMP] = .BKT_ADDR + BKT$C_OVERHDSZ;
: 735      0797 4      END;
: 736      0798 3      END;
: 737      0799 2      END;
: 738      0800 2
: 739      0801 2      ! We are storing in key buffer 2 the possible key to be inserted at the
: 740      0802 2      ! index level.
: 741      0803 2
: 742      0804 2
: 743      0805 2      RM$RECORD_KEY ( KEYBUF_ADDR(2) );
: 744      0806 2
: 745      0807 2      RETURN;
: 746      0808 2
: 747      0809 1      END;
```

```

.TITLE RM3SPLUDR
.IDENT \V04-000\

.EXTRN RM$MOVE, RM$RECORD_VBN
.EXTRN RM$RECORD_KEY, RM$REC_OVHD
.EXTRN RM$VBN_SIZE, RM$COMPARE_KEY
.EXTRN RM$COMPARE_REC, RM$GETNEXT_REC

.PSECT RM$RMS3,NOWRT, GBL, PIC,2
```

```
0850 8F BB 00000 RM$MOVE_KEY:
      PUSH  #M<R4,R6,R11>
```

; 0711

RM3SPLUDR
V04-000

RMSMOVE_KEY

J 16
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

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[RMS.SRC]RM3SPLUDR.B32;1

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56	48	A9	5B	50	D0	00004	MOVL	R0, R11	:		
			56	55	C2	00007	SUBL2	BKf_ADDR, R6	:	0766	
			10	00	ED	0000A	CMPZV	#0, #16, 72(IRAB), R6	:		
				0E	12	00010	BNEQ	1\$:		
		09	44	A9	03	E1	00012	BBC	#3, 68(IRAB), 1\$:	0768
			5C	03	D0	00017	MOVL	#3, AP	:	0771	
			56	58	A9	D0	0001A	MOVL	88(IRAB), REC_ADDR	:	0772
				26	11	0001E	BRB	3\$:	0766	
				5C	D4	00020	CLRL	AP	:	0776	
			56	5B	D0	00022	MOVL	ADDRESS, REC_ADDR	:	C777	
		1C	1C	A7	06	E1	00025	BBC	#6, 28(IDY_DFN), 3\$:	0784
			54	0098	C9	9E	0002A	MOVAB	152(IRAB), R4	:	0788
			5B		64	D1	0002F	CMPL	(R4), ADDRESS	:	
					12	1B	00032	BLEQU	3\$:	
				0000G	30	00034	BSBW	RM\$REC_OVHD	:	0792	
				01	A04B	95	00037	TSTB	1(R0)[ADDRESS]	:	
					05	12	0003B	BNEQ	2\$:	
			64		5B	D0	0003D	MOVL	ADDRESS, (R4)	:	0794
					04	11	00040	BRB	3\$:	
			64	0E	A5	9E	00042	MOVAB	14(R5), (R4)	:	0796
			50	00B4	CA	3C	00046	MOVZWL	180(IFAB), R0	:	0805
				60	B940	9F	0004B	PUSHAB	@96(IRAB)[R0]	:	
				0000G	30	0004F	BSBW	RM\$RECORD_KEY	:		
			5E		04	C0	00052	ADDL2	#4, SP	:	
				0850	8F	BA	00055	POPR	#^M<R4,R6,R11>	:	0809
					05	00059	RSB		:		

: Routine Size: 90 bytes, Routine Base: RM\$RMS3 + 0000

: 748 0810 1

RMSBUILD_KEY

```
750 0811 1 %SBTTL 'RMSBUILD_KEY'
751 0812 1 ROUTINE RMSBUILD_KEY (ADDRESS, KEYBUF) : RL$BUILD_KEY NOVALUE =
752 0813 1
753 0814 1 !++
754 0815 1
755 0816 1 FUNCTIONAL DESCRIPTION:
756 0817 1
757 0818 1 This routine builds a compressed key from the record into the
758 0819 1 given buffer, knowing that the front characters are valid from
759 0820 1 the previous expansion.
760 0821 1
761 0822 1 CALLING SEQUENCE:
762 0823 1 bsbw rmsbuild_key (address, keybuf)
763 0824 1
764 0825 1 INPUT PARAMETERS:
765 0826 1 - address in bucket which points to key compression overhead
766 0827 1 - key output buffer
767 0828 1
768 0829 1 IMPLICIT INPUTS:
769 0830 1 IDX_DFN - index definition for key size
770 0831 1
771 0832 1 OUTPUT PARAMETERS:
772 0833 1 none
773 0834 1
774 0835 1 IMPLICIT OUTPUTS:
775 0836 1 key is moved into appropriate key buffer
776 0837 1
777 0838 1 ROUTINE VALUE:
778 0839 1 none
779 0840 1
780 0841 1 SIDE EFFECTS:
781 0842 1 key is moved into appropriate key buffer
782 0843 1
783 0844 1 --
784 0845 1
785 0846 2 BEGIN
786 0847 2
787 0848 2 EXTERNAL REGISTER
788 0849 2 R_IDX_DFN_STR;
789 0850 2
790 0851 2 LOCAL
791 0852 2 TRUN_CHAR,
792 0853 2 LENGTH;
793 0854 2
794 0855 2 BIND
795 0856 2 REC_KEY = ADDRESS : REF BBLOCK;
796 0857 2
797 0858 2 MACRO
798 0859 2 KEY_LEN = 0,0,8,0 %;
799 0860 2 CMP_CNT = 1,0,8,0 %;
800 0861 2
801 0862 2 KEYBUF = .KEYBUF + .REC_KEY[ CMP_CNT ]; ! skip characters already moved
802 0863 2 TRUN_CHAR = .REC_KEY + .REC_KEY[ KEY_LEN ] + 1;
803 0864 2 LENGTH = .IDX_DFN[ IDX$B_KEYSZ ] - .REC_KEY[ CMP_CNT ];
804 0865 2 CH$COPY ( .REC_KEY[ KEY_LEN ], .REC_KEY + 2,
805 0866 2 ..TRUN_CHAR,
806 0867 2 .LENGTH, .KEYBUF );
```

RM3SPLUDR
V04-000

RM\$BUILD_KEY

L 16
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

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[RMS.SRC]RM3SPLUDR.B32;1

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: 807
: 808
: 809
: 810
0868 2
0869 2 RETURN;
0870 2
0871 1 END;

				3E	BB	00000	RM\$BUILD_KEY:		
							PUSHR	#^M<R1,R2,R3,R4,R5>	: 0812
							MOVL	REC KEY, R1	: 0862
							MOVZBL	1(RT), R0	:
							ADDL2	R0, KEYBUF	:
							MOVZBL	(R1), R0	: 0863
							MOVAB	1(R1)[R0], TRUN_CHAR	:
							MOVZBL	32(IDX_DFN), LENGTH	: 0864
							MOVZBL	1(R1), R0	:
							SUBL2	R0, LENGTH	:
							MOVZBL	(R1), R0	: 0865
							MOVCS	R0, 2(R1), (TRUN_CHAR), LENGTH, @KEYBUF	: 0867
							POPR	#^M<R1,R2,R3,R4,R5>	: 0871
							RSB		:

51 18 AE D0 00002
50 01 A1 9A 00006
1C AE 50 C0 0000A
50 61 9A 0000E
53 01 A140 9E 00011
52 20 A7 9A 00016
50 01 A1 9A 0001A
52 50 C2 0001E
50 61 9A 00021
52 50 2C 00024
1C BE 0002A
3E BA 0002C
05 0002E

52 63 02 A1

; Routine Size: 47 bytes, Routine Base: RM\$RMS3 + 005A

: 811 0872 1

RMSSPLIT_UDR

```
813 0873 1 %SBTTL 'RMSSPLIT_UDR'
814 0874 1 GLOBAL ROUTINE RMSSPLIT_UDR : RL$RABREG_4567 NOVALUE =
815 0875 1
816 0876 1 ++
817 0877 1
818 0878 1 FUNCTIONAL DESCRIPTION:
819 0879 1
820 0880 1
821 0881 1 CALLING SEQUENCE:
822 0882 1 BSBW RMSSPLIT_UDR()
823 0883 1
824 0884 1 INPUT PARAMETERS:
825 0885 1 none
826 0886 1
827 0887 1 IMPLICIT INPUTS:
828 0888 1 BDB pointer, BUFFER pointer, REC_ADDR = point of insert, IDX_DFN
829 0889 1 in IRAB -- curbdb, associated w/ bdb and bkt_addr
830 0890 1 pos_ins corresponding to rec_addr
831 0891 1 in RAB -- rsz of record
832 0892 1 in IFAB -- rfm
833 0893 1 BKT$B_NXTRECID = 0 in original bucket signals that this is
834 0894 1 a split due to a lack of id's in the bucket
835 0895 1
836 0896 1 OUTPUT PARAMETERS:
837 0897 1 none
838 0898 1
839 0899 1 IMPLICIT OUTPUTS:
840 0900 1 in IRAB --
841 0901 1 if 2 bkt split --
842 0902 1 IRB$W_SPLIT, offset to split point
843 0903 1 IRB$V_REC_W_LO -- set if split point is pos_insert and
844 0904 1 record goes w/ lo set
845 0905 1 new high key for original bucket in keybuffer 2
846 0906 1 number of new buckets = 1
847 0907 1 if original bucket was all rrv's, set IRB$V_EMPTY_BKT flag
848 0908 1 if new bucket is a continuation bkt., set IRB$V_CONT_BKT flag
849 0909 1 if 3 bkt split --
850 0910 1 same as above w/ these changes:
851 0911 1 IRB$W_SPLIT_1, offset to second split point
852 0912 1 number of new buckets = 2
853 0913 1 if right bucket is a continuation bkt, set IRB$V_CONT_R flag
854 0914 1 if 4 bkt split --
855 0915 1 same as above w/ these changes:
856 0916 1 IRB$W_SPLIT_2, offset to third split point
857 0917 1 number of new buckets = 3
858 0918 1
859 0919 1 ROUTINE VALUE:
860 0920 1 rmssuc
861 0921 1
862 0922 1 SIDE EFFECTS:
863 0923 1 AP is clobbered
864 0924 1
865 0925 1 --
866 0926 1
867 0927 2 BEGIN
868 0928 2
869 0929 2 EXTERNAL REGISTER
```

```
870 0930 2 COMMON RAB_STR,  
871 0931 2 R_REC_ADDR_STR,  
872 0932 2 R_IDX_DFN_STR,  
873 0933 2 COMMON_IO_STR;  
874 0934 2  
875 0935 2 LOCAL  
876 0936 2 SAVE_REC_W_LO,  
877 0937 2 NUM_RRVs,  
878 0938 2 POS_INSERT,  
879 0939 2 EOB,  
880 0940 2 RRV,  
881 0941 2 RHS,  
882 0942 2 LHS,  
883 0943 2 LAST : REF BBLOCK,  
884 0944 2 LAST_DIFF,  
885 0945 2 BKTSIZE,  
886 0946 2 REC_SIZE,  
887 0947 2 DIFFERENCE;  
888 0948 2  
889 0949 2 MACRO  
890 0950 2 NEED_RRV = NUM_RRVs<0,16> %,  
891 0951 2 NOT_NEED_RRV = -NUM_RRVs<16,16> %;  
892 0952 2  
893 0953 2 LABEL  
894 0954 2 DO_IT,  
895 0955 2 HALF,  
896 0956 2 NEXT;  
897 0957 2  
898 0958 2 DO_IT :  
899 0959 2  
900 0960 3 BEGIN  
901 0961 3  
902 0962 3 ! define a block so that we can have some common checks before returning  
903 0963 3 ! successfully  
904 0964 3 !  
905 0965 3 HALF :  
906 0966 3  
907 0967 4 BEGIN  
908 0968 4  
909 0969 4 !  
910 0970 4 ! define a block so that we can simulate a go-to (naughty, naughty)  
911 0971 4 ! if we have decided that we are positioning at the end of the bucket  
912 0972 4 ! & we're in somewhat of an ascending order, where the last record  
913 0973 4 ! inserted is a duplicate of the new record, skip over the 50-50 code  
914 0974 4 ! and go to the code to take duplicates into account  
915 0975 4 !  
916 0976 4 ! scan 1 -- calculate  
917 0977 4 ! size of existing rrv's and total number of rrv's needed to move the whole  
918 0978 4 ! bucket out ( worst case) as a side effect, adjust eob ptr to pt to the  
919 0979 4 ! rrv's instead of freespace assume not empty bucket until showed otherwise  
920 0980 4 !  
921 0981 4  
922 0982 4 [RAB[IRBSV_EMPTY_BKT] = 0;  
923 0983 4  
924 0984 4 ! new rec is tried 1st w/ hi set, then w/ lo set  
925 0985 4  
926 0986 4 [RAB[IRBSV_REC_W_LO] = 0;
```

```

927 0987 4 IRAB[IRBSV_NEW_BKTS] = 1; ! assume 2-bkt split until showed otherwise
928 0988 4 NUM_RRVs = 0; ! this zeroes NEED_RRV and NOT_NEED_RRV
929 0989 4 POS_INSERT = .REC_ADDR;
930 0990 4 REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
931 0991 4 EOB = .BKT_ADDR + .BKT_ADDR[BKT$W_FREESPACE];
932 0992 4 LAST = 0;
933 0993 4
934 0994 4 DO
935 0995 5 BEGIN
936 0996 5
937 0997 5 BUILTIN
938 0998 5 AP;
939 0999 5
940 1000 5 IF .REC_ADDR[IRC$V_RRV]
941 1001 5 THEN
942 1002 5 EXITLOOP;
943 1003 5
944 1004 5 AP = 3;
945 1005 5
946 1006 5 IF .BDB[BDB$L_VBN] EQLU RM$RECORD_VBN()
947 1007 5 THEN
948 1008 5 NEED_RRV = .NEED_RRV + 1
949 1009 5
950 1010 5 ! the records not requiring rrv's are counted also because in the
951 1011 5 ! case where we're splitting due to lack of id's, the lhs side will
952 1012 5 ! fit with the new record if any of the record being moved to the
953 1013 5 ! new bucket doesn't require an rrv. this will be checked when we
954 1014 5 ! check to see if the lhs will fit after the first point that the
955 1015 5 ! rhs fits.
956 1016 5
957 1017 5 ELSE
958 1018 5 NOT_NEED_RRV = .NOT_NEED_RRV + 1;
959 1019 5
960 1020 5 LAST = .REC_ADDR;
961 1021 5 RM$GETNEXT_REC()
962 1022 5 END
963 1023 4 UNTIL .REC_ADDR GEQU .EOB;
964 1024 4
965 1025 4 ! set split_2 and split_1 to be eob, so if there's less than 3 new buckets
966 1026 4 ! bkt_spl can use the value w/o having to recalculate it also set up the
967 1027 4 ! bucket size and the record size
968 1028 4
969 1029 4 IRAB[IRBSW_SPLIT_1] = IRAB[IRBSW_SPLIT_2] = .REC_ADDR - .BKT_ADDR;
970 1030 4 BKTSIZE = .IDX_DFN[IDX$B_DATABKTSZ]*512 - BKT$C_OVERHDSZ - 1;
971 1031 4
972 1032 4 REC_SIZE = .RAB[RAB$W_RSZ] + IRC$C_FIXOVHDSZ;
973 1033 4
974 1034 4 IF .IFAB[IFB$B_RFMORG] NEQ FAB$C_FIX
975 1035 4 THEN
976 1036 4 REC_SIZE = .REC_SIZE + 2;
977 1037 4
978 1038 4 ! if this is an update, may have to count in an rrv for the existing record
979 1039 4
980 1040 4 IF .IRAB[IRBSV_UPDATE]
981 1041 4 THEN
982 1042 5 BEGIN
983 1043 5
```

```
: 984      1044 5      IF .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 985      1045 5      THEN
: 986      1046 5          NEED_RRV = .NEED_RRV + 1
: 987      1047 5
: 988      1048 4      END;
: 989      1049 4
: 990      1050 4      RRV = .EOB - .REC_ADDR;          ! size of existing rrv's
: 991      1051 4      EOB = .REC_ADDR;              ! adjust eob
: 992      1052 4
: 993      1053 4      ! special case it, if the bucket was all rrv's
: 994      1054 4      !
: 995      1055 4
: 996      1056 4      IF .REC_ADDR EQLU .BKT_ADDR + BKT$C_OVERHDSZ
: 997      1057 4      THEN
: 998      1058 5          BEGIN
: 999      1059 5              ! bkt is all rrv's yet the record wouldn't fit so we need to
1000      1060 5              ! allocate another bkt ( 2 bkt split) yet special case it so as not
1001      1061 5              ! to make another idx entry only to update the existing one by
1002      1062 5              ! setting empty bucket flag
1003      1063 5              !
1004      1064 5              IRAB[IRB$W_SPLIT] = .REC_ADDR - .BKT_ADDR;
1005      1065 5              LEAVE DO_IT
1006      1066 5
1007      1067 5
1008      1068 4          END;                      ! { of special case an all-rrv bucket }
1009      1069 4
1010      1070 4      +
1011      1071 4      ! special case -- if we can detect a possible ascending order to these
1012      1072 4      ! records it probably will be better to do a straight point of insert split
1013      1073 4      ! this would put the new record in a bucket all by itself.
1014      1074 4      ! do this kind of split if and only if all the following conditions are met:
1015      1075 4      ! 1) the record is being inserted at the end of bucket
1016      1076 4      ! 2) the last record physically in the bkt is the last record to have
1017      1077 4      !     been inserted
1018      1078 4      ! 3) the last record and the new record do not have duplicate key values
1019      1079 4
1020      1080 4      ! note that if they are duplicates, we can still make an optimization by
1021      1081 4      ! skipping the 50-50 split code
1022      1082 4
1023      1083 4      ! note that last cannot be zero, since if it were we
1024      1084 4      ! would have an all rrv bkt
1025      1085 4      -
1026      1086 4
1027      1087 4      IF .POS_INSERT EQLU .REC_ADDR
1028      1088 4          AND
1029      1089 5          (((.LAST[IRC$B_ID] + 1) AND %X'FF') EQLU .BKT_ADDR[BKT$B_NXTRECID])
1030      1090 4      THEN
1031      1091 5          BEGIN
1032      1092 5              REC_ADDR = .LAST;
1033      1093 5
1034      1094 5              IF RM$COMPARE_REC(KEYBUF_ADDR(3), .IDX_DFN[IDX$B_KEYSZ], 0)
1035      1095 5              THEN
1036      1096 6                  BEGIN
1037      1097 6
1038      1098 6                  ! since we have detected a possible ascending order in the input
1039      1099 6                  ! let's try to optimize a little and split at the point of insert
1040      1100 6                  ! send the record by itself into the new bucket have to set up the
```



```
1041 1101 6      ! key value and the split point and that's it
1042 1102 6
1043 1103 6      RMSMOVE KEY(.REC_ADDR, .REC_ADDR);
1044 1104 6      IRAB[IRBSW_SPLIT] = .IRAB[IRBSW_POS_INS];
1045 1105 6      LEAVE DO_IT;
1046 1106 6
1047 1107 6      END
1048 1108 5      ELSE
1049 1109 5          LEAVE HALF;
1050 1110 5
1051 1111 5      ! ( end of trying to special case insertion of records in ascending
1052 1112 5      ! order )
1053 1113 5
1054 1114 4      END;
1055 1115 4
1056 1116 4      REC_ADDR = .BKT_ADDR + BKTSC_OVERHDSZ;
1057 1117 4      LAST_DIFF = 'x'7FFFFFFF;
1058 1118 4      LAST = 0;
1059 1119 4      SAVE_REC_W_LO = 0;
1060 1120 4
1061 1121 4      ! start from the beginning of the bucket and scan rightward. first find the
1062 1122 4      ! 1st place the rhs will fit in 1 bkt then, as long as the lhs will fit in
1063 1123 4      ! a bkt, try to find an optimal point if there is no point where the rhs
1064 1124 4      ! and lhs will both fit we can't do a 2-bkt split and this case will fall
1065 1125 4      ! out
1066 1126 4
1067 1127 4
1068 1128 4      WHILE 1
1069 1129 4      DO
1070 1130 5          BEGIN
1071 1131 5          RHS = .EOB - .REC_ADDR;
1072 1132 5
1073 1133 5          IF .REC_ADDR LEQU .POS_INSERT
1074 1134 5              AND
1075 1135 5              NOT .IRAB[IRBSV_REC_W_LO]
1076 1136 5          THEN
1077 1137 5              RHS = .RHS + .REC_SIZE;
1078 1138 5
1079 1139 5          ! the right hand side fits if there is enough room and there are id's
1080 1140 5          ! available. id's are always available in the new bucket in the update
1081 1141 5          ! situation, or if we're leaving at least 1 record behind in the old
1082 1142 5          ! bucket. note that nxtrecid is always zeroed if this is a split due to
1083 1143 5          ! lack of id's.
1084 1144 5
1085 1145 5
1086 1146 5          IF .RHS LSSU .BKTSIZE
1087 1147 5              AND
1088 1148 6              (.BKT_ADDR[BKTSB_NXTRECID] NEQ 0
1089 1149 6              OR
1090 1150 6              .IRAB[IRBSV_UPDATE]
1091 1151 6              OR
1092 1152 7              .REC_ADDR NEQA (.BKT_ADDR + BKTSC_OVERHDSZ)
1093 1153 6              OR
1094 1154 6              .IRAB[IRBSV_REC_W_LO])
1095 1155 5          THEN
1096 1156 6              BEGIN
1097 1157 6                  LHS = .REC_ADDR - (.BKT_ADDR + BKTSC_OVERHDSZ);
```

```

: 1098      1158 6
: 1099      1159 6
: 1100      1160 6
: 1101      1161 6
: 1102      1162 6
: 1103      1163 6
: 1104      1164 6
: 1105      1165 6
: 1106      1166 6
: 1107      1167 6
: 1108      1168 6
: 1109      1169 6
: 1110      1170 6
: 1111      1171 6
: 1112      1172 6
: 1113      1173 6
: 1114      1174 6
: 1115      1175 6
: 1116      1176 6
: 1117      1177 6
: 1118      1178 6
: 1119      1179 6
: 1120      1180 7
: 1121      1181 7
: 1122      1182 7
: 1123      1183 7
: 1124      1184 7
: 1125      1185 7
: 1126      1186 7
: 1127      1187 6
: 1128      1188 7
: 1129      1189 7
: 1130      1190 7
: 1131      1191 7
: 1132      1192 7
: 1133      1193 7
: 1134      1194 7
: 1135      1195 7
: 1136      1196 7
: 1137      1197 7
: 1138      1198 7
: 1139      1199 7
: 1140      1200 7
: 1141      1201 7
: 1142      1202 7
: 1143      1203 7
: 1144      1204 8
: 1145      1205 8
: 1146      1206 8
: 1147      1207 8
: 1148      1208 8
: 1149      1209 8
: 1150      1210 8
: 1151      1211 8
: 1152      1212 8
: 1153      1213 8
: 1154      1214 8

IF .REC_ADDR GEQU .POS_INSERT
AND
.IRAB[IRBSV_REC_W_LO]
THEN
LHS = .LHS + .REC_SIZE;

! will lhs fit ? lhs doesn't fit if there is no space in the
! bucket, or if there won't be any id's available in the bucket.
! if not & if there is no previous point at which it fit, goto 3-bkt
! split code if there is a previous place where we could have had a
! 2-bkt split, use it

IF .LHS + .RRV + (7*.NEED_RRV) GTRU .BKTSIZE
! id's will be available in the original bucket if we aren't
! out of id's to begin with, this is an update, any record
! being moved out doesn't need an rrv, or the new record is
! going in the new bucket
OR
(.BKT_ADDR[BKTSB_NXTRECID] EQL 0
AND
NOT .IRAB[IRBSV_UPDATED]
AND
.NOT_NEED_RRV EQL 0
AND
.IRAB[IRBSV_REC_W_LO])
THEN
BEGIN
IF .LAST EQL 0
THEN
EXITLOOP;

REC_ADDR = .LAST;

IF NOT .SAVE_REC_W_LO
THEN
IRAB[IRBSV_REC_W_LO] = 0;

! 2 bkt split is possible rec_addr points to the most
! optimal place since we had to back up, reset last to point
! to the record immediately before the split point
BEGIN
LOCAL
TMP;

TMP = .REC_ADDR;
REC_ADDR = .BKT_ADDR + BKTSC_OVERHDSZ;
LAST = .REC_ADDR;

WHILE .REC_ADDR NEQU .TMP
DO
```

```
: 1155      1215 9      BEGIN
: 1156      1216 9      LAST = .REC_ADDR;
: 1157      1217 9      RMSGETNEXT_REC();
: 1158      1218 8      END;
: 1159      1219 8
: 1160      1220 7      END;
: 1161      1221 7      RMSMOVE KEY(.LAST, .REC_ADDR);
: 1162      1222 7      IRAB[IRB$W_SPLIT] = .REC_ADDR - .BKT_ADDR;
: 1163      1223 7
: 1164      1224 7      ! treat another exception case of the new record going off into
: 1165      1225 7      ! a cont. bkt all by itself
: 1166      1226 7
: 1167      1227 7
: 1168      1228 7      IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_POS_INS]
: 1169      1229 7      THEN
: 1170      1230 7
: 1171      1231 7      IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_SPLIT_1]
: 1172      1232 7      THEN
: 1173      1233 7
: 1174      1234 7      IF NOT .IRAB[IRB$V_REC_W_LO]
: 1175      1235 7      THEN
: 1176      1236 8      BEGIN
: 1177      1237 8
: 1178      1238 8      BUILTIN
: 1179      1239 8      AP;
: 1180      1240 8
: 1181      1241 8      AP = 3;
: 1182      1242 8
: 1183      1243 8      IF NOT RMSCOMPARE KEY(KEYBUF_ADDR(2),
: 1184      1244 8      KEYBUF_ADDR(3),
: 1185      1245 8      .IDX_DFN[IDX$B_KEYSZ])
: 1186      1246 8      THEN
: 1187      1247 8      IRAB[IRB$V_CONT_BKT] = 1;
: 1188      1248 8
: 1189      1249 7      END;
: 1190      1250 7
: 1191      1251 7      LEAVE DO_IT
: 1192      1252 7
: 1193      1253 6      END;
: 1194      1254 6      ! { end of lhs doesn't fit anymore }
: 1195      1255 6      ! lhs fits also, calculate the magic ratio
: 1196      1256 6
: 1197      1257 8      DIFFERENCE = (.LHS*.BKTSIZE) - (.RHS*(.BKTSIZE - (7*.NEED_RRV) -
: 1198      1258 6      .RRV));
: 1199      1259 6
: 1200      1260 6      IF .DIFFERENCE GEQ 0
: 1201      1261 6      THEN
: 1202      1262 7      BEGIN
: 1203      1263 7
: 1204      1264 7      ! found the 1st point at which the magic ratio is positive
: 1205      1265 7      ! was the last point more optimal, if so use it
: 1206      1266 7
: 1207      1267 7
: 1208      1268 7      IF ABS(.DIFFERENCE) GTRU ABS(.LAST_DIFF)
: 1209      1269 7      THEN
: 1210      1270 8      BEGIN
: 1211      1271 8
```

```

: 1212      1272  8      IF .REC_ADDR EQLU .LAST
: 1213      1273  8      THEN
: 1214      1274  8          IRAB[IRBSV_REC_W_LO] = 0
: 1215      1275  8      ELSE
: 1216      1276  9          (REC_ADDR = .LAST;
: 1217      1277  9              IF .REC_ADDR LSSU .POS_INSERT
: 1218      1278  9              THEN
: 1219      1279  9                  IRAB[IRBSV_REC_W_LO] = 0);
: 1220      1280  8
: 1221      1281  8
: 1222      1282  8          LAST = 0;
: 1223      1283  7          END;
: 1224      1284  7
: 1225      1285  7      ! 2-bkt split is possible rec_addr points to the most
: 1226      1286  7      ! optimal place
: 1227      1287  7
: 1228      1288  7
: 1229      1289  7      IF .LAST EQL 0
: 1230      1290  7      THEN      ! just backed up rec_addr, need to recalc last
: 1231      1291  8          BEGIN
: 1232      1292  8
: 1233      1293  8              LOCAL
: 1234      1294  8                  TMP;
: 1235      1295  8
: 1236      1296  8              TMP = .REC_ADDR;
: 1237      1297  8              REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
: 1238      1298  8              LAST = .REC_ADDR;
: 1239      1299  8
: 1240      1300  8              WHILE .REC_ADDR NEQU .TMP
: 1241      1301  8                  DO
: 1242      1302  9                      BEGIN
: 1243      1303  9                          LAST = .REC_ADDR;
: 1244      1304  9                          RMSGETNEXT_REC();
: 1245      1305  8                          END;
: 1246      1306  8
: 1247      1307  7                  END;
: 1248      1308  7
: 1249      1309  7      RMSMOVE KEY(.LAST, .REC_ADDR);
: 1250      1310  7      IRAB[IRBSW_SPLIT] = .REC_ADDR - .BKT_ADDR;
: 1251      1311  7
: 1252      1312  7      ! treat another exception case of the new record going off into
: 1253      1313  7      ! a cont. bkt all by itself
: 1254      1314  7
: 1255      1315  7
: 1256      1316  7      IF .IRAB[IRBSW_SPLIT] EQLU .IRAB[IRBSW_POS_INS]
: 1257      1317  7      THEN
: 1258      1318  7
: 1259      1319  7          IF .IRAB[IRBSW_SPLIT] EQLU .IRAB[IRBSW_SPLIT_1]
: 1260      1320  7          THEN
: 1261      1321  7
: 1262      1322  7              IF NOT .IRAB[IRBSV_REC_W_LO]
: 1263      1323  7              THEN
: 1264      1324  8                  BEGIN
: 1265      1325  8
: 1266      1326  8                      BUILTIN
: 1267      1327  8                          AP;
: 1268      1328  8
```

```

: 1269      1329  8      AP = 3;
: 1270      1330  8
: 1271      1331  8      IF NOT RM$COMPARE_KEY(KEYBUF_ADDR(2),
: 1272      1332  8          KEYBUF_ADDR(3),
: 1273      1333  8          .IDX_DFN[IDX$B_KEY$Z])
: 1274      1334  8      THEN
: 1275      1335  8          IRAB[IRB$V_CONT_BKT] = 1;
: 1276      1336  8
: 1277      1337  7      END;
: 1278      1338  7
: 1279      1339  7          LEAVE DO_IT
: 1280      1340  7
: 1281      1341  6          END;
: 1282      1342  6
: 1283      1343  6      ! the magic ratio isn't positive yet, so save all the context and
: 1284      1344  6      ! move on to the next record
: 1285      1345  6
: 1286      1346  6      LAST_DIFF = .DIFFERENCE;
: 1287      1347  6      LAST = .REC_ADDR;
: 1288      1348  6
: 1289      1349  6      IF .IRAB[IRB$V_REC_W_LO]
: 1290      1350  6      THEN
: 1291      1351  6          SAVE_REC_W_LO = 1;
: 1292      1352  6
: 1293      1353  5      END;
: 1294      1354  5          ! { end of rhs fits, is this a good point? }
: 1295      1355  5      ! go on to the next record
: 1296      1356  5
: 1297      1357  5      NEXT :
: 1298      1358  6      BEGIN
: 1299      1359  6
: 1300      1360  6      IF .REC_ADDR EQLU .POS_INSERT
: 1301      1361  6      AND
: 1302      1362  6      NOT .IRAB[IRB$V_REC_W_LO]
: 1303      1363  6      THEN
: 1304      1364  7          BEGIN
: 1305      1365  7
: 1306      1366  7      ! if this is an update and we pass the record, check to see if it
: 1307      1367  7      ! needed an rrv
: 1308      1368  7
: 1309      1369  7      IF .IRAB[IRB$V_UPDATE]
: 1310      1370  7      THEN
: 1311      1371  8          BEGIN
: 1312      1372  8
: 1313      1373  8          IF .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 1314      1374  8          THEN
: 1315      1375  8              NEED_RRV = .NEED_RRV - 1;
: 1316      1376  8
: 1317      1377  7          END;
: 1318      1378  7
: 1319      1379  7      IRAB[IRB$V_REC_W_LO] = 1;
: 1320      1380  7      RM$MOVE_KEY(.REC_ADDR, .REC_ADDR);
: 1321      1381  7
: 1322      1382  7      IF .REC_ADDR EQLU .EOB
: 1323      1383  7      THEN
: 1324      1384  7          LEAVE NEXT
: 1325      1385  7
```

```
: 1326      1386 7      ELSE IF RMSCOMPARE_REC(KEYBUF_ADDR(2), .IDX_DFN[IDX$B_KEYSZ], 0)
: 1327      1387 7      THEN
: 1328      1388 7          LEAVE NEXT;
: 1329      1389 7
: 1330      1390 6      END;          ! ( end of at position for insert for the 1st time )
: 1331      1391 6
: 1332      1392 6      ! fool move key a little by always clearing rec_w_lo to always get the
: 1333      1393 6      ! key associated w/ the record at pos_ins. (I think it is the key of
: 1334      1394 6      ! the record we are pointing to, not the one at pos_ins...)
: 1335      1395 6
: 1336      1396 7      BEGIN
: 1337      1397 7
: 1338      1398 7      LOCAL
: 1339      1399 7          TMP : BYTE;
: 1340      1400 7
: 1341      1401 7      TMP = .IRAB[IRB$B_SPL_BITS];
: 1342      1402 7      IRAB[IRB$V_REC_W_LO] = 0;
: 1343      1403 7      RMSMOVE KEY(.REC_ADDR, .REC_ADDR);
: 1344      1404 7      IRAB[IRB$B_SPL_BITS] = .TMP
: 1345      1405 6      END;
: 1346      1406 6
: 1347      1407 6      DO
: 1348      1408 7          BEGIN
: 1349      1409 7
: 1350      1410 7          BUILTIN
: 1351      1411 7              AP;
: 1352      1412 7
: 1353      1413 7          IF .REC_ADDR EQLU .EOB
: 1354      1414 7          THEN
: 1355      1415 7              EXITLOOP;
: 1356      1416 7
: 1357      1417 7          AP = 3;
: 1358      1418 7
: 1359      1419 7          IF .BDB[BDB$L_VBN] EQLU RMSRECORD_VBN()
: 1360      1420 7          THEN
: 1361      1421 7              NEED_RRV = .NEED_RRV - 1
: 1362      1422 7          ELSE
: 1363      1423 7              NOT_NEED_RRV = .NOT_NEED_RRV - 1;
: 1364      1424 7
: 1365      1425 7          RMSGETNEXT_REC();
: 1366      1426 7
: 1367      1427 7          IF .REC_ADDR EQLU .EOB
: 1368      1428 7          THEN
: 1369      1429 7              EXITLOOP;
: 1370      1430 7
: 1371      1431 7          END
: 1372      1432 7
: 1373      1433 7      ! compare_rec returns 0 if a match
: 1374      1434 7
: 1375      1435 6      UNTIL RMSCOMPARE_REC(KEYBUF_ADDR(2), .IDX_DFN[IDX$B_KEYSZ], 0);
: 1376      1436 6
: 1377      1437 6      ! if the key compares brought us up to the pos of insert, see if the
: 1378      1438 6      ! key of the new record matches. if it does, have to include it w/ the
: 1379      1439 6      ! lhs
: 1380      1440 6
: 1381      1441 6
: 1382      1442 6      IF .REC_ADDR EQLU .POS_INSERT
```

```
: 1383      1443  6      THEN
: 1384      1444  7          BEGIN
: 1385      1445  7
: 1386      1446  7          BUILTIN
: 1387      1447  7              AP;
: 1388      1448  7
: 1389      1449  7          AP = 3;
: 1390      1450  7
: 1391      1451  7          IF NOT RM$COMPARE_KEY(KEYBUF_ADDR(2),
: 1392      1452  7              KEYBUF_ADDR(3),
: 1393      1453  7              .IDX_DFRC[IDX$B_KEYSZ])
: 1394      1454  7              THEN
: 1395      1455  8                  BEGIN
: 1396      1456  8                      IRAB[IRB$V_REC_W_LO] = 1;
: 1397      1457  8
: 1398      1458  8                      IF .IRAB[IRB$V_UPDATE]
: 1399      1459  8                          AND
: 1400      1460  8                          .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 1401      1461  8                          THEN
: 1402      1462  8                              NEED_RRV = .NEED_RRV - 1;
: 1403      1463  7                              END;
: 1404      1464  7
: 1405      1465  6                      END;
: 1406      1466  6
: 1407      1467  6          IF .REC_ADDR GTRU .POS_INSERT
: 1408      1468  6              THEN
: 1409      1469  7                  BEGIN
: 1410      1470  7                      IRAB[IRB$V_REC_W_LO] - 1;
: 1411      1471  7
: 1412      1472  7                      IF .IRAB[IRB$V_UPDATE]
: 1413      1473  7                          AND
: 1414      1474  7                          .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 1415      1475  7                          THEN
: 1416      1476  7                              NEED_RRV = .NEED_RRV - 1;
: 1417      1477  6                              END;
: 1418      1478  6
: 1419      1479  5          END;
: 1420      1480  4          END;
: 1421      1481  4
: 1422      1482  3      END;
: 1423      1483  3
: 1424      1484  3      ! define a new block here so local storage can be redefined
: 1425      1485  3      !
: 1426      1486  4      BEGIN
: 1427      1487  4
: 1428      1488  4      MACRO
: 1429      1489  4          BEG_CHAIN = LHS %,
: 1430      1490  4          END_CHAIN = RHS %,
: 1431      1491  4          NUM_DUPS = NUM_RRVs %,
: 1432      1492  4          DUPS = RRV %;
: 1433      1493  4
: 1434      1494  4      BUILTIN
: 1435      1495  4          AP;
: 1436      1496  4
: 1437      1497  4      ! must be a 3 or 4 bucket split or we detected ascending order and the new
: 1438      1498  4      ! record was a dupe. we'll optimize here to the extent of trying to keep a
: 1439      1499  4      ! dup chain around the new record together and in the middle bucket
```

```
: 1440      1500  4  | note that in all the cases that follow the new record is going into the
: 1441      1501  4  | middle bucket. therefore, the "lhs" will always fit, since it can only
: 1442      1502  4  | get smaller ( or stay the same size, in the degenerate case). also note
: 1443      1503  4  | that in any of these case, the left hand bucket may be empty of data
: 1444      1504  4  | records (have only rrv's in it) if the first split point is at the
: 1445      1505  4  | beginning and all data records get moved out
: 1446      1506  4  |
: 1447      1507  4  | IRAB[IRBSV_NEW_BKTS] = 2;    ! assume 3-bkt split until shown otherwise
: 1448      1508  4  | IRAB[IRBSV_REC_W_LO] = 0;
: 1449      1509  4  |
: 1450      1510  4  | initialize key buffer 2 with the contents of key buffer 3 (the value
: 1451      1511  4  | of the primary key of the record being inserted). This is necessary
: 1452      1512  4  | when the new record is at the beginning of the bucket and is going into
: 1453      1513  4  | a bucket all by itself and there were already 255 records in their
: 1454      1514  4  | original bucket and they all need rrv's therefore they all move into the
: 1455      1515  4  | next bucket. At any rate, that seems to be the only case where key buffer
: 1456      1516  4  | 2 is not correct coming into here and will be set correctly before
: 1457      1517  4  | leaving.
: 1458      1518  4  |
: 1459      1519  4  | RMSMOVE(.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2));
: 1460      1520  4  |
: 1461      1521  4  | ! find beginning and end of this possible dups chain equal to the key value
: 1462      1522  4  | of the record being inserted.
: 1463      1523  4  |
: 1464      1524  4  | REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
: 1465      1525  5  | BEGIN
: 1466      1526  5  |
: 1467      1527  5  | LOCAL
: 1468      1528  5  |     STATUS;
: 1469      1529  5  |
: 1470      1530  5  | WHILE STATUS = RMSCOMPARE_REC(KEYBUF_ADDR(3), .IRAB[IRBSB_KEYSZ], 0)
: 1471      1531  5  | DO
: 1472      1532  6  |     BEGIN
: 1473      1533  6  |         IF .REC_ADDR LSSU .POS_INSERT
: 1474      1534  6  |         THEN
: 1475      1535  6  |             BEGIN
: 1476      1536  7  |                 AP = 0;
: 1477      1537  7  |                 RMSRECORD_KEY(KEYBUF_ADDR(2));
: 1478      1538  7  |                 END;
: 1479      1539  6  |             IF .REC_ADDR EQLU .EOB
: 1480      1540  6  |             OR
: 1481      1541  6  |             .STATUS LSS 0
: 1482      1542  6  |             THEN
: 1483      1543  6  |                 BEGIN
: 1484      1544  7  |                     ! !!!! SPLIT TYPE 3 !!!! no duplicates found for simplicity, do a
: 1485      1545  7  |                     ! 3-bkt split at the point of insert w/ new record in its own
: 1486      1546  7  |                     ! bucket
: 1487      1547  7  |                     !
: 1488      1548  7  |                     ! IRAB[IRBSW_SPLIT] = IRAB[IRBSW_SPLIT_1] = .IRAB[IRBSW_POS_INS];
: 1489      1549  7  |                     ! LEAVE DO_IT
: 1490      1550  7  |                     !
: 1491      1551  7  |                     ! ( end of didn't find a duplicate, put record in its own bucket )
: 1492      1552  7  |                     !
: 1493      1553  7  |                     !
: 1494      1554  7  |                     !
: 1495      1555  7  |                     !
: 1496      1556  6  |     END;
```



```
1497 1557 6
1498 1558 6      RM$GETNEXT_REC();
1499 1559 5      END;                                ! { end of while no duplicate has been found }
1500 1560 5
1501 1561 4      END;                                ! { end of block defining status for while loop }
1502 1562 4
1503 1563 4      ! found the beginning of the dups chain, now find the end
1504 1564 4      !
1505 1565 4      NUM_DUPS = 0;
1506 1566 4      BEG_CHAIN = .REC_ADDR;
1507 1567 4
1508 1568 4      DO
1509 1569 5          BEGIN
1510 1570 5              NUM_DUPS = .NUM_DUPS + 1;
1511 1571 5              RM$GETNEXT_REC();
1512 1572 5
1513 1573 5              IF .REC_ADDR EQLU .EOB
1514 1574 5                  THEN
1515 1575 5                      EXITLOOP;
1516 1576 5
1517 1577 5              END
1518 1578 4      UNTIL RM$COMPARE_REC(KEYBUF_ADDR(3), .IRAB[IRB$B_KEYSZ], 0);
1519 1579 4                      ! compare_rec returns 0 if keys match
1520 1580 4
1521 1581 4      END_CHAIN = .REC_ADDR;
1522 1582 4
1523 1583 4      ! found the beginning and the end of the chain calculate its size if we got
1524 1584 4      ! here via an update, we never called rm$srch_by_key to set dups_seen
1525 1585 4      ! for us. so let us do that now if necessary
1526 1586 4      !
1527 1587 4
1528 1588 4      IF .POS_INSERT GTRU .BEG_CHAIN
1529 1589 4      THEN
1530 1590 4          IRAB[IRB$V_DUPS_SEEN] = 1;
1531 1591 4
1532 1592 4      DUPS = .END_CHAIN - .BEG_CHAIN;
1533 1593 4      DUPS = .DUPS + .REC_SIZE;
1534 1594 4
1535 1595 4      IF .DUPS LSSU .BKTSIZE
1536 1596 4
1537 1597 4          ! if there are 255 dups on a put, there won't be enough id's in the
1538 1598 4          ! new bucket even if there is enough space for them.
1539 1599 4
1540 1600 4          AND
1541 1601 5              (.IRAB[IRB$V_UPDATE]
1542 1602 5              OR
1543 1603 5              .NUM_DUPS < 0, 8> LEQU 254)
1544 1604 4      THEN
1545 1605 5          BEGIN
1546 1606 5
1547 1607 5              !+
1548 1608 5              !!!!! SPLIT TYPE 1 !!!!!
1549 1609 5              ! duplicates found and fortunately, they all fit
1550 1610 5              ! in one bucket so 3-bkt split w/ all of the dups in the middle bucket
1551 1611 5              ! because of the optimization used for dups being inserted "in order"
1552 1612 5              ! this can still be a 2-bkt split if the new record is being inserted
1553 1613 5              ! at the end of the bucket
```

```

: 1554      1614 5      | 22-jan-79 if loa forced us to think that a
: 1555      1615 5      | bkt w/ all dups had to be split ( only on put) be smart and just put
: 1556      1616 5      | new record by itself a better solution would be not to split at all,
: 1557      1617 5      | but at this date it's rather inconceivable
: 1558      1618 5      | 23-jan-79 it's not only loa
: 1559      1619 5      | that can fool us, the bkt might have had a lot of rrv's
: 1560      1620 5      |
: 1561      1621 5      |
: 1562      1622 5      | IRAB[IRBSW_SPLIT] = .BEG_CHAIN - .BKT_ADDR;
: 1563      1623 5      | IRAB[IRBSW_SPLIT_1] = .END_CHAIN - .BKT_ADDR;
: 1564      1624 5      |
: 1565      1625 5      | IF .END_CHAIN EQLU .EOB
: 1566      1626 5      | THEN
: 1567      1627 6      | BEGIN
: 1568      1628 6      | IRAB[IRBSV_NEW_BKTS] = 1;
: 1569      1629 6      |
: 1570      1630 7      | IF .BEG_CHAIN EQLU (.BKT_ADDR + BKTSC_OVERHDSZ)
: 1571      1631 6      | THEN
: 1572      1632 7      | BEGIN
: 1573      1633 7      | IRAB[IRBSW_SPLIT_1] = .IRAB[IRBSW_SPLIT_2];
: 1574      1634 7      | IRAB[IRBSW_SPLIT] = .IRAB[IRBSW_POS_INS];
: 1575      1635 7      | IRAB[IRBSV_CONT_BKT] = 1;
: 1576      1636 7      | END
: 1577      1637 7      |
: 1578      1638 6      | END
: 1579      1639 5      | ELSE
: 1580      1640 6      | BEGIN
: 1581      1641 6      |
: 1582      1642 6      | IF .IRAB[IRBSW_SPLIT] EQLU BKTSC_OVERHDSZ<0, 16>
: 1583      1643 6      | THEN
: 1584      1644 6      | IRAB[IRBSV_EMPTY_BKT] = 1;
: 1585      1645 6      |
: 1586      1646 6      | ! Only force the record into the low bucket if it is not the
: 1587      1647 6      | ! first one in the duplicate chain.
: 1588      1648 6      | !
: 1589      1649 6      |
: 1590      1650 6      | IF .END_CHAIN GEQU .POS_INSERT
: 1591      1651 6      | AND .IRAB[IRBSW_SPLIT] NEQU .IRAB[IRBSW_POS_INS]
: 1592      1652 6      | THEN
: 1593      1653 6      | IRAB[IRBSV_REC_W_LO] = 1;
: 1594      1654 5      | END;
: 1595      1655 5      |
: 1596      1656 5      | LEAVE DO_IT
: 1597      1657 5      |
: 1598      1658 4      | END; ! ( end of duplicates found and they fit in one bucket )
: 1599      1659 4      |
: 1600      1660 4      | ! if we had 255 dups above we dropped thru to here and this next test
: 1601      1661 4      | ! will fail because it can only happen on an update so the all dups case
: 1602      1662 4      | ! will fall thru to split type 2, which will put the new record by itself.
: 1603      1663 4      | ! consider oddball update case in which there are dups before and after
: 1604      1664 4      | ! position of insert. ( note that if this case doesn't apply, the duplicates
: 1605      1665 4      | ! were only before or after -- and didn't fit w/ record -- so new record
: 1606      1666 4      | ! will end up by itself. for code flow purposes, leave that till later).
: 1607      1667 4      |
: 1608      1668 4      |
: 1609      1669 4      | IF .IRAB[IRBSV_DUPS_SEEN]
: 1610      1670 4      | AND
```

```
1611      1671  4      .END_CHAIN GTRU .POS_INSERT
1612      1672  4      THEN
1613      1673  5      BEGIN
1614      1674  5
1615      1675  5      IF .DUPS - (.POS_INSERT - .BEG_CHAIN) LSSU .BKTSIZE
1616      1676  5      THEN
1617      1677  5          | if high dups will fit w/ record, put them in a bucket together
1618      1678  5          |
1619      1679  5          BEGIN
1620      1680  6
1621      1681  6          +
1622      1682  6          |!!!! SPLIT TYPE 4 !!!!!
1623      1683  6          | 3 bkt split where middle bkt is a continuation bkt containing
1624      1684  6          | new record and dups following it
1625      1685  6          |
1626      1686  6          |!!!! AND SPLIT TYPE 4B !!!!! however, if the hi set consists
1627      1687  6          | solely of duplicates, we can still have a 2-bkt split case that
1628      1688  6          | would not have been picked up by the previous algorithm ( since
1629      1689  6          | it won't divide dups).
1630      1690  6          |
1631      1691  6          |
1632      1692  6          IRAB[IRBSV_CONT_BKT] = 1;
1633      1693  6          IRAB[IRBSW_SPLIT] = .IRAB[IRBSW_POS_INS];
1634      1694  6
1635      1695  6          IF .END_CHAIN EQLU .EOB
1636      1696  6          THEN
1637      1697  6              IRAB[IRBSV_NEW_BKTS] = 1
1638      1698  6          ELSE
1639      1699  6              IRAB[IRBSW_SPLIT_1] = .END_CHAIN - .BKT_ADDR;
1640      1700  6
1641      1701  6          REC_ADDR = .BEG_CHAIN;
1642      1702  6          AP = 0;
1643      1703  6          RMSRECORD_KEY(KEYBUF_ADDR(2));
1644      1704  6          LEAVE DO_IT
1645      1705  6
1646      1706  6          END;
1647      1707  5
1648      1708  5      | try to fit new record w/ before-dups in middle bucket
1649      1709  5      |
1650      1710  5
1651      1711  5      IF .DUPS - (.END_CHAIN - .POS_INSERT) LSSU .BKTSIZE
1652      1712  5      THEN
1653      1713  5          BEGIN
1654      1714  6
1655      1715  6          +
1656      1716  6          |!!!! SPLIT TYPE 5 !!!!!
1657      1717  6          | 3 or 4 bkt split ( depending on status of
1658      1718  6          | high set) where left-middle bkt is new record w/ before-dups
1659      1719  6          | and right-middle bkt, if it is needed, is a continuation bkt
1660      1720  6          | w/ the after-dups. it is needed iff the dups aren't the whole hi
1661      1721  6          | set it still is a continuation bkt.
1662      1722  6          |
1663      1723  6          |***** NOTE FROM NOV-7-78
1664      1724  6          | This case doesn't take into account the fact that the
1665      1725  6          | whole bucket may be dups. In the case of all dups, we could
1666      1726  6          | end up generating an empty bucket when we don't have to (if
1667      1727  6          |
```

```
: 1668      1728 6      | no RRV's) or a relatively useless bucket (some RRV's). In any
: 1669      1729 6      | event we could end up generating an extra bucket when we
: 1670      1730 6      | don't have to
: 1671      1731 6      |
: 1672      1732 6      |
: 1673      1733 6      | IRAB[IRB$W_SPLIT] = .BEG_CHAIN - .BKT_ADDR;
: 1674      1734 6      | IRAB[IRB$W_SPLIT_1] = .IRAB[IRB$W_POS_INS];
: 1675      1735 6      |
: 1676      1736 6      | IF .IRAB[IRB$W_SPLIT] EQLU BKT$C_OVERHDSZ<0, 16>
: 1677      1737 6      | THEN
: 1678      1738 6      |     IRAB[IRB$V_EMPTY_BKT] = 1;
: 1679      1739 6      |
: 1680      1740 6      | IRAB[IRB$V_REC_W_LO] = 1;
: 1681      1741 6      |
: 1682      1742 6      | IF .END_CHAIN LSSU .EOB
: 1683      1743 6      | THEN
: 1684      1744 7      |     BEGIN
: 1685      1745 7      |         IRAB[IRB$V_NEW_BKTS] = 3;
: 1686      1746 7      |         IRAB[IRB$W_SPLIT_2] = .END_CHAIN - .BKT_ADDR;
: 1687      1747 7      |     END
: 1688      1748 6      | ELSE
: 1689      1749 6      |     IRAB[IRB$V_CONT_R] = 1;
: 1690      1750 6      |
: 1691      1751 6      | LEAVE DO_IT
: 1692      1752 6      |
: 1693      1753 5      | END;
: 1694      1754 5      |
: 1695      1755 5      | ! { end of oddball update case w/ dups on both sides of new record }
: 1696      1756 5      |
: 1697      1757 4      | END;
: 1698      1758 4      |
: 1699      1759 4      |
: 1700      1760 4      | +
: 1701      1761 4      | !!!! SPLIT TYPE 2 !!!!!
: 1702      1762 4      | the new record must go all by itself therefore,
: 1703      1763 4      | this is a 3-bkt split if there are no after-dups or no hi set and a 4-bkt
: 1704      1764 4      | split if both of those exist even more exceptional, this can still be a
: 1705      1765 4      | 2-bkt split if there is no hi set at all ---- i.e., eob = end of the dups
: 1706      1766 4      | chain
: 1707      1767 4      |
: 1708      1768 4      | IRAB[IRB$W_SPLIT] = IRAB[IRB$W_SPLIT_1] = .IRAB[IRB$W_POS_INS];
: 1709      1769 4      |
: 1710      1770 4      | IF .IRAB[IRB$V_DUPS_SEEN]
: 1711      1771 4      | THEN
: 1712      1772 5      |     BEGIN
: 1713      1773 5      |         IRAB[IRB$V_CONT_BKT] = 1;
: 1714      1774 5      |         REC_ADDR = .BEG_CHAIN;
: 1715      1775 5      |         AP = 0;
: 1716      1776 5      |         RMS$RECORD_KEY(KEYBUF_ADDR(2));
: 1717      1777 4      |     END;
: 1718      1778 4      |
: 1719      1779 4      | IF .POS_INSERT EQLU .EOB
: 1720      1780 4      | THEN
: 1721      1781 4      |     IRAB[IRB$V_NEW_BKTS] = 1
: 1722      1782 4      | ELSE
: 1723      1783 4      |
: 1724      1784 4      |     IF .POS_INSERT LSSU .END_CHAIN
```

```
: 1725      1785  4      THEN
: 1726      1786  5      BEGIN
: 1727      1787  5
: 1728      1788  5      IF .END_CHAIN LSSU .EOB
: 1729      1789  5      THEN
: 1730      1790  5          IRAB[IRBSV_NEW_BKTS] = 3
: 1731      1791  5      ELSE
: 1732      1792  5          IRAB[IRBSV_CONT_R] = 1;
: 1733      1793  5
: 1734      1794  5      IRAB[IRBSW_SPLIT_2] = .END_CHAIN - .BKT_ADDR;
: 1735      1795  4      END;
: 1736      1796  4
: 1737      1797  3      END;
: 1738      1798  3      ! { end of block defining local symbols }
: 1739      1799  2      END;
: 1740      1800  2      ! { end of do_it }
: 1741      1801  2      ! if the first split point is at the beginning of the data, this means that
: 1742      1802  2      ! all data records will be moved out and only rrv's will be left in the
: 1743      1803  2      ! original bucket ..... therefore, we can mark this bucket as empty
: 1744      1804  2
: 1745      1805  2
: 1746      1806  2      IF .IRAB[IRBSW_SPLIT] EQLU BKT$C_OVERHDSZ<0, 16>
: 1747      1807  2      AND
: 1748      1808  2      NOT .IRAB[IRBSV_REC_W_LO]
: 1749      1809  2      THEN
: 1750      1810  2          IRAB[IRBSV_EMPTY_BKT] = 1;
: 1751      1811  2
: 1752      1812  2      RETURN;
: 1753      1813  2
: 1754      1814  1      END;
:                               ! { end of routine }
```

				OC	BB	00000	RM\$SPLIT_UDR::		
							PUSHR	#^M<R2,R3>	0874
							SUBL2	#32, SP	
							BICB2	#72, 68(IRAB)	0986
44	A9	02	44	01	48	8F 8A 00005	INSV	#1, #1, #2, 68(IRAB)	0987
						7E D4 00010	CLRL	NUM_RRVs	0988
						56 DD 00012	PUSHL	REC_ADDR	0989
				52	0E	A5 9E 00014	MOVAB	14(R5), R2	0990
				56		52 D0 00018	MOVL	R2, REC_ADDR	
				50	04	A5 3C 0001B	MOVZWL	4(BKT_ADDR), R0	0991
						6045 9F 0001F	PUSHAB	(R0)[BKT_ADDR]	
						7E D4 00022	CLRL	LAST	0992
		20		66		03 E0 00024	BBS	#3, (REC_ADDR), 4\$	1000
				5C		03 D0 00028	MOVL	#3, AP	1004
						0000G 30 0002B	BSBW	RM\$RECORD_VBN	1006
				50	1C	A4 D1 0002E	CMPL	28(BDB), R0	
						05 12 00032	BNEQ	2\$	
					OC	AE B6 00034	INCW	NUM_RRVs	1008
						03 11 00037	BRB	3\$	
					OE	AE B6 00039	INCW	NUM_RRVs+2	1018
				6E		56 D0 0003C	MOVL	REC_ADDR, LAST	1020
						0000G 30 0003F	BSBW	RM\$GETNEXT_REC	1021

	04	AE		56	D1	00042	CMPL	REC_ADDR, EOB	1023	
				DC	1F	00046	BLSSU	1\$		
50		56		55	C3	00048	4\$:	SUBL3	BKT_ADDR, REC_ADDR, R0	1029
	4E	A9		50	B0	0004C		MOVW	R0, 78(IRAB)	
	4C	A9		50	B0	00050		MOVW	R0, 76(IRAB)	
51		51	17	A7	9A	00054		MOVZBL	23(IDX_DFN), R1	1030
		51		09	78	00058		ASHL	#9, R1 - R1	
	1C	AE	F1	A1	9E	0005C		MOVAB	-15(R1), BKTSIZE	
	2C	AE	22	A8	3C	00061		MOVZWL	34(RAB), REC_SIZE	1032
	2C	AE		07	C0	00066		ADDL2	#7, REC_SIZE	
		01	50	AA	91	0006A		CPMB	80(IFABT), #1	1034
				04	13	0006E		BEQL	5\$	
	2C	AE		02	C0	00070		ADDL2	#2, REC_SIZE	1036
0A		06		03	E1	00074	5\$:	BBC	#3, 6(IRAB), 6\$	1040
	78	A9	1C	A4	D1	00079		CMPL	28(BDB), 120(IRAB)	1044
				03	12	C007E		BNEQ	6\$	
			0C	AE	B6	00080		INCW	NUM_RRVS	1046
20	AE	04		56	C3	00083	6\$:	SUBL3	REC_ADDR, EOB, RRV	1050
		04		56	D0	00089		MOVL	REC_ADDR, EOB	1051
		52		56	D1	0008D		CMPL	REC_ADDR, R2	1056
				03	12	00090		BNEQ	7\$	
				02FD	31	00092		BRW	47\$	
		56	08	AE	D1	00095	7\$:	CMPL	POS_INSERT, REC_ADDR	1087
				3B	12	00099		BNEQ	9\$	
51		6E		01	C1	0009B		ADDL3	#1, LAST, R1	1089
		50		61	9A	0009F		MOVZBL	(R1), R0	
				50	D6	000A2		INCL	R0	
	06	A5		50	91	000A4		CPMB	R0, 6(BKT_ADDR)	
				2C	12	C00A8		BNEQ	9\$	
		56		6E	D0	000AA		MOVL	LAST, REC_ADDR	1092
				7E	D4	000AD		CLRL	-(SP)	1094
		7E	20	A7	9A	000AF		MOVZBL	32(IDX_DFN), -(SP)	
		50	00B4	CA	3C	000B3		MOVZWL	180(IFAB), R0	
			60	B940	3F	000B8		PUSHAW	296(IRAB)[R0]	
				0000G	30	000BC		BSBW	RMSCOMPARE_REC	
		5E		0C	C0	000BF		ADDL2	#12, SP	
		03		50	E8	000C1		BLBS	R0, 8\$	
				025B	31	000C5		BRW	43\$	
		50		56	D0	000C8	8\$:	MOVL	REC_ADDR, R0	1103
				FEA9	30	000CB		BSBW	RMSMOVE_KEY	
	4A	A9	48	A9	B0	000CE		MOVW	72(IRABT), 74(IRAB)	1104
				044B	31	000D3		BRW	72\$	1105
		56	0E	A5	9E	000D6	9\$:	MOVAB	14(R5), REC_ADDR	1116
	24	AE	7FFFFFFF	8F	D0	000DA		MOVL	#2147483647, LAST_DIFF	1117
				6E	D4	000E2		CLRL	LAST	1118
			28	AE	D4	000E4		CLRL	SAVE_REC_W_LO	1119
10	AE	04		56	C3	000E7	10\$:	SUBL3	REC_ADDR, EOB, RHS	1131
		08		56	D1	000ED		CMPL	REC_ADDR, POS_INSERT	1133
				0A	1A	000F1		BGTRU	11\$	
05		44		03	E0	000F3		BBS	#3, 68(IRAB), 11\$	1135
		10		AE	C0	000F8		ADDL2	REC_SIZE, RHS	1137
		1C		AE	D1	000FD	11\$:	CMPL	RHS, BKTSIZE	1146
				03	1F	00102		BLSSU	13\$	
				0139	31	00104	12\$:	BRW	32\$	
			06	A5	95	00107	13\$:	TSTB	6(BKT_ADDR)	1148
				13	12	0010A		BNEQ	14\$	
0E		06		03	E0	0010C		BBS	#3, 6(IRAB), 14\$	1150

		50	0E	A5	9E	00111	MOVAB	14(R5), R0	1152
		50		56	D1	00115	CMPL	REC_ADDR, R0	
				05	12	00118	BNEQ	14\$	
	E5	44	A9	03	E1	0011A	BBC	#3, 68(IRAB), 12\$	1154
	50		56	55	C3	0011F	SUBL3	BKT_ADDR, REC_ADDR, R0	1157
			52	A0	9E	00123	MOVAB	-14(R0), LHS	
		08	AE	56	D1	00127	CMPL	REC_ADDR, POS_INSERT	1159
				09	1F	0012B	BLSSU	15\$	
	04	44	A9	03	E1	0012D	BBC	#3, 68(IRAB), 15\$	1161
			52	AE	C0	00132	ADDL2	REC_SIZE, LHS	1163
	50		52	20	AE	C1	00136	RRV, LHS, R0	1172
			51	0C	AE	3C	0013B	NUM_RRV, R1	
14	AE		51	07	C5	0013F	MULL3	#7, R1, 20(SP)	
			50	14	AE	C0	00144	ADDL2	20(SP), R0
		1C	AE	50	D1	00148	CMPL	R0, BKT_SIZE	
				14	1A	0014C	BGTRU	6\$	
			06	A5	95	0014E	TSTB	6(BKT_ADDR)	1180
				38	12	00151	BNEQ	20\$	
	33	06	A9	03	E0	00153	BBS	#3, 6(IRAB), 20\$	1182
				0E	AE	B5	00158	TSTW	NUM_RRV+2
				2E	12	0015B	BNEQ	20\$	1184
	29	44	A9	03	E1	0015D	BBC	#3, 68(IRAB), 20\$	1186
				6E	D5	00162	TSTL	LAST	1190
				03	12	00164	BNEQ	17\$	
				01BA	31	00166	BRW	43\$	
		56		6E	D0	00169	MOVL	LAST, REC_ADDR	1194
		04	28	AE	E8	0016C	BLBS	SAVE_REC_W LO, 18\$	1196
		44	A9	08	8A	00170	BICB2	#8, 68(IRAB)	1198
				56	D0	00174	MOVL	REC_ADDR, TMP	1209
				0E	A5	9E	00177	MOVAB	14(R5), REC_ADDR
				56	D0	0017B	MOVL	REC_ADDR, LAST	1210
				56	D1	0017E	CMPL	REC_ADDR, TMP	1211
				6C	13	00181	BEQL	28\$	1213
		6E		56	D0	00183	MOVL	REC_ADDR, LAST	1216
				0000G	30	00186	BSBW	RMSGETNEXT_REC	1217
				F3	11	00189	BRB	19\$	1213
	50		1C	AE	C5	0018B	MULL3	BKT_SIZE, LHS, R0	1257
	51	1C	AE	14	AE	C3	00190	SUBL3	20(SP), BKT_SIZE, R1
	51	20	AE	51	C3	00196	SUBL3	R1, RRV, R1	1258
				51	AE	C4	0019B	MULL2	RHS, R1
18	AE		50	51	C1	0019F	ADDL3	R1, R0, DIFFERENCE	1257
				03	18	001A4	BGEQ	21\$	1260
				0086	31	001A6	BRW	31\$	
		51	18	AE	D0	001A9	MOVL	DIFFERENCE, R1	1268
				03	18	001AD	BGEQ	22\$	
		51		51	CE	001AF	MNEGL	R1, R1	
		50	24	AE	D0	001B2	MOVL	LAST_DIFF, R0	
				03	18	001B6	BGEQ	23\$	
		50		50	CE	001B8	MNEGL	R0, R0	
		50		51	D1	001BB	CMPL	R1, R0	
				14	1B	001BE	BLEQU	26\$	
		6E		56	D1	001C0	CMPL	REC_ADDR, LAST	1272
				09	13	001C3	BEQL	24\$	
		56		6E	D0	001C5	MOVL	LAST, REC_ADDR	1276
		08	AE	56	D1	001C8	CMPL	REC_ADDR, POS_INSERT	1278
				04	1E	001CC	BGEQU	25\$	
		44	A9	08	8A	001CE	BICB2	#8, 68(IRAB)	1280

			6E	D4	001D2	25\$:	CLRL	LAST	1282		
			6E	D5	001D4	26\$:	TSTL	LAST	1289		
			17	12	001D6		BNEQ	28\$			
		53	56	D0	001D8		MOVL	REC_ADDR, TMP	1296		
		56	0E	A5	9E	001DB	MOVAB	14(R5), REC_ADDR	1297		
		6E	56	D0	001DF		MOVL	REC_ADDR, LAST	1298		
		53	56	D1	001E2	27\$:	CMPL	REC_ADDR, TMP	1300		
			08	13	001E5		BEQL	28\$			
		6E	56	D0	001E7		MOVL	REC_ADDR, LAST	1303		
			0000G	30	001EA		BSBW	RMSGETNEXT_REC	1304		
			F3	11	001ED		BRB	27\$	1300		
		50	6E	D0	001EF	28\$:	MOVL	LAST, R0	1309		
			FD82	30	001F2		BSBW	RMSMOVE_KEY			
4A	A9		56	A3	001F5		SUBW3	BKT_ADDR, REC_ADDR, 74(IRAB)	1310		
		48	A9	4A	A9	B1	001FA	CMPW	74(IRAB), 72(IRAB)	1316	
			0C	12	001FF		BNEQ	29\$			
		4C	A9	4A	A9	B1	00201	CMPW	74(IRAB), 76(IRAB)	1319	
			05	12	00206		BNEQ	29\$			
	03		44	A9	03	E1	00208	BBC	#3, 68(IRAB), 30\$	1322	
					0311	31	0020D	BRW	72\$		
			5C		05	D0	00210	MOVL	#3, AP	1329	
			50	00B4	CA	3C	00213	MOVZWL	180(IFAB), R0	1332	
			53	60	B940	3E	00218	MOVAV	296(IRAB)[R0], R3		
	51		50	60	A9	C1	0021D	ADDL3	96(IRAB), R0, R1	1331	
			50	20	A7	9A	00222	MOVZBL	32(IDX DFN), R0		
					0000G	30	00226	BSBW	RMSCOMPARE_KEY		
			E1	50	E8	00229	BLBS	R0, 29\$			
					01EF	31	0022C	BRW	55\$	1335	
		24	AE	18	AE	D0	0022F	MOVL	DIFFERENCE, LAST_DIFF	1346	
			6E		56	D0	00234	MOVL	REC_ADDR, LAST	1347	
	04		44	A9	03	E1	00237	BBC	#3, 68(IRAB), 32\$	1349	
			28	AE	01	D0	0023C	MOVL	#1, SAVE_REC_W_LO	1351	
			08	AE	56	D1	00240	CMPL	REC_ADDR, POS_INSERT	1360	
					3F	12	00244	BNEQ	36\$		
	3A		44	A9	03	E0	00246	BBS	#3, 68(IRAB), 36\$	1362	
	0A		06	A9	03	E1	0024B	BBC	#3, 6(IRAB), 33\$	1369	
			78	A9	1C	A4	D1	00250	CMPL	28(BDB), 120(IRAB)	1373
					03	12	00255	BNEQ	33\$		
					0C	AE	B7	00257	DECW	NUM_RRVS	1375
			44	A9	08	88	0025A	BISB2	#8, 68(IRAB)	1379	
			50		56	D0	0025E	MOVL	REC_ADDR, R0	1380	
					FD13	30	00261	BSBW	RMSMOVE_KEY		
		04	AE		56	D1	00264	CMPL	REC_ADDR, EOB	1382	
					03	12	00268	BNEQ	35\$		
					FE7A	31	0026A	BRW	10\$		
					7E	D4	0026D	CLRL	-(SP)	1386	
			7E		A7	9A	0026F	MOVZBL	32(IDX DFN), -(SP)		
			50	00B4	CA	3C	00273	MOVZWL	180(IFAB), R0		
				60	B940	9F	00278	PUSHAB	296(IRAB)[R0]		
					0000G	30	0027C	BSBW	RMSCOMPARE_REC		
			5E		0C	C0	0027F	ADDL2	#12, SP		
			E5		50	E8	00282	BLBS	R0, 34\$		
			53	44	A9	90	00285	MOVB	68(IRAB), TMP	1401	
		44	A9		08	8A	00289	BICB2	#8, 68(IRAB)	1402	
			50		56	D0	0028D	MOVL	REC_ADDR, R0	1403	
					FCE4	30	00290	BSBW	RMSMOVE_KEY		
		44	A9		53	90	00293	MOVB	TMP, 68(IRAB)	1404	

04	AE	56	D1	00297	37\$:	CMPL	REC_ADDR, EOB	1413	
		35	13	00298		BEQL	40\$		
	5C	03	D0	0029D		MOVL	#3, AP	1417	
		0000G	30	002A0		BSBW	RM\$RECORD_VBN	1419	
	50	1C	A4	D1	002A3	CMPL	28(BDB), R0		
			05	12	002A7	BNEQ	38\$		
		0C	AE	B7	002A9	DECW	NUM_RRVS	1421	
			03	11	002AC	BRB	39\$		
		0E	AE	B7	002AE	38\$:	DECW	NUM_RRVS+2	1423
		0000G	30	002B1	39\$:	BSBW	RM\$GETNEXT_REC	1425	
04	AE	56	D1	002B4		CMPL	REC_ADDR, EOB	1427	
		18	13	002B8		BEQL	40\$		
		7E	D4	002BA		CLRL	-(SP)	1435	
	7E	20	A7	9A	002BC	MOVZBL	32(IDX_DFN), -(SP)		
	50	00B4	CA	3C	002C0	MOVZWL	180(IFAB), R0		
		60	B940	9F	002C5	PUSHAB	@96(IRAB)[R0]		
			0000G	30	002C9	BSBW	RM\$COMPARE_REC		
	5E		0C	C0	002CC	ADDL2	#12, SP		
	C5		50	E9	002CF	BLBC	R0, 37\$		
08	AE	56	D1	002D2	40\$:	CMPL	REC_ADDR, POS_INSERT	1442	
		2F	12	002D6		BNEQ	41\$		
	5C		03	D0	002D8	MOVL	#3, AP	1449	
	50	00B4	CA	3C	002DB	MOVZWL	180(IFAB), R0	1452	
	53	60	B940	3E	002E0	MOVAW	@96(IRAB)[R0], R3		
51	50	60	A9	C1	002E5	ADDL3	96(IRAB), R0, R1	1451	
	50	20	A7	9A	002EA	MOVZBL	32(IDX_DFN), R0		
			0000G	30	002EE	BSBW	RM\$COMPARE_KEY		
	13		50	E8	002F1	BLBS	R0, 41\$		
44	A9		08	88	002F4	BISB2	#8, 68(IRAB)	1456	
0A	06		03	E1	002F8	BBC	#3, 6(IRAB), 41\$	1458	
	78	A9	1C	A4	D1	002FD	CMPL	28(BDB), 120(IRAB)	1460
				03	12	00302	BNEQ	41\$	
		0C	AE	B7	00304	DECW	NUM_RRVS	1462	
08	AE	56	D1	00307	41\$:	CMPL	REC_ADDR, POS_INSERT	1467	
		13	1B	0030B		BLEQU	42\$		
44	A9		08	88	0030D	BISB2	#8, 68(IRAB)	1470	
0A	06		03	E1	00311	BBC	#3, 6(IRAB), 42\$	1472	
	78	A9	1C	A4	D1	00316	CMPL	28(BDB), 120(IRAB)	1474
				03	12	0031B	BNEQ	42\$	
		0C	AE	B7	0031D	DECW	NUM_RRVS	1476	
			FDC4	31	00320	BRW	10\$	1128	
	01		02	F0	00323	INSV	#2, #1, #2, 68(IRAB)	1507	
44	A9		08	8A	00329	BICB2	#8, 68(IRAB)	1508	
	50	00B4	CA	3C	0032D	MOVZWL	180(IFAB), R0	1519	
		60	B940	9F	00332	PUSHAB	@96(IRAB)[R0]		
		60	B940	3F	00336	PUSHAW	@96(IRAB)[R0]		
	7E	20	A7	9A	0033A	MOVZBL	32(IDX_DFN), -(SP)		
			0000G	30	0033E	BSBW	RM\$MOVE		
	5E		0C	C0	00341	ADDL2	#12, SP		
28	AE	0E	A5	9E	00344	MOVAB	14(R5), 40(SP)	1524	
	56	28	AE	D0	00349	MOVL	40(SP), REC_ADDR		
			7E	D4	0034D	44\$:	CLRL	-(SP)	1530
	7E	00A6	C9	9A	0034F	MOVZBL	166(IRAB), -(SP)		
	50	00B4	CA	3C	00354	MOVZWL	180(IFAB), R0		
		60	B940	3F	00359	PUSHAW	@96(IRAB)[R0]		
			0000G	30	0035D	BSBW	RM\$COMPARE_REC		
	5E		0C	C0	00360	ADDL2	#12, SP		

			53		50	D0	00363	MOVL	R0, STATUS			
			35		53	E9	00366	BLBC	STATUS, 50\$			
		08	AE		56	D1	00369	CMPL	REC_ADDR, POS_INSERT	1534		
					11	1E	0036D	BGEQU	45\$			
					5C	D4	0036F	CLRL	AP	1537		
			50		CA	3C	00371	MOVZWL	180(IFAB), R0	1538		
				00B4	B940	9F	00376	PUSHAB	@96(IRAB)[R0]			
				60	0000G	30	0037A	BSBW	RMS\$RECORD_KEY			
					04	C0	0037D	ADDL2	#4, SP			
		04	5E		56	D1	00380	45\$: CMPL	REC_ADDR, EOB	1541		
			AE		04	13	00384	BEQL	46\$			
					53	D5	00386	TSTL	STATUS	1543		
					0F	18	00388	BGEQ	49\$			
			50		A9	3C	0038A	46\$: MOVZWL	72(IRAB), R0	1551		
		4C	A9		50	B0	0038E	MOVW	R0, 76(IRAB)			
		4A	A9		50	B0	00392	47\$: MOVW	R0, 74(IRAB)			
					0188	31	00396	48\$: BRW	72\$	1552		
					0000G	30	00399	49\$: BSBW	RMS\$GETNEXT_REC	1558		
					AF	11	0039C	BRB	44\$	1530		
				0C	AE	D4	0039E	50\$: CLRL	NUM_RRVS	1565		
			52		56	D0	003A1	MOVL	REC_ADDR, LHS	1566		
				0C	AE	D6	003A4	51\$: INCL	NUM_RRVS	1570		
					0000G	30	003A7	BSBW	RMS\$GETNEXT_REC	1571		
		04	AE		56	D1	003AA	CMPL	REC_ADDR, EOB	1573		
					19	13	003AE	BEQL	52\$			
					7E	D4	003B0	CLRL	-(SP)	1578		
			7E		C9	9A	003B2	MOVZBL	166(IRAB), -(SP)			
			50		00B4	CA	003B7	MOVZWL	180(IFAB), R0			
					60	B940	3F	003BC	PUSHAW	@96(IRAB)[R0]		
					0000G	30	003C0	BSBW	RMS\$COMPARE_REC			
					0C	C0	003C3	ADDL2	#12, SP			
			5E		50	E9	003C6	BLBC	R0, 51\$			
		10	AE		56	D0	003C9	52\$: MOVL	REC_ADDR, RHS	1581		
			52		08	AE	003CD	CMPL	POS_INSERT, LHS	1588		
					05	1B	003D1	BLEQU	53\$			
			44	A9	80	8F	003D3	BISB2	#128, 68(IRAB)	1590		
	20	AE	10	AE	52	C3	003D8	53\$: SUBL3	LHS, RHS, RRV	1592		
			20	AE	2C	C0	003DE	ADDL2	REC_SIZE, RRV	1593		
			1C	AE	20	AE	003E3	CMPL	RRV, BKTSIZE	1595		
					59	1E	003E8	BGEQU	58\$			
		07	06	A9	03	E0	003EA	BBS	#3, 6(IRAB), 54\$	1601		
			FE	8F	0C	AE	003EF	CMPB	NUM_RRVS, #254	1603		
					4D	1A	003F4	BGTRU	58\$			
	4A	A9		52	55	A3	003F6	54\$: SUBW3	BKT_ADDR, LHS, 74(IRAB)	1622		
	4C	A9		10	55	A3	003FB	SUBW3	BKT_ADDR, RHS, 76(IRAB)	1623		
			04	AE	10	AE	00401	CMPL	RHS, EOB	1625		
					1C	12	00406	BNEQ	56\$			
44	A9		02	01	01	F0	00408	INSV	#1, #1, #2, 68(IRAB)	1628		
				28	AE	52	0040E	CMPL	LHS, 40(SP)	1630		
					82	12	00412	BNEQ	48\$			
	4C	A9		4E	A9	B0	00414	MOVW	78(IRAB), 76(IRAB)	1633		
	4A	A9		48	A9	B0	00419	MOVW	72(IRAB), 74(IRAB)	1634		
	44	A9			10	88	0041E	55\$: BISB2	#16, 68(IRAB)	1635		
					6F	11	00422	BRB	62\$	1630		
			0E	4A	A9	B1	00424	56\$: CMPW	74(IRAB), #14	1642		
					05	12	00428	BNEQ	57\$			
	44	A9		40	8F	88	0042A	BISB2	#64, 68(IRAB)	1644		

	08	AE	10	AE	D1	0042F	57\$:	CMPL	RHS, POS_INSERT	1650
				5D	1F	00434		BLSSU	62\$	
	48	A9	4A	A9	B1	00436		CMPL	74(IRAB), 72(IRAB)	1651
				56	13	00438		BEQL	62\$	
	44	A9		08	88	0043D		BISB2	#8, 68(IRAB)	1653
				50	11	00441		BRB	62\$	1656
			44	A9	95	00443	58\$:	TSTB	68(IRAB)	1669
				03	19	00446		BLSS	59\$	
				0080	31	00448		BRW	65\$	
	08	AE	10	AE	D1	0044B	59\$:	CMPL	RHS, POS_INSERT	1671
				79	1B	00450		BLEQU	65\$	
50		52	08	AE	C3	00452		SUBL3	POS_INSERT, LHS, R0	1675
		50	20	AE	C0	00457		ADDL2	RRV, R0	
	1C	AE		50	D1	0045B		CMPL	R0, BKTSIZE	
				34	1E	0045F		BGEQU	63\$	
	44	A9		10	88	00461		BISB2	#16, 68(IRAB)	1693
	4A	A9	48	A9	B0	00465		MOVW	72(IRAB), 74(IRAB)	1694
	04	AE	10	AE	D1	0046A		CMPL	RHS, EOB	1696
				08	12	0046F		BNEQ	60\$	
44	A9	02		01	F0	00471		INSV	#1, #1, #2, 68(IRAB)	1698
				06	11	00477		BRB	61\$	
	4C	A9	10	AE	A3	00479	60\$:	SUBW3	BKT_ADDR, RHS, 76(IRAB)	1700
				56	D0	0047F	61\$:	MOVL	LHS, REC_ADDR	1702
				5C	D4	00482		CLRL	AP	1703
			50	CA	3C	00484		MOVZWL	180(IFAB), R0	1704
			60	B940	9F	00489		PUSHAB	296(IRAB)[R0]	
				0000G	30	0048D		BSBW	RMSRECORD_KEY	
				04	C0	00490		ADDL2	#4, SP	
				6C	11	00493	62\$:	BRB	67\$	1705
	50	08	10	AE	C3	00495	63\$:	SUBL3	RHS, POS_INSERT, R0	1712
			20	AE	C0	0049B		ADDL2	RRV, R0	
	1C	AE		50	D1	0049F		CMPL	R0, BKTSIZE	
				26	1E	004A3		BGEQU	65\$	
				55	A3	004A5		SUBW3	BKT_ADDR, LHS, 74(IRAB)	1733
	4C	A9	48	A9	B0	004AA		MOVW	72(IRAB), 76(IRAB)	1734
			4A	A9	B1	004AF		CMPL	74(IRAB), #14	1736
				05	12	004B3		BNEQ	64\$	
	44	A9	40	8F	88	004B5		BISB2	#64, 68(IRAB)	1738
	44	A9		08	88	004BA	64\$:	BISB2	#8, 68(IRAB)	1740
	04	AE	10	AE	D1	004BE		CMPL	RHS, EOB	1742
				4C	1F	004C3		BLSSU	69\$	
	44	A9		20	88	004C5		BISB2	#32, 68(IRAB)	1749
				56	11	004C9		BRB	72\$	1751
			48	A9	3C	004CB	65\$:	MOVZWL	72(IRAB), R0	1768
	4C	A9		50	B0	004CF		MOVW	R0, 76(IRAB)	
	4A	A9		50	B0	004D3		MOVW	R0, 74(IRAB)	
			44	A9	95	004D7		TSTB	68(IRAB)	1770
				18	18	004DA		BGEQ	66\$	
	44	A9		10	88	004DC		BISB2	#16, 68(IRAB)	1773
				52	D0	004E0		MOVL	LHS, REC_ADDR	1774
				5C	D4	004E3		CLRL	AP	1775
			50	CA	3C	004E5		MOVZWL	180(IFAB), R0	1776
			60	B940	9F	004EA		PUSHAB	296(IRAB)[R0]	
				0000G	30	004EE		BSBW	RMSRECORD_KEY	
				04	C0	004F1		ADDL2	#4, SP	
	04	AE	08	AE	D1	004F4	66\$:	CMPL	POS_INSERT, EOB	1779
				08	12	004F9		BNEQ	68\$	

RM3SPLUDR
V04-000

RM\$SPLIT_UDR

K 2
16-Sep-1984 02:03:28 VAX-11 Bliss-32 V4.0-742
14-Sep-1984 13:01:40 [RMS.SRC]RM3SPLUDR.B32;1

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44	A9	02	01	01	F0	004FB	INSV	#1, #1, #2, 68(IRAB)
				1E	11	00501	BRB	72\$
		10	AE	08	AE	D1 00503	CMPL	POS_INSERT, RHS
				17	1E	00508	BGEQU	72\$
		04	AE	10	AE	D1 0050A	CMPL	RHS, EOB
				06	1E	0050F	BGEQU	70\$
		44	A9		06	88 00511	BISB2	#6, 68(IRAB)
				04	11	00515	BRB	71\$
		44	A9		20	88 00517	BISB2	#32, 68(IRAB)
4E	A9	10	AE		55	A3 0051B	SUBW3	BKT_ADDR, RHS, 78(IRAB)
			OE	4A	A9	B1 00521	CMPL	74(IRAB), #14
					0A	12 00525	BNEQ	73\$
	05	44	A9		03	E0 00527	BBS	#3, 68(IRAB), 73\$
		44	A9	40	8F	88 0052C	BISB2	#64, 68(IRAB)
			5E		30	C0 00531	ADDL2	#48, SP
					0C	8A 00534	POPR	#^M<R2,R3>
					05	00536	RSB	

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:

; Routine Size: 1335 bytes, Routine Base: RM\$RMS3 + 0089

; 1755 1815 1

```
1757 1816 1 %SBTTL 'RMSSPLIT_UDR_3'
1758 1817 1 GLOBAL ROUTINE RMSSPLIT_UDR_3(RECSZ) : RLSRABREG_4567 NOVALUE =
1759 1818 1
1760 1819 1 ++
1761 1820 1
1762 1821 1 FUNCTIONAL DESCRIPTION:
1763 1822 1 This routine calculates bucket splits for prologue 3 version files.
1764 1823 1
1765 1824 1 CALLING SEQUENCE:
1766 1825 1 BSBW RMSSPLIT_UDR_3(RECSZ)
1767 1826 1
1768 1827 1 INPUT PARAMETERS:
1769 1828 1 RECSZ - packed record size including overhead
1770 1829 1
1771 1830 1 IMPLICIT INPUTS:
1772 1831 1 BDB pointer
1773 1832 1 BUFFER pointer
1774 1833 1 REC_ADDR -- point of insert
1775 1834 1 RAB -- to be passed to RMSMOVE_KEY
1776 1835 1 IDX_DFN
1777 1836 1 in IRAB -- CURBDB, associated with bdb and bkt_addr
1778 1837 1 POS_INS corresponding to REC_ADDR
1779 1838 1 key_buffer address
1780 1839 1 in IFAB -- key buffer size
1781 1840 1 BKT$B_NXTRECID = 0 in original bucket signals that this is
1782 1841 1 a split due to a lack of id's in the bucket
1783 1842 1
1784 1843 1 OUTPUT PARAMETERS:
1785 1844 1 none
1786 1845 1
1787 1846 1 IMPLICIT OUTPUTS:
1788 1847 1 in IRAB --
1789 1848 1 if 2 bkt split --
1790 1849 1 IRBSW_SPLIT, offset to split point
1791 1850 1 IRBSV_REC_W_LO -- set if split point is pos_insert and
1792 1851 1 record goes with lo set
1793 1852 1 key buffer 2 - new high key for original bucket, i.e. key to be
1794 1853 1 inserted at the index level
1795 1854 1 key buffer 4 - old high key
1796 1855 1 number of new buckets = 1
1797 1856 1 if original bucket was all rrv's, set IRBSV_EMPTY_BKT flag
1798 1857 1 if new bucket is a continuation bkt., set IRBSV_CONT_BKT flag
1799 1858 1 if 3 bkt split --
1800 1859 1 same as above with these changes:
1801 1860 1 key buffer 3 - implicitly it contains second key to be inserted
1802 1861 1 at the index level
1803 1862 1 IRBSW_SPLIT_1, offset to second split point
1804 1863 1 number of new buckets = 2
1805 1864 1 if right bucket is a continuation bkt, set IRBSV_CONT_R flag
1806 1865 1 if 4 bkt split --
1807 1866 1 same as above with these changes:
1808 1867 1 IRBSW_SPLIT_2, offset to third split point
1809 1868 1 number of new buckets = 3
1810 1869 1
1811 1870 1 ROUTINE VALUE:
1812 1871 1 rmssuc
1813 1872 1
```

```
1814 1873 1 | SIDE EFFECTS:
1815 1874 1 |   AP is clobbered
1816 1875 1 |
1817 1876 1 | --
1818 1877 1 |
1819 1878 2 | BEGIN
1820 1879 2 |
1821 1880 2 |   EXTERNAL REGISTER
1822 1881 2 |     COMMON_RAB_STR,
1823 1882 2 |     R_REC_ADDR_STR,
1824 1883 2 |     R_IDX_DFN_STR,
1825 1884 2 |     COMMON_IO_STR;
1826 1885 2 |
1827 1886 2 |   LOCAL
1828 1887 2 |     SAVE_REC_W_LO,
1829 1888 2 |     NEED_RRV,
1830 1889 2 |     POS_INSERT,
1831 1890 2 |     EOB,
1832 1891 2 |     RRV,
1833 1892 2 |     RMS,
1834 1893 2 |     LMS,
1835 1894 2 |     LAST : REF BBLOCK,
1836 1895 2 |     LAST_DIFF,
1837 1896 2 |     BKTSIZE,
1838 1897 2 |     DIFFERENCE;
1839 1898 2 |
1840 1899 2 |   LITERAL
1841 1900 2 |     RRV_SIZE = 9;
1842 1901 2 |
1843 1902 2 |   LABEL
1844 1903 2 |     DO_IT,
1845 1904 2 |     HALF,
1846 1905 2 |     NEXT;
1847 1906 2 |
1848 1907 2 | DO_IT :
1849 1908 2 |
1850 1909 3 |   BEGIN
1851 1910 3 |
1852 1911 3 |     ! define a block so that we can have some common checks before returning
1853 1912 3 |     ! successfully
1854 1913 3 |     !
1855 1914 3 |   HALF :
1856 1915 3 |
1857 1916 4 |   BEGIN
1858 1917 4 |
1859 1918 4 |     !+
1860 1919 4 |     ! Define a block so that we can simulate a go-to (naughty, naughty),
1861 1920 4 |     ! if we have decided that we are positioning at the end of the bucket
1862 1921 4 |     ! & we're in somewhat of an ascending order, where the last record
1863 1922 4 |     ! inserted is a duplicate of the new record, skip over the 50-50 code
1864 1923 4 |     ! and go to the code to take duplicates into account.
1865 1924 4 |
1866 1925 4 |     ! scan 1 -- Calculate size of existing rrv's and total number of rrv's
1867 1926 4 |     ! needed to move the whole bucket out ( worst case). As a side effect,
1868 1927 4 |     ! adjust eob pointer to point to the rrv's instead of freespace. Assume
1869 1928 4 |     ! not empty bucket until showed otherwise.
1870 1929 4 |     !-
```

```
1871 1930 4
1872 1931 4 IRAB[IRBSV_EMPTY_BKT] = 0;
1873 1932 4
1874 1933 4 ! new rec is tried 1st with hi set, then with lo set
1875 1934 4
1876 1935 4 IRAB[IRBSV_REC_W_LO] = 0;
1877 1936 4 IRAB[IRBSV_NEW_BKTS] = 1; ! assume 2-bkt split until showed otherwise
1878 1937 4 NEED_RRV = 0;
1879 1938 4 POS_INSERT = .REC_ADDR;
1880 1939 4 REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
1881 1940 4 EOB = .BKT_ADDR + .BKT_ADDR[BKT$W_FREESPACE];
1882 1941 4 LAST = 0;
1883 1942 4
1884 1943 4 DO
1885 1944 5 BEGIN
1886 1945 5
1887 1946 5 BUILTIN
1888 1947 5 AP;
1889 1948 5
1890 1949 5 IF .REC_ADDR[IR$V_RRV]
1891 1950 5 THEN
1892 1951 5 EXITLOOP;
1893 1952 5
1894 1953 5 AP = 3;
1895 1954 5
1896 1955 5 IF .BDB[BDB$L_VBN] EQLU RMS$RECORD_VBN()
1897 1956 5 THEN
1898 1957 5 NEED_RRV = .NEED_RRV + 1;
1899 1958 5
1900 1959 5 LAST = .REC_ADDR;
1901 1960 5
1902 1961 5 ! If the front compression of the current record is zero, save its
1903 1962 5 ! address as the last noncompressed key. This may prevent a bucket
1904 1963 5 ! scan when it comes time to extract and re-expand the key of the
1905 1964 5 ! last record in the bucket.
1906 1965 5
1907 1966 5 IF .IDX_DFN[IDX$V_KEY_COMPR]
1908 1967 5 THEN
1909 1968 6 BEGIN
1910 1969 6
1911 1970 6 IF .(.REC_ADDR + RMS$REC_OVHD() + 1) < 0.8 > EQLU 0
1912 1971 6 THEN
1913 1972 6 IRAB[IRBS$L_LST_NCMP] = .REC_ADDR;
1914 1973 5
1915 1974 5 END;
1916 1975 5
1917 1976 5 RMS$GETNEXT_REC()
1918 1977 4 UNTIL .REC_ADDR GEQU .EOB; ! end of first scan
1919 1978 4
1920 1979 4 ! Now that we have the address of the last record in the bucket, store
1921 1980 4 ! the key of that record in key buffer 4, to be used by index updating.
1922 1981 4
1923 1982 4 IF .LAST NEQU 0
1924 1983 4 THEN
1925 1984 5 BEGIN
1926 1985 5
1927 1986 5 LOCAL
```

```
: 1928      1987  5      TMP_ADDR;
: 1929      1988  5
: 1930      1989  5      BUILTIN
: 1931      1990  5      AP;
: 1932      1991  5
: 1933      1992  5      TMP_ADDR = .REC_ADDR;
: 1934      1993  5      REC_ADDR = .LAST;
: 1935      1994  5      AP = 0;          ! overhead and compressed form
: 1936      1995  5      RMSRECORD_KEY(KEYBUF_ADDR(4));
: 1937      1996  5      REC_ADDR = .TMP_ADDR;
: 1938      1997  4      END;
: 1939      1998  4
: 1940      1999  4      ! Set SPLIT_2 and SPLIT_1 to be EOB, so if there are less than 3 new
: 1941      2000  4      ! buckets BKT_SPL can use the value without having to recalculate it.
: 1942      2001  4
: 1943      2002  4      IRAB[IRBSW_SPLIT_1] = IRAB[IRBSW_SPLIT_2] = .REC_ADDR - .BKT_ADDR;
: 1944      2003  4
: 1945      2004  4      ! Set up the bucket size
: 1946      2005  4
: 1947      2006  4      BKTSIZE = .IDX_DFN[IDX$B_DATBKTSZ]*512 - BKT$C_OVERHDSZ - BKT$C_DATBKTOVH;
: 1948      2007  4
: 1949      2008  4      ! If this is an update, may have to count in an rrv for the existing record
: 1950      2009  4
: 1951      2010  4
: 1952      2011  4      IF .IRAB[IRBSV_UPDATE]
: 1953      2012  4      THEN
: 1954      2013  5          BEGIN
: 1955      2014  5
: 1956      2015  5          IF .BDB[BDB$L_VBN] EQLU .IRAB[IRBSL_PUTUP_VBN]
: 1957      2016  5          THEN
: 1958      2017  5              NEED_RRV = .NEED_RRV + 1;
: 1959      2018  4          END;
: 1960      2019  4
: 1961      2020  4      RRV = .EOB - .REC_ADDR;          ! size of existing rrv's
: 1962      2021  4      EOB = .REC_ADDR;          ! adjust eob
: 1963      2022  4
: 1964      2023  4      ! special case it, if the bucket was all rrv's
: 1965      2024  4
: 1966      2025  4
: 1967      2026  4      IF .REC_ADDR EQLU .BKT_ADDR + BKT$C_OVERHDSZ
: 1968      2027  4      THEN
: 1969      2028  5          BEGIN
: 1970      2029  5
: 1971      2030  5          ! Bkt is all rrv's yet the record wouldn't fit so we need to
: 1972      2031  5          ! allocate another bkt ( 2 bkt split). Yet special case it so as not
: 1973      2032  5          ! to make another idx entry, only to update the existing one by
: 1974      2033  5          ! setting empty bucket flag.
: 1975      2034  5
: 1976      2035  5          IRAB[IRBSW_SPLIT] = .REC_ADDR - .BKT_ADDR;
: 1977      2036  5          LEAVE DO_IT
: 1978      2037  5
: 1979      2038  4          END;          ! end { of special case an all-rrv bucket }
: 1980      2039  4
: 1981      2040  4
: 1982      2041  4      * BLOCK 1 *
: 1983      2042  4      ! Special Case -- If we can detect a possible ascending order to these
: 1984      2043  4      ! records it probably will be better to do a straight point of insert split
```



```
: 1985 2044 4 | which would put the new record in a bucket all by itself.
: 1986 2045 4 | Do this kind of split if and only if all the following conditions are met:
: 1987 2046 4 |   1) the record is being inserted at the end of bucket
: 1988 2047 4 |   2) the last record physically in the bkt is the last record to have
: 1989 2048 4 |       been inserted
: 1990 2049 4 |   3) the last record and the new record do not have duplicate key values
: 1991 2050 4 |
: 1992 2051 4 | Note that if they are duplicates, we can still make an optimization by
: 1993 2052 4 | skipping the 50-50 split code.
: 1994 2053 4 |
: 1995 2054 4 | Note that LAST cannot be zero, since if it were we would have an all
: 1996 2055 4 | rrv bkt.
: 1997 2056 4 |
: 1998 2057 4 |
: 1999 2058 4 | IF .POS_INSERT EQLU .REC_ADDR
: 2000 2059 4 |     AND
: 2001 2060 5 |     (((.LAST[IRCSW_ID] + 1) AND %X'FFFF') EQLU .BKT_ADDR[BKTSW_NXTRECID])
: 2002 2061 4 | THEN
: 2003 2062 5 |     BEGIN
: 2004 2063 5 |
: 2005 2064 5 |     REC_ADDR = .LAST;
: 2006 2065 5 |
: 2007 2066 5 |     | Check for duplicates:
: 2008 2067 5 |     | If the key is compressed, and the new key has a length of zero, then
: 2009 2068 5 |     | we know it is a duplicate of the previous one.
: 2010 2069 5 |     | If the key is not compressed, then compare the new key (key buffer 3)
: 2011 2070 5 |     | with the previous key.
: 2012 2071 5 |     |
: 2013 2072 5 |
: 2014 2073 5 |     IF .IDX_DFN[IDX$V_KEY_COMPR]
: 2015 2074 5 |     THEN
: 2016 2075 6 |         BEGIN
: 2017 2076 6 |
: 2018 2077 6 |         IF .(.IRAB[IRBSL_RECBUF])<0,8> NEQU 0
: 2019 2078 6 |         THEN
: 2020 2079 7 |             BEGIN
: 2021 2080 7 |
: 2022 2081 7 |             | Since we have detected a possible ascending order in the
: 2023 2082 7 |             | input, let's try to optimize a little and split at the point
: 2024 2083 7 |             | of insert. Send the record by itself into the new bucket
: 2025 2084 7 |             | and store the new high key of the old bucket in keybuf2,
: 2026 2085 7 |             | the high key of the new bucket in keybuf4, and split point.
: 2027 2086 7 |             |
: 2028 2087 7 |             RMSMOVE(.IDX_DFN[IDX$B_KEYSZ],
: 2029 2088 7 |                     KEYBUF_ADDR(4),
: 2030 2089 7 |                     KEYBUF_ADDR(2));
: 2031 2090 7 |             RMSMOVE(.IDX_DFN[IDX$B_KEYSZ],
: 2032 2091 7 |                     KEYBUF_ADDR(3),
: 2033 2092 7 |                     KEYBUF_ADDR(4));
: 2034 2093 7 |             IRAB[IRBSW_SPLIT] = .IRAB[IRBSW_POS_INS];
: 2035 2094 7 |             LEAVE DO_IT;
: 2036 2095 7 |             END
: 2037 2096 6 |         ELSE
: 2038 2097 6 |             LEAVE HALF
: 2039 2098 6 |         END
: 2040 2099 5 |     ELSE
: 2041 2100 6 |         BEGIN
```

```
2042 2101 6
2043 2102 6
2044 2103 6
2045 2104 6
2046 2105 6
2047 2106 6
2048 2107 6
2049 2108 6
2050 2109 6
2051 2110 6
2052 2111 6
2053 2112 6
2054 2113 6
2055 2114 6
2056 2115 7
2057 2116 7
2058 2117 7
2059 2118 7
2060 2119 7
2061 2120 7
2062 2121 7
2063 2122 7
2064 2123 7
2065 2124 7
2066 2125 6
2067 2126 6
2068 2127 6
2069 2128 6
2070 2129 6
2071 2130 6
2072 2131 4
2073 2132 4
2074 2133 4
2075 2134 4
2076 2135 4
2077 2136 4
2078 2137 4
2079 2138 4
2080 2139 4
2081 2140 4
2082 2141 4
2083 2142 4
2084 2143 4
2085 2144 4
2086 2145 4
2087 2146 4
2088 2147 4
2089 2148 4
2090 2149 5
2091 2150 5
2092 2151 5
2093 2152 5
2094 2153 5
2095 2154 5
2096 2155 5
2097 2156 5
2098 2157 5

LOCAL
REC_OVHD;

BUILTIN
AP;

REC_OVHD = RMSREC_OVHD(0);
AP = 3; ! Contiguous compare of keys

IF RMSCOMPARE_KEY ( .REC_ADDR + .REC_OVHD,
                    KEYBUF_ADDR(3),
                    .IDX_DFN[IDX$B_KEYSZ] )
THEN
BEGIN
RMSMOVE(.IDX_DFN[IDX$B_KEYSZ],
        KEYBUF_ADDR(4),
        KEYBUF_ADDR(2));
RMSMOVE(.IDX_DFN[IDX$B_KEYSZ],
        KEYBUF_ADDR(3),
        KEYBUF_ADDR(4));
IRAB[IRB$W_SPLIT] = .IRAB[IRB$W_POS_INS];
LEAVE DO_IT;
END
ELSE
LEAVE HALF
END

! * end of BLOCK 1 *
END;

REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
IRAB[IRB$L_LST_NCMP] = .REC_ADDR;
LAST_DIFF = %X7FFFFFFF;
LAST = 0;
SAVE_REC_W_LO = 0;

! * BLOCK 2 *
! Start from the beginning of the bucket and scan rightward. First find the
! 1st place the rhs will fit in 1 bkt then, as long as the lhs will fit in
! a bkt, try to find an optimal point. If there is no point where the rhs
! and lhs will both fit, we can't do a 2-bkt split and this case will fall
! out.

WHILE 1
DO
BEGIN
RHS = .EOB - .REC_ADDR;

IF .REC_ADDR LEQU .POS_INSERT
AND
NOT .IRAB[IRB$V_REC_W_LO]
THEN
RHS = .RHS + .RECSZ;
```

```
2099 2158 5 | If the primary key is compressed, then the righthand side total must
2100 2159 5 | include the count of characters currently front compressed off the
2101 2160 5 | key of the record which will be first in the right bucket.
2102 2161 5
2103 2162 5 IF .IDX_DFN[IDX$V_KEY_COMP]
2104 2163 5 THEN
2105 2164 5 | If the point of insertion of the new (updated) record is the same
2106 2165 5 | as that of the current split point, and the new (updated) record
2107 2166 5 | is to go in the new (right) bucket, the the number of front
2108 2167 5 | compressed characters to be added to the righthand total comes
2109 2168 5 | from the currently compressed key of the new (updated) record.
2110 2169 5 | This key will be found in keybuffer 5, if the current operation
2111 2170 5 | is an $UPDATE, or in a record buffer, if the current operation is
2112 2171 5 | a $PUT.
2113 2172 5
2114 2173 6 IF (.REC_ADDR EQLA .POS_INSERT)
2115 2174 5 AND
2116 2175 5 NOT .IRAB[IRB$V_REC_W_LO]
2117 2176 5 THEN
2118 2177 5 IF .IRAB[IRB$V_UPDATE]
2119 2178 5 THEN
2120 2179 5 RHS = .RHS + .(KEYBUF_ADDR(5) + 1)<0,8>
2121 2180 5 ELSE
2122 2181 5 RHS = .RHS + .(.IRAB[IRB$L_RECBUF] + 1)<0,8>
2123 2182 5
2124 2183 5 | If the current split point is not at the point of insertion of
2125 2184 5 | the new (updated) record, or if it is but the new (updated)
2126 2185 5 | record is to go in the old (left) bucket, then the first record
2127 2186 5 | in the new (right) bucket will be the current record, and the
2128 2187 5 | number of characters currently front compressed off its key is
2129 2188 5 | added to the righthand side total.
2130 2189 5
2131 2190 5 ELSE
2132 2191 5 IF .REC_ADDR LSSA .EOB
2133 2192 5 THEN
2134 2193 5 RHS = .RHS + .(.REC_ADDR + RM$REC_OVHD(0) + 1)<0,8>;
2135 2194 5
2136 2195 5 * BLOCK 3 *
2137 2196 5 | The right hand side fits if there is enough room and there are id's
2138 2197 5 | available. Id's are always available in the new bucket in the update
2139 2198 5 | situation, or if we're leaving at least 1 record behind in the old
2140 2199 5 | bucket. note that nxtrecid is always zeroed if this is a split due to
2141 2200 5 | lack of id's.
2142 2201 5
2143 2202 5
2144 2203 5 IF .RHS LSSU .BKTSIZE
2145 2204 5 AND
2146 2205 6 (.BKT_ADDR[BKT$W_NXTRECID] NEQ 0
2147 2206 6 OR
2148 2207 6 .IRAB[IRB$V_UPDATE]
2149 2208 6 OR
2150 2209 7 .REC_ADDR NEQA (.BKT_ADDR + BKT$C_OVERHDSZ)
2151 2210 6 OR
2152 2211 6 .IRAB[IRB$V_REC_W_LO])
2153 2212 5 THEN
2154 2213 6 BEGIN
2155 2214 6 LHS = .REC_ADDR - (.BKT_ADDR + BKT$C_OVERHDSZ);
```

```
2156 2215 6
2157 2216 6
2158 2217 6
2159 2218 6
2160 2219 6
2161 2220 6
2162 2221 6
2163 2222 6
2164 2223 6
2165 2224 6
2166 2225 6
2167 2226 6
2168 2227 6
2169 2228 6
2170 2229 6
2171 2230 6
2172 2231 6
2173 2232 6
2174 2233 6
2175 2234 6
2176 2235 6
2177 2236 6
2178 2237 7
2179 2238 7
2180 2239 7
2181 2240 7
2182 2241 7
2183 2242 6
2184 2243 7
2185 2244 7
2186 2245 7
2187 2246 7
2188 2247 7
2189 2248 7
2190 2249 7
2191 2250 7
2192 2251 7
2193 2252 7
2194 2253 7
2195 2254 7
2196 2255 7
2197 2256 7
2198 2257 7
2199 2258 7
2200 2259 8
2201 2260 8
2202 2261 8
2203 2262 8
2204 2263 8
2205 2264 8
2206 2265 8
2207 2266 8
2208 2267 8
2209 2268 8
2210 2269 8
2211 2270 9
2212 2271 9

IF .REC_ADDR GEQU .POS_INSERT
AND
.IRAB[IRBSV_REC_W_LO]
THEN
LHS = .LHS + .RECSZ;

! * BLOCK 4 *
! will lhs fit ? lhs doesn't fit if there is no space in the
! bucket, or if there won't be any id's available in the bucket.
! if not & if there is no previous point at which it fit, goto 3-bkt
! split code if there is a previous place where we could have had a
! 2-bkt split, use it

IF .LHS + .RRV + (RRV_SIZE * .NEED_RRV) GTRU .BKTSIZE
! Id's will be available in the original bucket if we aren't
! out of id's to begin with, if this is an update,
! or if the new record is going in the new bucket
OR
(.BKT_ADDR[BKTSW_NXTRECID] EQL 0
AND
NOT .IRAB[IRBSV_UPDATE]
AND
.IRAB[IRBSV_REC_W_LO])
THEN
BEGIN
IF .LAST EQL 0
THEN
EXITLOOP;

REC_ADDR = .LAST;

IF NOT .SAVE_REC_W_LO
THEN
IRAB[IRBSV_REC_W_LO] = 0;

! 2 bkt split is possible rec_addr points to the most
! optimal place since we had to back up, reset last to point
! to the record immediately before the split point
BEGIN
LOCAL
TMP;

TMP = .REC_ADDR;
REC_ADDR = .BKT_ADDR + BKTSC_OVERHDSZ;
LAST = .REC_ADDR;

WHILE .REC_ADDR NEQU .TMP
DO
BEGIN
LAST = .REC_ADDR;
```

```
2213 2272 9
2214 2273 9
2215 2274 9
2216 2275 9
2217 2276 9
2218 2277 9
2219 2278 9
2220 2279 9
2221 2280 9
2222 2281 10
2223 2282 10
2224 2283 10
2225 2284 10
2226 2285 10
2227 2286 9
2228 2287 9
2229 2288 9
2230 2289 8
2231 2290 8
2232 2291 7
2233 2292 7
2234 2293 7
2235 2294 7
2236 2295 7
2237 2296 7
2238 2297 7
2239 2298 7
2240 2299 7
2241 2300 7
2242 2301 7
2243 2302 7
2244 2303 7
2245 2304 7
2246 2305 7
2247 2306 7
2248 2307 8
2249 2308 8
2250 2309 8
2251 2310 8
2252 2311 8
2253 2312 8
2254 2313 8
2255 2314 8
2256 2315 8
2257 2316 8
2258 2317 8
2259 2318 8
2260 2319 8
2261 2320 8
2262 2321 8
2263 2322 8
2264 2323 8
2265 2324 7
2266 2325 7
2267 2326 7
2268 2327 7
2269 2328 6

      ! If the front compression of the current record is zero,
      ! save its address as the last noncompressed key. This may
      ! prevent a bucket scan when it comes time to extract and
      ! re-expand the key of the last record in the bucket
      ! immediately before the split point.
      IF .IDX_DFN[IDX$V_KEY_COMPR]
      THEN
        BEGIN
          IF .(.REC_ADDR + RMSREC_OVHD() + 1) < 0.8> EQLU 0
          THEN
            IRAB[IRB$L_LST_NCMP] = .REC_ADDR;
          END;
        RMSGETNEXT_REC();
      END;
      END;
      RMSMOVE_KEY(.LAST, .REC_ADDR);
      IRAB[IRB$W_SPLIT] = .REC_ADDR - .BKT_ADDR;

      ! treat another exception case of the new record going off into
      ! a cont. bkt all by itself
      IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_POS_INS]
      THEN
        IF .IRAB[IRB$W_SPLIT] EQLU .IRAB[IRB$W_SPLIT_1]
        THEN
          IF NOT .IRAB[IRB$V_REC_W_LO]
          THEN
            BEGIN
              BUILTIN
                AP;
              AP = 3;

              ! If the new last key in the bucket equals the key
              ! to be inserted in the new bucket, then we have a
              ! continuation bucket.
              IF NOT RMSCOMPARE_KEY ( KEYBUF_ADDR(2),
                                      KEYBUF_ADDR(3),
                                      .IDX_DFN[IDX$B_KEYSZ] )
              THEN
                IRAB[IRB$V_CONT_BKT] = 1;
              END;
            LEAVE DO_IT
          END;
          ! end of * BLOCK 4 * (LHS does not fit)
```

```
: 2270      2329  6
: 2271      2330  6      ! lhs fits also, calculate the magic ratio
: 2272      2331  6
: 2273      2332  6      DIFFERENCE = (.LHS * .BKTSIZE) -
: 2274      2333  6      (.RHS * (.BKTSIZE - (RRV_SIZE * .NEED_RRV) - .RRV));
: 2275      2334  6
: 2276      2335  6      ! * BLOCK 5 *
: 2277      2336  6
: 2278      2337  6
: 2279      2338  6      IF .DIFFERENCE GEQ 0
: 2280      2339  6      THEN
: 2281      2340  7          BEGIN
: 2282      2341  7
: 2283      2342  7          ! found the 1st point at which the magic ratio is positive
: 2284      2343  7          ! was the last point more optimal, if so use it
: 2285      2344  7
: 2286      2345  7
: 2287      2346  7      IF ABS(.DIFFERENCE) GTRU ABS(.LAST_DIFF)
: 2288      2347  7      THEN
: 2289      2348  8          BEGIN
: 2290      2349  8
: 2291      2350  8          IF .REC_ADDR EQLU .LAST
: 2292      2351  8          THEN
: 2293      2352  8              IRAB[IRBSV_REC_W_LO] = 0
: 2294      2353  8          ELSE
: 2295      2354  9              BEGIN
: 2296      2355  9                  REC_ADDR = .LAST;
: 2297      2356  9
: 2298      2357  9                  IF .REC_ADDR LSSU .POS_INSERT
: 2299      2358  9                  THEN
: 2300      2359  9                      IRAB[IRBSV_REC_W_LO] = 0;
: 2301      2360  8                  END;
: 2302      2361  8
: 2303      2362  8                  LAST = 0;
: 2304      2363  7                  END;
: 2305      2364  7
: 2306      2365  7      ! 2-bkt split is possible rec_addr points to the most
: 2307      2366  7      ! optimal place
: 2308      2367  7
: 2309      2368  7
: 2310      2369  7      IF .LAST EQL 0
: 2311      2370  7      THEN      ! just backed up rec_addr, need to recalc last
: 2312      2371  8          BEGIN
: 2313      2372  8
: 2314      2373  8          LOCAL
: 2315      2374  8              TMP;
: 2316      2375  8
: 2317      2376  8          TMP = .REC_ADDR;
: 2318      2377  8          REC_ADDR = .BKT_ADDR + BKTSC_OVERHDSZ;
: 2319      2378  8          LAST = .REC_ADDR;
: 2320      2379  8
: 2321      2380  8          WHILE .REC_ADDR NEQU .TMP
: 2322      2381  8          DO
: 2323      2382  9              BEGIN
: 2324      2383  9                  LAST = .REC_ADDR;
: 2325      2384  9
: 2326      2385  9                  ! If the front compression of the current record is
```

```
: 2327      2386 9      : zero, save its address as the last noncompressed key.  
: 2328      2387 9      : This may prevent a bucket scan when it comes time to  
: 2329      2388 9      : extract and re-expand the key of the last record in  
: 2330      2389 9      : the bucket immediately before the split point.  
: 2331      2390 9      :  
: 2332      2391 9      IF .IDX_DFN[IDX$V_KEY_COMPR]  
: 2333      2392 9      THEN  
: 2334      2393 10      BEGIN  
: 2335      2394 10      IF .(.REC_ADDR + RMSREC_OVHD() + 1) < 0.8 > EQLU 0  
: 2336      2395 10      THEN  
: 2337      2396 10      IRAB[IRBSL_LST_NCMP] = .REC_ADDR;  
: 2338      2397 10      END;  
: 2339      2398 9      RMSGETNEXT_REC();  
: 2340      2399 9      END;  
: 2341      2400 9      END;  
: 2342      2401 8      RMSMOVE KEY(.LAST, .REC_ADDR);  
: 2343      2402 8      IRAB[IRBSW_SPLIT] = .REC_ADDR - .BKT_ADDR;  
: 2344      2403 7      : treat another exception case of the new record going off into  
: 2345      2404 7      : a cont. bkt all by itself  
: 2346      2405 7      :  
: 2347      2406 7      IF .IRAB[IRBSW_SPLIT] EQLU .IRAB[IRBSW_POS_INS]  
: 2348      2407 7      THEN  
: 2349      2408 7      IF .IRAB[IRBSW_SPLIT] EQLU .IRAB[IRBSW_SPLIT_1]  
: 2350      2409 7      THEN  
: 2351      2410 7      IF NOT .IRAB[IRBSV_REC_W_LO]  
: 2352      2411 7      THEN  
: 2353      2412 7      BEGIN  
: 2354      2413 7      BUILTIN  
: 2355      2414 7      AP;  
: 2356      2415 7      AP = 3;  
: 2357      2416 7      IF NOT RMSCOMPARE_KEY ( KEYBUF_ADDR(2),  
: 2358      2417 7      KEYBUF_ADDR(3),  
: 2359      2418 7      .IDX_DFN[IDX$B_KEYSZ] )  
: 2360      2419 7      THEN  
: 2361      2420 8      IRAB[IRBSV_CONT_BKT] = 1;  
: 2362      2421 8      END;  
: 2363      2422 8      LEAVE DO_IT  
: 2364      2423 8      END;  
: 2365      2424 8      ! end of * BLOCK 5 *  
: 2366      2425 8      : the magic ratio isn't positive yet, so save all the context and  
: 2367      2426 8      : move on to the next record  
: 2368      2427 8      LAST_DIFF = .DIFFERENCE;  
: 2369      2428 8  
: 2370      2429 8  
: 2371      2430 8  
: 2372      2431 8  
: 2373      2432 8  
: 2374      2433 7  
: 2375      2434 7  
: 2376      2435 7  
: 2377      2436 7  
: 2378      2437 6  
: 2379      2438 6  
: 2380      2439 6  
: 2381      2440 6  
: 2382      2441 6  
: 2383      2442 6
```

```

: 2384      2443 6      LAST = .REC_ADDR;
: 2385      2444 6
: 2386      2445 6      IF .IRAB[IRBSV_REC_W_LO]
: 2387      2446 6      THEN
: 2388      2447 6          SAVE_REC_W_LO = 1;
: 2389      2448 6
: 2390      2449 5      END;                ! end of * BLOCK 3 *
: 2391      2450 5
: 2392      2451 5      ! Go get the next record, but special case when we are at the position
: 2393      2452 5      ! of insert.
: 2394      2453 5
: 2395      2454 5      NEXT :
: 2396      2455 6          BEGIN
: 2397      2456 6
: 2398      2457 6          IF .REC_ADDR EQLU .POS_INSERT
: 2399      2458 6              AND
: 2400      2459 6              NOT .IRAB[IRBSV_REC_W_LO]
: 2401      2460 6          THEN
: 2402      2461 7              BEGIN
: 2403      2462 7
: 2404      2463 7                  ! If this is an update, check to see if it needed an rrv, since
: 2405      2464 7                  ! the record will go in the left bucket.
: 2406      2465 7
: 2407      2466 7                  IF .IRAB[IRBSV_UPDATE]
: 2408      2467 7                  THEN
: 2409      2468 8                      BEGIN
: 2410      2469 8
: 2411      2470 8                          IF .BDB[BDB$L_VBN] EQLU .IRAB[IRBSL_PUTUP_VBN]
: 2412      2471 8                          THEN
: 2413      2472 8                              NEED_RRV = .NEED_RRV - 1;
: 2414      2473 8
: 2415      2474 7                      END;
: 2416      2475 7
: 2417      2476 7                  ! Force record to low bucket, and put in key buffer 2 the key
: 2418      2477 7                  ! of the record we are inserting (currently in keybuffer 3).
: 2419      2478 7
: 2420      2479 7                  IRAB[IRBSV_REC_W_LO] = 1;
: 2421      2480 7                  RMSMOVE(.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2));
: 2422      2481 7
: 2423      2482 7                  ! If we are inserting at the end of the bucket, or if the record
: 2424      2483 7                  ! at position of insert has a different key from that to be inserted,
: 2425      2484 7                  ! leave NEXT so that no other record goes to the left bucket (so far).
: 2426      2485 7                  ! If the key is a duplicate, then keep them together in the left
: 2427      2486 7                  ! bucket.
: 2428      2487 7
: 2429      2488 7          IF .REC_ADDR EQLU .EOB
: 2430      2489 7          THEN
: 2431      2490 7              LEAVE NEXT
: 2432      2491 7          ELSE
: 2433      2492 8              BEGIN
: 2434      2493 8
: 2435      2494 8                  BUILTIN
: 2436      2495 8                      AP;
: 2437      2496 8
: 2438      2497 8                  LOCAL
: 2439      2498 8                      CURR_KEY,
: 2440      2499 8                      REC_OVHD;
```



```
: 2441      2500      8
: 2442      2501      8      REC_OVHD = RM$REC_OVHD(0);
: 2443      2502      8
: 2444      2503      8      ! When the key is compressed, we must build it first in key
: 2445      2504      8      ! buffer 5, and then compare. This build is easy because we
: 2446      2505      8      ! can take the front chars from the key to be inserted.
: 2447      2506      8
: 2448      2507      8
: 2449      2508      8      IF .IDX_DFN[IDX$V_KEY_COMP]
: 2450      2509      8      THEN
: 2451      2510      9          BEGIN
: 2452      2511      9              CURR_KEY = KEYBUF_ADDR(5);
: 2453      2512      9              PMSMOVE ( (.REC_ADDR + .REC_OVHD + 1)<0,8>,
: 2454      2513      9                  KEYBUF_ADDR(2),
: 2455      2514      9                  .CURR_KEY );
: 2456      2515      9              RM$BUILD_KEY ( .REC_ADDR + .REC_OVHD, .CURR_KEY );
: 2457      2516      9          END
: 2458      2517      8      ELSE
: 2459      2518      8          CURR_KEY = .REC_ADDR + .REC_OVHD;
: 2460      2519      8          AP = 3;      ! Contiguous compare of keys
: 2461      2520      8
: 2462      2521      8          IF RM$COMPARE_KEY ( .CURR_KEY,
: 2463      2522      8              KEYBUF_ADDR(2),
: 2464      2523      8              .IDX_DFN[IDX$B_KEYSZ] )
: 2465      2524      8          THEN
: 2466      2525      8              LEAVE NEXT;
: 2467      2526      7          END;
: 2468      2527      7
: 2469      2528      6          END;      ! end of { at position for insert for the 1st time }
: 2470      2529      6
: 2471      2530      6      ! Now RMS will scan the bucket starting from the current record
: 2472      2531      6      ! position and keeping duplicates together, since RMS does not want to
: 2473      2532      6      ! split the bucket in the middle of a duplicate chain. Before scanning
: 2474      2533      6      ! RMS obtains the size of the current record, saves its address in
: 2475      2534      6      ! IRB$LST_NCMP, if its key is zero front compressed, and saves the
: 2476      2535      6      ! key of the current record in keybuffer 2
: 2477      2536      6
: 2478      2537      7      BEGIN
: 2479      2538      7
: 2480      2539      7      LOCAL
: 2481      2540      7          REC_OVHD,
: 2482      2541      7          S_REC_SIZE,
: 2483      2542      7          NOT_DUP;
: 2484      2543      7
: 2485      2544      7      NOT_DUP = 0;      ! assume duplicates
: 2486      2545      7
: 2487      2546      7      ! Determine the size of the current record.
: 2488      2547      7
: 2489      2548      7      REC_OVHD = RM$REC_OVHD(0; S_REC_SIZE);
: 2490      2549      7
: 2491      2550      7      ! Save the address of the current record if its key is zero front
: 2492      2551      7      ! compressed.
: 2493      2552      7
: 2494      2553      7      IF .IDX_DFN[IDX$V_KEY_COMP]
: 2495      2554      7          AND
: 2496      2555      7          (.REC_ADDR + .REC_OVHD + 1)<0,8> EQLU 0
: 2497      2556      7      THEN
```

```
: 2498      2557 7      IRAB[IRBS$L_LST_NCMP] = .REC_ADDR;
: 2499      2558 7
: 2500      2559 7      ! Move the key of the current record into keybuffer 2. Fool RMSMOVE_KEY
: 2501      2560 7      ! a little by always clearing REC_W_LO so that we get in key buffer-2
: 2502      2561 7      ! the key associated with the record we are pointing to.
: 2503      2562 7
: 2504      2563 8      BEGIN
: 2505      2564 8
: 2506      2565 8      LOCAL
: 2507      2566 8      TMP : BYTE;
: 2508      2567 8
: 2509      2568 8      TMP = .IRAB[IRBS$B SPL_BITS];
: 2510      2569 8      IRAB[IRBS$V REC_W_CO] = 0;
: 2511      2570 8      RMSMOVE KEY(.REC_ADDR, .REC_ADDR);
: 2512      2571 8      IRAB[IRBS$B SPL_BITS] = .TMP
: 2513      2572 7      END;
: 2514      2573 7
: 2515      2574 7      ! Position to the next record which does not contain a key duplicate to
: 2516      2575 7      ! that of the current record (whose key has been saved in keybuffer 2).
: 2517      2576 7
: 2518      2577 7      DO
: 2519      2578 8      BEGIN
: 2520      2579 8
: 2521      2580 8      BUILTIN
: 2522      2581 8      AP;
: 2523      2582 8
: 2524      2583 8      IF .REC_ADDR EQLU .EOB
: 2525      2584 8      THEN
: 2526      2585 8      EXITLOOP;
: 2527      2586 8
: 2528      2587 8      AP = 3;
: 2529      2588 8
: 2530      2589 8      IF .BDB[BDB$L_VBN] EQLU RMS$RECORD_VBN()
: 2531      2590 8      THEN
: 2532      2591 8      NEED_RRV = .NEED_RRV - 1;
: 2533      2592 8
: 2534      2593 8      REC_ADDR = .REC_ADDR + .REC_OVHD + .S_REC_SIZE;      ! get next rec
: 2535      2594 8
: 2536      2595 8      IF .REC_ADDR EQLU .EOB
: 2537      2596 8      THEN
: 2538      2597 8      EXITLOOP;
: 2539      2598 8
: 2540      2599 8      REC_OVHD = RMS$REC_OVHD(0; S_REC_SIZE);
: 2541      2600 8
: 2542      2601 8      IF .IDX_DFN[IDX$V_KEY_COMPR]
: 2543      2602 8      THEN
: 2544      2603 9      BEGIN
: 2545      2604 9
: 2546      2605 9      IF .(.REC_ADDR + .REC_OVHD)<0,8> NEQU 0
: 2547      2606 9      THEN
: 2548      2607 9      NOT_DUP = 1;
: 2549      2608 9      END
: 2550      2609 8      ELSE
: 2551      2610 9      BEGIN
: 2552      2611 9      AP = 3;      ! Contiguous compare of keys
: 2553      2612 9
: 2554      2613 9      IF RMS$COMPARE_KEY ( .REC_ADDR + .REC_OVHD,
```

```

: 2555      2614 9      KEYBUF_ADDR(2),
: 2556      2615 9      .IDX_DFN[IDX$B_KEYSZ] )
: 2557      2616 9      THEN
: 2558      2617 9      NOT_DUP = 1;
: 2559      2618 8      END;
: 2560      2619 8
: 2561      2620 8      ! If RMS is currently positioned to the point of insertion of the
: 2562      2621 8      ! updated record, and if the key of the next record matches the
: 2563      2622 8      ! key of the previous record, then the updated record must go
: 2564      2623 8      ! into the old (left) bucket.
: 2565      2624 8
: 2566      2625 8      IF .REC_ADDR EQLU .POS_INSERT
: 2567      2626 8      AND
: 2568      2627 8      NOT .NOT_DUP
: 2569      2628 8      AND
: 2570      2629 8      .IRAB[IRB$V_UPDATE]
: 2571      2630 8      THEN
: 2572      2631 9      BEGIN
: 2573      2632 9
: 2574      2633 9      IRAB[IRB$V_REC_W_LO] = 1;
: 2575      2634 9
: 2576      2635 9      IF .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
: 2577      2636 9      THEN
: 2578      2637 9      NEED_RRV = .NEED_RRV - 1;
: 2579      2638 8      END;
: 2580      2639 8
: 2581      2640 8      END
: 2582      2641 8
: 2583      2642 8      ! Loop until a non-duplicate record is found
: 2584      2643 8
: 2585      2644 7      UNTIL .NOT_DUP;
: 2586      2645 7
: 2587      2646 6      END;
: 2588      2647 6      ! end of block defining NOT_DUP
: 2589      2648 6
: 2590      2649 6      ! If the key compares brought us up to the pos of insert, see if the
: 2591      2650 6      ! key of the new record matches the key of the record before the
: 2592      2651 6      ! position of insert. If it does, have to include the new record with
: 2593      2652 6      ! the lhs.
: 2594      2653 6
: 2595      2654 6      IF .REC_ADDR EQLU .POS_INSERT
: 2596      2655 6      THEN
: 2597      2656 7      BEGIN
: 2598      2657 7
: 2599      2658 7      BUILTIN
: 2600      2659 7      AP;
: 2601      2660 7
: 2602      2661 7      AP = 3;
: 2603      2662 7
: 2604      2663 7      IF NOT RMSCOMPARE_KEY( KEYBUF_ADDR(2),
: 2605      2664 7      KEYBUF_ADDR(3),
: 2606      2665 7      .IDX_DFN[IDX$B_KEYSZ] )
: 2607      2666 7      THEN
: 2608      2667 8      BEGIN
: 2609      2668 8      IRAB[IRB$V_REC_W_LO] = 1;
: 2610      2669 8
: 2611      2670 8      IF .IRAB[IRB$V_UPDATE]
```

```
2612 2671 8          AND
2613 2672 8          .BDB[BDB$L_VBN] EQLU .IRAB[IRB$L_PUTUP_VBN]
2614 2673 8          THEN
2615 2674 8          NEED_RRV = .NEED_RRV - 1;
2616 2675 7          END;
2617 2676 7
2618 2677 6          END;
2619 2678 6
2620 2679 5          END;          ! end of NEXT
2621 2680 4          END;          ! end of * BLOCK 2 *
2622 2681 4
2623 2682 3          END;          ! end of HALF
2624 2683 3
2625 2684 3          ! define a new block here so local storage can be redefined
2626 2685 3          !
2627 2686 4          BEGIN
2628 2687 4
2629 2688 4          MACRO
2630 2689 4          BEG_CHAIN = LHS %,
2631 2690 4          END_CHAIN = RHS %,
2632 2691 4          DUPS = RRV %;
2633 2692 4
2634 2693 4          LOCAL
2635 2694 4          FIRST_KEY_EXPANSION;
2636 2695 4
2637 2696 4          BUILTIN
2638 2697 4          AP;
2639 2698 4
2640 2699 4          ! If we end up with a duplicate chain here, we need to account for the
2641 2700 4          ! the fact that the first record which would end up in a new bucket
2642 2701 4          ! will have it's first key expanded fully. Initialize the expansion
2643 2702 4          ! amount to 0.
2644 2703 4
2645 2704 4          FIRST_KEY_EXPANSION = 0;
2646 2705 4
2647 2706 4          ! Must be a 3 or 4 bucket split or we detected ascending order and the new
2648 2707 4          ! record was a dupe. We'll optimize here to the extent of trying to keep a
2649 2708 4          ! dup chain around the new record together and in the middle bucket.
2650 2709 4          ! Note that in all the cases that follow the new record is going into the
2651 2710 4          ! middle bucket. Therefore, the 'lhs' will always fit, since it can only
2652 2711 4          ! get smaller (or stay the same size, in the degenerate case). Also note
2653 2712 4          ! that in any of these cases, the left hand bucket may be empty of data
2654 2713 4          ! records (have only rrv's in it) if the first split point is at the
2655 2714 4          ! beginning and all data records get moved
2656 2715 4
2657 2716 4          IRAB[IRB$V_NEW_BKTS] = 2;    ! assume 3-bkt split until shown otherwise
2658 2717 4          IRAB[IRB$V_REC_W_LO] = 0;
2659 2718 4
2660 2719 4          ! Initialize key buffer 2 with the contents of key buffer 3 (the value
2661 2720 4          ! of the primary key of the record being inserted). This is necessary
2662 2721 4          ! when the new record is at the beginning of the bucket and is going into
2663 2722 4          ! a bucket all by itself so that all the records in the bucket need rrv's
2664 2723 4          ! since they all move into the next bucket.
2665 2724 4          ! At any rate, that seems to be the only case where key buffer 2 is not
2666 2725 4          ! correct coming into here and will be set correctly before leaving.
2667 2726 4
2668 2727 4          RMSMOVE(.IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2));
```

```
: 2669      2728  4
: 2670      2729  4      ! Find beginning and end of this possible dups chain equal to the key value
: 2671      2730  4      ! of the record being inserted.
: 2672      2731  4
: 2673      2732  4      REC_ADDR = .BKT_ADDR + BKT$C_OVERHDSZ;
: 2674      2733  5      BEGIN
: 2675      2734  5
: 2676      2735  5      LOCAL
: 2677      2736  5          STATUS,
: 2678      2737  5          REC_OVHD,
: 2679      2738  5          S_REC_SIZE,
: 2680      2739  5          CORR_KEY;
: 2681      2740  5
: 2682      2741  5      WHILE 1
: 2683      2742  5      DO
: 2684      2743  6          BEGIN
: 2685      2744  6
: 2686      2745  6          REC_OVHD = RM$REC_OVHD(0; S_REC_SIZE);
: 2687      2746  6
: 2688      2747  6          ! If the key is compressed, it must be rebuilt into keybuffer 5 first
: 2689      2748  6          !
: 2690      2749  6
: 2691      2750  6          IF .IDX_DFN[IDX$V_KEY_COMPR]
: 2692      2751  6          THEN
: 2693      2752  7              BEGIN
: 2694      2753  7                  CURR_KEY = KEYBUF_ADDR(5);
: 2695      2754  7                  RM$BUILD_KEY ( .REC_ADDR + .REC_OVHD, .CURR_KEY );
: 2696      2755  7                  END
: 2697      2756  6          ELSE
: 2698      2757  6              ! Otherwise, we are already pointing to the beginning of the key
: 2699      2758  6              !
: 2700      2759  6              CURR_KEY = .REC_ADDR + .REC_OVHD;
: 2701      2760  6          AP = 3;      ! Contiguous compare of keys
: 2702      2761  6          STATUS = RM$COMPARE_KEY ( .CURR_KEY,
: 2703      2762  6              KEYBUF_ADDR(3),
: 2704      2763  6              .IDX_DFN[IDX$B_KEYSZ] );
: 2705      2764  6
: 2706      2765  6          IF NOT .STATUS      ! If key matched, found beginning of chain
: 2707      2766  6          THEN
: 2708      2767  6              EXITLOOP;
: 2709      2768  6
: 2710      2769  6          IF .REC_ADDR LSSU .POS_INSERT
: 2711      2770  6          THEN
: 2712      2771  6              RM$MOVE( .IDX_DFN[IDX$B_KEYSZ], .CURR_KEY, KEYBUF_ADDR(2) );
: 2713      2772  6
: 2714      2773  6          IF .REC_ADDR EQLU .EOB
: 2715      2774  6              OR
: 2716      2775  6              .STATUS LSS 0
: 2717      2776  6          THEN
: 2718      2777  7              BEGIN
: 2719      2778  7
: 2720      2779  7              ! !!!! SPLIT TYPE 3 !!!!
: 2721      2780  7              ! No duplicates found. For simplicity, do a 3-bkt split at the
: 2722      2781  7              ! point of insert with the new record in its own bucket.
: 2723      2782  7              !
: 2724      2783  7              IRAB[IRB$W_SPLIT] = IRAB[IRB$W_SPLIT_1] = .IRAB[IRB$W_POS_INS];
: 2725      2784  7              LEAVE DO_IT
```

```
: 2726      2785 7
: 2727      2786 7      ! { end of didn't find a duplicate, put record in its own bucket }
: 2728      2787 7
: 2729      2788 6      END;
: 2730      2789 6
: 2731      2790 6      REC_ADDR = .REC_ADDR + .REC_OVHD + .S_REC_SIZE;
: 2732      2791 5      END;      ! { end of while no duplicate has been found }
: 2733      2792 5
: 2734      2793 4      END;      ! { end of block defining status for while loop }
: 2735      2794 4
: 2736      2795 4      ! Found the beginning of the dups chain, now find the end.
: 2737      2796 4
: 2738      2797 4      BEG_CHAIN = .REC_ADDR;
: 2739      2798 4
: 2740      2799 5      BEGIN
: 2741      2800 5
: 2742      2801 5      LOCAL
: 2743      2802 5          NOT_DUP,
: 2744      2803 5          REC_OVHD,
: 2745      2804 5          S_REC_SIZE;
: 2746      2805 5
: 2747      2806 5      NOT_DUP = 0;      ! assume more duplicates
: 2748      2807 5      REC_OVHD = RM$REC_OVHD(0; S_REC_SIZE);
: 2749      2808 5
: 2750      2809 5      ! Ok, keep track of how much the first key would expand if placed
: 2751      2810 5      ! at the beginning of a new bucket.
: 2752      2811 5
: 2753      2812 5      IF .IDX_DFN[IDX$V_KEY_COMPR]
: 2754      2813 5      THEN
: 2755      2814 5          FIRST_KEY_EXPANSION = .(.REC_ADDR + .REC_OVHD + 1)<0,8>;
: 2756      2815 5
: 2757      2816 5      DO
: 2758      2817 6          BEGIN
: 2759      2818 6
: 2760      2819 6          REC_ADDR = .REC_ADDR + .REC_OVHD + .S_REC_SIZE;
: 2761      2820 6          IF .REC_ADDR EQ[U] .EOB
: 2762      2821 6          THEN
: 2763      2822 6              EXITLOOP;
: 2764      2823 6
: 2765      2824 6          REC_OVHD = RM$REC_OVHD(0; S_REC_SIZE);
: 2766      2825 6
: 2767      2826 6          IF .IDX_DFN[IDX$V_KEY_COMPR]
: 2768      2827 6          THEN
: 2769      2828 7              BEGIN
: 2770      2829 7
: 2771      2830 7                  IF .(.REC_ADDR + .REC_OVHD)<0,8> NEQU 0
: 2772      2831 7                  THEN
: 2773      2832 7                      NOT_DUP = 1
: 2774      2833 7                  END
: 2775      2834 6              ELSE
: 2776      2835 7                  BEGIN
: 2777      2836 7
: 2778      2837 7                      AP = 3;      ! Contiguous compare of keys
: 2779      2838 7
: 2780      2839 7                      IF RM$COMPARE_KEY ( .REC_ADDR + .REC_OVHD,
: 2781      2840 7                          KEYBUF_ADDR(3),
: 2782      2841 7                          .IDX_DFN[IDX$B_KEYSZ] )
```

```
2783 2842 7      THEN
2784 2843 7      NOT_DUP = 1
2785 2844 6      END;
2786 2845 6
2787 2846 6      END
2788 2847 6
2789 2848 5      UNTIL .NOT_DUP
2790 2849 4      END;          ! end of found end of dups chain
2791 2850 4
2792 2851 4      END_CHAIN = .REC_ADDR;
2793 2852 4
2794 2853 4      ! Found the beginning and the end of the chain. Calculate its size.
2795 2854 4      ! If we got here via an update, we never called RMSSRCH_BY_KEY to set
2796 2855 4      ! DUPS_SEEN for us, so let us do that now if necessary. Also be sure
2797 2856 4      ! to factor in the amount of key expansion that the first key would
2798 2857 4      ! undergo if placed first in a new bucket. If the keys aren't
2799 2858 4      ! compressed, don't sweat it -- FIRST_KEY_EXPANSION was initialized
2800 2859 4      ! to zero, and only changed if key compression is in effect.
2801 2860 4
2802 2861 4      IF .POS_INSERT GTRU .BEG_CHAIN
2803 2862 4      THEN
2804 2863 4          IRAB[IRBSV_DUPS_SEEN] = 1;
2805 2864 4
2806 2865 4      DUPS = .END_CHAIN - .BEG_CHAIN;
2807 2866 4      DUPS = .DUPS + .RECSZ + .FIRST_KEY_EXPANSION;
2808 2867 4
2809 2868 4      IF .DUPS LSSU .BKTSIZE
2810 2869 4      THEN
2811 2870 5          BEGIN
2812 2871 5
2813 2872 5          +
2814 2873 5          !!!!! SPLIT TYPE 1 !!!!!
2815 2874 5          Duplicates found and fortunately, they all fit in one bucket,
2816 2875 5          so do a 3-bkt split with all of the dups in the middle bucket.
2817 2876 5          Because of the optimization used for dups being inserted "in order"
2818 2877 5          this can still be a 2-bkt split if the new record is being inserted
2819 2878 5          at the end of the bucket.
2820 2879 5
2821 2880 5          22-Jan-79 If LOA forced us to think that a bkt with all dups had to
2822 2881 5          be split ( only on put) be smart and just put new record by itself.
2823 2882 5          A better solution would be not to split at all, but at this date
2824 2883 5          it's rather inconceivable.
2825 2884 5
2826 2885 5          23-Jan-79 It's not only LOA that can fool us, the bkt might have
2827 2886 5          had a lot of rrv's.
2828 2887 5          -
2829 2888 5
2830 2889 5          IRAB[IRBSW_SPLIT] = .BEG_CHAIN - .BKT_ADDR;
2831 2890 5          IRAB[IRBSW_SPLIT_1] = .END_CHAIN - .BKT_ADDR;
2832 2891 5
2833 2892 5          IF .END_CHAIN EQLU .EOB
2834 2893 5          THEN
2835 2894 6              BEGIN
2836 2895 6                  IRAB[IRBSV_NEW_BKTS] = 1;
2837 2896 6
2838 2897 6                  IF .BEG_CHAIN EQLU (.BKT_ADDR + BKTSC_OVERHDSZ)
2839 2898 6                  THEN
```

```
2840      2899 7      BEGIN
2841      2900 7      IRAB[IRBSW_SPLIT_1] = .IRAB[IRBSW_SPLIT_2];
2842      2901 7      IRAB[IRBSW_SPLIT] = .IRAB[IRBSW_POS_INS];
2843      2902 7      IRAB[IRBSV_CONT_BKT] = 1;
2844      2903 7      END
2845      2904 7
2846      2905 6      END
2847      2906 5      ELSE
2848      2907 6      BEGIN
2849      2908 6
2850      2909 6      IF .IRAB[IRBSW_SPLIT] EQLU BKTSC_OVERHDSZ<0, 16>
2851      2910 6      THEN
2852      2911 6          IRAB[IRBSV_EMPTY_BKT] = 1;
2853      2912 6
2854      2913 6      ! Only force record into the low bucket if it is not the first
2855      2914 6      ! one in a duplicate chain.
2856      2915 6
2857      2916 6
2858      2917 6      IF .END_CHAIN GEQU .POS_INSERT
2859      2918 6      AND .IRAB[IRBSW_SPLIT] NEQU .IRAB[IRBSW_POS_INS]
2860      2,19 6      THEN
2861      2920 6          IRAB[IRBSV_REC_W_LO] = 1;
2862      2921 5      END;
2863      2922 5
2864      2923 5      LEAVE DO_IT
2865      2924 5
2866      2925 4      END;      ! { end of duplicates found and they fit in one bucket }
2867      2926 4
2868      2927 4      ! This next test can only happen on an update so the all dupes case
2869      2928 4      ! will fall thru to split type 2, which will put the new record by itself.
2870      2929 4      ! Consider oddball update case in which there are dups before and after
2871      2930 4      ! position of insert. ( note that if this case doesn't apply, the duplicates
2872      2931 4      ! were only before or after -- and didn't fit with record -- so new record
2873      2932 4      ! will end up by itself. For code flow purposes, leave that till later).
2874      2933 4
2875      2934 4
2876      2935 4      IF .IRAB[IRBSV_DUPS_SEEN]
2877      2936 4      AND
2878      2937 4      .END_CHAIN GTRU .POS_INSERT
2879      2938 4      THEN
2880      2939 5      BEGIN
2881      2940 5
2882      2941 5      IF .DUPS - (.POS_INSERT - .BEG_CHAIN) LSSU .BKTSIZE
2883      2942 5      THEN
2884      2943 5
2885      2944 5      ! if high dups will fit with record, put them in a bucket together
2886      2945 5
2887      2946 6      BEGIN
2888      2947 6
2889      2948 6      !+
2890      2949 6      !!!!! SPLIT TYPE 4 !!!!!
2891      2950 6      ! 3 bkt split where middle bkt is a continuation bkt containing
2892      2951 6      ! new record and dups following it
2893      2952 6
2894      2953 6      !!!!! AND SPLIT TYPE 4B !!!!! however, if the hi set consists
2895      2954 6      ! solely of duplicates, we can still have a 2-bkt split case that
2896      2955 6      ! would not have been picked up by the previous algorithm ( since
```



```
: 2897      2956 6      ! it won't divide dups).
: 2898      2957 6      !-
: 2899      2958 6
: 2900      2959 6      IRAB[IRBSV_CONT_BKT] = 1;
: 2901      2960 6      IRAB[IRBSW_SPLIT] = .IRAB[IRBSW_POS_INS];
: 2902      2961 6
: 2903      2962 6      IF .END_CHAIN EQLU .EOB
: 2904      2963 6      THEN
: 2905      2964 6          IRAB[IRBSV_NEW_BKTS] = 1
: 2906      2965 6      ELSE
: 2907      2966 6          IRAB[IRBSW_SPLIT_1] = .END_CHAIN - .BKT_ADDR;
: 2908      2967 6
: 2909      2968 6      REC_ADDR = .BEG_CHAIN;
: 2910      2969 6      RMSMOVE ( .IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2) );
: 2911      2970 6      LEAVE DO_IT
: 2912      2971 6
: 2913      2972 5      END;
: 2914      2973 5
: 2915      2974 5      ! try to fit new record with before-dups in middle bucket
: 2916      2975 5      !
: 2917      2976 5
: 2918      2977 5      IF .DUPS - (.END_CHAIN - .POS_INSERT) LSSU .BKTSIZE
: 2919      2978 5      THEN
: 2920      2979 6          BEGIN
: 2921      2980 6
: 2922      2981 6          !+
: 2923      2982 6          !!!!! SPLIT TYPE 5 !!!!!
: 2924      2983 6          ! 3 or 4 bkt split ( depending on status of
: 2925      2984 6          ! high set) where left-middle bkt is new record with before-dups
: 2926      2985 6          ! and right-middle bkt, if it is needed, is a continuation bkt
: 2927      2986 6          ! with the after-dups. it is needed if the dups aren't the whole hi
: 2928      2987 6          ! set it still is a continuation bkt.
: 2929      2988 6
: 2930      2989 6          ***** NOTE FROM NOV-7-78
: 2931      2990 6          ! This case doesn't take into account the fact that the
: 2932      2991 6          ! whole bucket may be dups. In the case of all dups, we could
: 2933      2992 6          ! end up generating an empty bucket when we don't have to (if
: 2934      2993 6          ! no RRV's) or a relatively useless bucket (some RRV's). In any
: 2935      2994 6          ! event we could end up generating an extra bucket when we
: 2936      2995 6          ! don't have to
: 2937      2996 6          !-
: 2938      2997 6
: 2939      2998 6      IRAB[IRBSW_SPLIT] = .BEG_CHAIN - .BKT_ADDR;
: 2940      2999 6      IRAB[IRBSW_SPLIT_1] = .IRAB[IRBSW_POS_INS];
: 2941      3000 6
: 2942      3001 6      IF .IRAB[IRBSW_SPLIT] EQLU BKT$C_OVERHDSZ<0, 16>
: 2943      3002 6      THEN
: 2944      3003 6          IRAB[IRBSV_EMPTY_BKT] = 1;
: 2945      3004 6
: 2946      3005 6      IRAB[IRBSV_REC_W_LO] = 1;
: 2947      3006 6
: 2948      3007 6      IF .END_CHAIN LSSU .EOB
: 2949      3008 6      THEN
: 2950      3009 7          BEGIN
: 2951      3010 7          IRAB[IRBSV_NEW_BKTS] = 3;
: 2952      3011 7          IRAB[IRBSW_SPLIT_2] = .END_CHAIN - .BKT_ADDR;
: 2953      3012 7          END
```

```

2954 3013 6      ELSE
2955 3014 6          IRAB[IRBSV_CONT_R] = 1;
2956 3015 6
2957 3016 6      LEAVE DO_IT
2958 3017 6
2959 3018 5      END;
2960 3019 5
2961 3020 5      ! { end of oddball update case with dups on both sides of new record }
2962 3021 5      !
2963 3022 4      END;
2964 3023 4
2965 3024 4
2966 3025 4      !+
2967 3026 4      !!!!! SPLIT TYPE 2 !!!!!
2968 3027 4      ! the new record must go all by itself therefore,
2969 3028 4      ! this is a 3-bkt split if there are no after-dups or no hi set and a 4-bkt
2970 3029 4      ! split if both of those exist even more exceptional, this can still be a
2971 3030 4      ! 2-bkt split if there is no hi set at all ---- i.e., eob = end of the dups
2972 3031 4      ! chain
2973 3032 4      !-
2974 3033 4      IRAB[IRBSW_SPLIT] = IRAB[IRBSW_SPLIT_1] = .IRAB[IRBSW_POS_INS];
2975 3034 4
2976 3035 4      IF .IRAB[IRBSV_DUPS_SEEN]
2977 3036 4      THEN
2978 3037 5          BEGIN
2979 3038 5              IRAB[IRBSV_CONT_BKT] = 1;
2980 3039 5              REC_ADDR = .BEG_CHAIN;
2981 3040 5              RMSMOVE ( .IDX_DFN[IDX$B_KEYSZ], KEYBUF_ADDR(3), KEYBUF_ADDR(2) );
2982 3041 4              END;
2983 3042 4
2984 3043 4      IF .POS_INSERT EQLU .EOB
2985 3044 4      THEN
2986 3045 4          IRAB[IRBSV_NEW_BKTS] = 1
2987 3046 4      ELSE
2988 3047 4          IF .POS_INSERT LSSU .END_CHAIN
2989 3048 4          THEN
2990 3049 4              BEGIN
2991 3050 5
2992 3051 5                  IF .END_CHAIN LSSU .EOB
2993 3052 5                  THEN
2994 3053 5                      IRAB[IRBSV_NEW_BKTS] = 3
2995 3054 5                  ELSE
2996 3055 5                      IRAB[IRBSV_CONT_R] = 1;
2997 3056 5
2998 3057 5                  IRAB[IRBSW_SPLIT_2] = .END_CHAIN - .BKT_ADDR;
2999 3058 5                  END;
3000 3059 4
3001 3060 4      END;
3002 3061 3      ! { end of block defining local symbols }
3003 3062 3
3004 3063 2      END;
3005 3064 2      ! { end of do_it }
3006 3065 2
3007 3066 2      ! if the first split point is at the beginning of the data, this means that
3008 3067 2      ! all data records will be moved out and only rrv's will be left in the
3009 3068 2      ! original bucket ..... therefore, we can mark this bucket as empty
3010 3069 2
```

```

. { end of routine }

```

PC	BB	00000	RM\$SPLIT	UDR 3::	
				PUSHR	#*M<R2,R3>
				SUBL2	#36, SP
				BICB2	#72, 68(IRAB)
				INSV	#1, #1, #2, 68(IRAB)
				CLRL	NEED_RRV
				PUSHL	REC_ADDR
				MOVAB	14(R5), 40(SP)
				MOVL	40(SP), REC_ADDR
				MOVZWL	4(BKT_ADDR), R0
				PUSHAB	(R0)[BKT_ADDR]
				CLRL	LAST
			1\$:	BBS	#3, (REC_ADDR), 4\$
				MOVL	#3, AP
				BSBW	RM\$RECORD_VBN
				CMPL	28(BDB), R0
				BNFQ	2\$
				INCL	NEED_RRV
			2\$:	MOVL	REC_ADDR, LAST
				BBC	#6, -28(IDX_DFN), 3\$
				BSBW	RM\$REC_OVHD
				TSTB	1(R0)[REC_ADDR]
				BNEQ	3\$
				MOVL	REC_ADDR, 152(IRAB)
			3\$:	BSBW	RM\$GETNEXT_REC
				CMPL	REC_ADDR, EOB
				BLSSU	1\$
			4\$:	TSTL	LAST
				BEQL	5\$
				MOVL	REC_ADDR, TMP_ADDR
				MOVL	LAST, REC_ADDR
				CLRL	AP
				MOVZWL	180(IFAB), R0
				MULL2	#3, R0
				PUSHAB	296(IRAB)[R0]
				BSBW	RM\$RECORD_KEY
				ADDL2	#4, SP
				MOVL	TMP_ADDR, REC_ADDR
			5\$:	SUBL3	BKT_ADDR, REC_ADDR, R0
				MOVW	R0, -78(IRAB)
				MOVW	R0, 76(IRAB)
				MOVZBL	23(IDX_DFN), R1
				ASHL	#9, R1, R1

10	AE	04	AE	2C	AE	D4	0015F	CLRL	SAVE_REC_W_LO	:	2137
		08	AE		56	C3	00162	SUBL3	REC_ADDR, EOR, RHS	:	2150
					56	D1	00168	CMPL	REC_ADDR, POS_INSERT	:	2152
					0A	1A	0016C	BGTRU	16\$:	
	05	44	A9		03	E0	0016E	BBS	#3, 68(IRAB), 16\$:	2154
		10	AE	40	AE	C0	00173	ADDL2	RECSZ, RHS	:	2156
	3E	1C	A7		06	E1	00178	BBC	#6, 28(IDX DFN), 20\$:	2162
		08	AE		56	D1	0017D	CMPL	REC_ADDR, POS_INSERT	:	2173
					24	12	00181	BNEQ	19\$:	
	1F	44	A9		03	E0	00183	BBS	#3, 68(IRAB), 19\$:	2175
	0C	06	A9		03	E1	00188	BBC	#3, 6(IRAB), 17\$:	2177
		50		00B4	CA	3C	0018D	MOVZWL	180(IFAB), R0	:	2179
		50		60	B940	DE	00192	MOVAL	296(IRAB)[R0], R0	:	
					04	11	00197	BRB	18\$:	
		50		68	A9	D0	00199	MOVL	104(IRAB), R0	:	2181
		51		01	A0	9A	0019D	MOVZBL	1(R0), R1	:	
		10	AE		51	C0	001A1	ADDL2	R1, RHS	:	
					14	11	001A5	BRB	20\$:	2177
		04	AE		56	D1	001A7	CMPL	REC_ADDR, EOB	:	2191
					0E	1E	001AB	BGEQU	20\$:	
					51	D4	001AD	CLRL	R1	:	2193
				0000G	30	001AF		BSBW	RMSREC OVHD	:	
		53		01	A046	9A	001B2	MOVZBL	1(R0)[REC_ADDR], R3	:	
		10	AE		53	C0	001B7	ADDL2	R3, RHS	:	
		1C	AE	10	AE	D1	001BB	CMPL	RHS, BKTSIZE	:	2203
					03	1F	001C0	BLSSU	22\$:	
				0157	31	001C2		BRW	44\$:	
				06	A5	B5	001C5	TSTW	6(BKT_ADDR)	:	2205
					13	12	001C8	BNEQ	23\$:	
	0E	06	A9		03	E0	001CA	BBS	#3, 6(IRAB), 23\$:	2207
		50		0E	A5	9E	001CF	MOVAB	14(R5), R0	:	2209
		50			56	D1	001D3	CMPL	REC_ADDR, R0	:	
					05	12	001D6	BNEQ	23\$:	
	E5	44	A9		03	E1	001D8	BBC	#3, 68(IRAB), 21\$:	2211
	50	56			55	C3	001DD	SUBL3	BKT_ADDR, REC_ADDR, R0	:	2214
		52		F2	A0	9E	001E1	MOVAB	-14(R0), LHS	:	
		08	AE		56	D1	001E5	CMPL	REC_ADDR, POS_INSERT	:	2216
					09	1F	001E9	BLSSU	24\$:	
	04	44	A9		03	E1	001EB	BBC	#3, 68(IRAB), 24\$:	2218
		52		40	AE	C0	001F0	ADDL2	RECSZ, LHS	:	2220
	50	52		18	AE	C1	001F4	ADDL3	RRV, LHS, R0	:	2230
30	AE	0C	AE		09	C5	001F9	MULL3	#9, NEED RRV, 48(SP)	:	
		50		30	AE	C0	001FF	ADDL2	48(SP), R0	:	
		1C	AE		50	D1	00203	CMPL	R0, BKTSIZE	:	
					0F	1A	00207	BGTRU	25\$:	
				06	A5	B5	00209	TSTW	6(BKT_ADDR)	:	2237
					46	12	0020C	BNEQ	30\$:	
	41	06	A9		03	E0	0020E	BBS	#3, 6(IRAB), 30\$:	2239
	3C	44	A9		03	E1	00213	BBC	#3, 68(IRAB), 30\$:	2241
					6E	D5	00218	TSTL	LAST	:	2245
					03	12	0021A	BNEQ	26\$:	
				0280	31	0021C		BRW	62\$:	
		56			6E	D0	0021F	MOVL	LAST, REC_ADDR	:	2249
		04		2C	AE	E8	00222	BLBS	SAVE_REC_W_LO, 27\$:	2251
		44	A9		08	8A	00226	BICB2	#8, 68(IRAB)	:	2253
		53			56	D0	0022A	MOVL	REC_ADDR, TMP	:	2264
		56		0E	A5	9E	0022D	MOVAB	14(R5), REC_ADDR	:	2265

		6E	56	D0	00231	MOVL	REC_ADDR, LAST	2266		
		53	56	D1	00234	CMPL	REC_ADDR, TMP	2268		
			75	13	00237	BEQL	38\$			
		6E	56	D0	00239	MOVL	REC_ADDR, LAST	2271		
	OE	A7	06	E1	0023C	BBC	#6, -28(IDX_DFN), 29\$	2279		
			0000G	30	00241	BSBW	RMSREC_OVHD	2283		
			01	A046	95	00244	TSTB	1(R0)[REC_ADDR]		
			05	12	00248	BNEQ	29\$			
		0098	56	D0	0024A	MOVL	REC_ADDR, 152(IRAB)	2285		
			0000G	30	0024F	BSBW	RMSGETNEXT_REC	2288		
			E0	11	00252	BRB	28\$	2268		
	51	52	1C	AE	C5	00254	MULL3	BKTSIZE, LHS, R1	2332	
	50	AE	30	AE	C3	00259	SUBL3	48(SP), BKTSIZE, R0	2333	
	50	AE		50	C3	0025F	SUBL3	R0, RRV, R0		
		50	10	AE	C4	00264	MULL2	RHS, R0		
24	AE	51	50	C1	00268	ADDL3	R0, R1, DIFFERENCE			
			03	18	0026D	BGEQ	31\$	2338		
			0099	31	0026F	BRW	43\$			
		51	24	AE	D0	00272	MOVL	DIFFERENCE, R1	2346	
			03	18	00276	BGEQ	32\$			
		51		51	CE	00278	MNEGL	R1, R1		
		50	28	AE	D0	0027B	MOVL	LAST_DIFF, R0		
			03	18	0027F	BGEQ	33\$			
		50		50	CE	00281	MNEGL	R0, R0		
		50		51	D1	00284	CMPL	R1, R0		
				14	1B	00287	BLEQU	36\$		
		6E		56	D1	00289	CMPL	REC_ADDR, LAST	2350	
				09	13	0028C	BEQL	34\$		
		56		6E	D0	0028E	MOVL	LAST, REC_ADDR	2355	
	08	AE		56	D1	00291	CMPL	REC_ADDR, POS_INSERT	2357	
				04	1E	00295	BGEQU	35\$		
	44	A9		08	8A	00297	BICB2	#8, 68(IRAB)	2359	
				6E	D4	0029B	CLRL	LAST	2362	
				6E	D5	0029D	TSTL	LAST	2369	
				2A	12	0029F	BNEQ	40\$		
		53		56	D0	002A1	MOVL	REC_ADDR, TMP	2376	
		56	OE	A5	9E	002A4	MOVAB	14(R5), REC_ADDR	2377	
		6E		56	D0	002A8	MOVL	REC_ADDR, LAST	2378	
		53		56	D1	002AB	CMPL	REC_ADDR, TMP	2380	
				1B	13	002AE	BEQL	40\$		
		6E		56	D0	002B0	MOVL	REC_ADDR, LAST	2383	
	OE	A7		06	E1	002B3	BBC	#6, -28(IDX_DFN), 39\$	2391	
				0000G	30	002B8	BSBW	RMSREC_OVHD	2395	
			01	A046	95	002BB	TSTB	1(R0)[REC_ADDR]		
				05	12	002BF	BNEQ	39\$		
		0098		56	D0	002C1	MOVL	REC_ADDR, 152(IRAB)	2397	
				0000G	30	002C6	BSBW	RMSGETNEXT_REC	2400	
				E0	11	002C9	BRB	37\$	2380	
		50		6E	D0	002CB	MOVL	LAST, R0	2405	
				F76F	30	002CE	BSBW	RMSMOVE_KEY		
4A	A9	56		55	A3	002D1	SUBW3	BKT_ADDR, REC_ADDR, 74(IRAB)	2406	
		48	A9	4A	A9	B1	002D6	CMPL	74(IRAB), 72(IRAB)	2412
				0C	12	002DB	BNEQ	41\$		
		4C	A9	4A	A9	B1	002DD	CMPL	74(IRAB), 76(IRAB)	2415
				05	12	002E2	BNEQ	41\$		
	03	44	A9	03	E1	002E4	BBC	#3, 68(IRAB), 42\$	2418	
				044F	31	002E9	BRW	96\$		

	5C		03	D0	002EC	42\$:	MOVL	#3, AP	2425	
	50	00B4	CA	3C	002EF		MOVZWL	180(IFAB), R0	2428	
	53	60	B940	3E	002F4		MOVAW	@96(IRAB)[R0], R3		
51	50	60	A9	C1	002F9		ADDL3	96(IRAB), R0, R1	2427	
	50	20	A7	9A	002FE		MOVZBL	32(IDX_DFN), R0		
			0000G	30	00302		BSBW	RMSCOMPARE_KEY		
	E1		50	E8	00305		BLBS	R0, 41\$		
			0321	31	00308		BRW	78\$	2431	
	28	AE	24	AE	D0	0030B	43\$:	MOVL	DIFFERENCE, LAST_DIFF	2442
	6E		56	D0	00310		MOVL	REC_ADDR, LAST	2443	
04	44	A9		03	E1	00313		BBC	#3, 68(IRAB), 44\$	2445
	2C	AE		01	D0	00318		MOVL	#1, SAVE_REC_W_LO	2447
	08	AE		56	D1	0031C	44\$:	CMPL	REC_ADDR, POS_INSERT	2457
				03	13	00320		BEQL	46\$	
			008D	31	00322	45\$:	BRW	52\$		
F8	44	A9		03	E0	00325	46\$:	BBS	#3, 68(IRAB), 45\$	2459
0A	06	A9		03	E1	0032A		BBC	#3, 6(IRAB), 47\$	2466
	78	A9	1C	A4	D1	0032F		CMPL	28(BDB), 120(IRAB)	2470
				03	12	00334		BNEQ	47\$	
			0C	AE	D7	00336		DECL	NC2D_RRV	2472
	44	A9		08	88	00339	47\$:	BISB2	#8, 88(IRAB)	2479
	50		00B4	CA	3C	0033D		MOVZWL	180(IFAB), R0	2480
			60	B940	9F	00342		PUSHAB	@96(IRAB)[R0]	
			60	B940	3F	00346		PUSHAW	@96(IRAB)[R0]	
		7E	20	A7	9A	0034A		MOVZBL	32(IDX_DFN), -(SP)	
				0000G	30	0034E		BSBW	RMSMOVE	
		5E		0C	C0	00351		ADDL2	#12, SP	
	04	AE		56	D1	00354		CMPL	REC_ADDR, EOB	2488
				03	12	00358		BNEQ	49\$	
			FE05	31	0035A	48\$:	BRW	15\$		
			51	D4	0035D	49\$:	CLRL	R1	2501	
			0000G	30	0035F		BSBW	RMSREC_OVHD		
53	30	AE		50	D0	00362		MOVL	R0, REC_OVHD	
29		56	30	AE	C1	00366		ADDL3	REC_OVHD, REC_ADDR, R3	2512
	1C	A7		06	E1	0036B		BBC	#6, 28(IDX_DFN), 50\$	2508
		50	00B4	CA	3C	00370		MOVZWL	180(IFAB), R0	2511
		51	60	B940	DE	00375		MOVAL	@96(IRAB)[R0], CURR_KEY	
				51	DD	0037A		PUSHL	CURR_KEY	2514
			60	B940	9F	0037C		PUSHAB	@96(IRAB)[R0]	2513
		7E	01	A3	9A	00380		MOVZBL	1(R3), -(SP)	2512
				0000G	30	00384		BSBW	RMSMOVE	
		5E		08	C0	00387		ADDL2	#8, SP	
		6E		51	D0	0038A		MOVL	CURR_KEY, (SP)	2515
			34	BE46	9F	0038D		PUSHAB	@REC_OVHD[REC_ADDR]	
				F706	30	00391		BSBW	RMSBUILD_KEY	
		5E		08	C0	00394		ADDL2	#8, SP	
				03	11	00397		BRB	51\$	2508
		51		53	D0	00399	50\$:	MOVL	R3, CURR_KEY	2518
		5C		03	D0	0039C	51\$:	MOVL	#3, AP	2519
		53	00B4	CA	3C	0039F		MOVZWL	180(IFAB), R3	2522
		53	60	A9	C0	003A4		ADDL2	96(IRAB), R3	
		50	20	A7	9A	003A8		MOVZBL	32(IDX_DFN), R0	2521
				0000G	30	003AC		BSBW	RMSCOMPARE_KEY	
		A8		50	E8	003AF		BLBS	R0, 48\$	
			20	AE	D4	003B2	52\$:	CLRL	NOT_DUP	2544
				51	D4	003B5		CLRL	R1	2548
				0000G	30	003B7		BSBW	RMSREC_OVHD	

	14	AE	50	D0	003BA	MOVL	R0, REC_OVHD		
	30	AE	51	D0	003BE	MOVL	R1, 48(SP)		
OF	1C	A7	06	E1	003C2	BBC	#6, 28(IDX_DFN), 53\$	2553	
50	14	AE	01	C1	003C7	ADDL3	#1, REC_OVHD, R0	2555	
			6046	95	003CC	TSTB	(R0)[REC_ADDR]		
			05	12	003CF	BNEQ	53\$		
	0098	C9	56	D0	003D1	MOVL	REC_ADDR, 152(IRAB)	2557	
		53	44	A9	90 003D6	53\$:	MOV8	68(IRAB), TMP	2568
	44	A9	08	8A	003DA		BICB2	#8, 68(IRAB)	2569
		50	56	D0	003DE		MOVL	REC_ADDR, R0	2570
			F65C	30	003E1		BSBW	RMSMOVE_KEY	
	44	A9	53	90	003E4		MOV8	TMP, 68(IRAB)	2571
	04	AE	56	D1	003E8	54\$:	CMPL	REC_ADDR, EOB	2583
			79	13	003EC		BEQL	60\$	
		5C	03	D0	003EE		MOVL	#3, AP	2587
			0000G	30	003F1		BSBW	RMSREC_OVHD	2589
		50	1C	A4	D1 003F4		CMPL	28(BDB), R0	
			03	12	003F8		BNEQ	55\$	
			0C	AE	D7 003FA		DECL	NEED_RRV	2591
50		56	14	AE	C1 003FD	55\$:	ADDL3	REC_OVHD, REC_ADDR, R0	2593
56		50	30	AE	C1 00402		ADDL3	S REC_SIZE, R0, REC_ADDR	
	04	AE	56	D1	00407		CMPL	REC_ADDR, EOB	2595
			5A	13	0040B		BEQL	60\$	
			51	D4	0040D		CLRL	R1	2599
			0000G	30	0040F		BSBW	RMSREC_OVHD	
	14	AE	50	D0	00412		MOVL	R0, REC_OVHD	
	30	AE	51	D0	00416		MOVL	R1, 48(SP)	
08	1C	A7	06	E1	0041A		BBC	#6, 28(IDX_DFN), 56\$	2601
			14	BE46	95 0041F		TSTB	@REC_OVHD[REC_ADDR]	2605
			21	13	00423		BEQL	58\$	
			1B	11	00425		BRB	57\$	2607
		5C	03	D0	00427	56\$:	MOVL	#3, AP	2611
		53	00B4	CA	3C 0042A		MOVZWL	180(IFAB), R3	2614
		53	60	A9	C0 0042F		ADDL2	96(IRAB), R3	
51		56	14	AE	C1 00433		ADDL3	REC_OVHD, REC_ADDR, R1	2613
		50	20	A7	9A 00438		MOVZBL	32(IDX_DFN), R0	
			0000G	30	0043C		BSBW	RMSCOMPARE_KEY	
		04	50	E9	0043F		BLBC	R0, 58\$	
	20	AE	01	D0	00442	57\$:	MOVL	#1, NOT_DUP	2617
	08	AE	56	D1	00446	58\$:	CMPL	REC_ADDR, POS_INSERT	2625
			17	12	0044A		BNEQ	59\$	
		17	20	AE	E8 0044C		BLBS	NOT_DUP, 60\$	2627
	06	A9	03	E1	00450		BBC	#3, 6(IRAB), 59\$	2629
0E	44	A9	08	88	00455		BISB2	#8, 68(IRAB)	2633
	78	A9	1C	A4	D1 00459		CMPL	28(BDB), 120(IRAB)	2635
			03	12	0045E		BNEQ	59\$	
			0C	AE	D7 00460		DECL	NEED_RRV	2637
		81	20	AE	E9 00463	59\$:	BLBC	NOT_DUP, 54\$	2644
	08	AE	56	D1	00467	60\$:	CMPL	REC_ADDR, POS_INSERT	2654
			2F	12	0046B		BNEQ	61\$	
		5C	03	D0	0046D		MOVL	#3, AP	2661
		50	00B4	CA	3C 00470		MOVZWL	180(IFAB), R0	2664
		53	60	B940	3E 00475		MOVAV	@96(IRAB)[R0], R3	
51		50	60	A9	C1 0047A		ADDL3	96(IRAB), R0, R1	2663
		50	20	A7	9A 0047F		MOVZBL	32(IDX_DFN), R0	
			0000G	30	00483		BSBW	RMSCOMPARE_KEY	
		13	50	E8	00486		BLBS	R0, 61\$	

		44	A9	08	88	00489	BISB2	#8, 68(IRAB)	2668
	0A	06	A9	03	E1	0048D	BBC	#3, 6(IRAB), 61\$	2670
		78	A9	1C	A4	D1 00492	CMPL	28(BDB), 120(IRAB)	2672
				03	12	00497	BNEQ	61\$	
				0C	AE	D7 00499	DECL	NEED_RRV	2674
				FCC3	31	0049C	BRW	15\$	2147
				24	AE	D4 0049F	CLRL	FIRST_KEY_EXPANSION	2704
44	A9	02	01	02	F0	004A2	INSV	#2, #T, #2, 68(IRAB)	2716
		44	A9	08	8A	004A8	BICB2	#8, 68(IRAB)	2717
		50	50	00B4	CA	3C 004AC	MOVZWL	180(IFAB), R0	2727
				60	B940	9F 004B1	PUSHAB	@96(IRAB)[R0]	
				60	B940	3F 004B5	PUSHAW	@96(IRAB)[R0]	
			7E	20	A7	9A 004B9	MOVZBL	32(IDX_DFN), -(SP)	
				0000G	30	004BD	BSBW	RM\$MOVE	
			5E	0C	C0	004C0	ADDL2	#12, SP	
		20	AE	0E	A5	9E 004C3	MOVAB	14(R5), 32(SP)	2732
		56	56	20	AE	D0 004C8	MOVL	32(SP), REC_ADDR	
				51	D4	004CC	CLRL	R1	2745
				0000G	30	004CE	BSBW	RM\$REC_OVHD	
		28	AE	50	D0	004D1	MOVL	R0, REC_OVHD	
		30	AE	51	D0	004D5	MOVL	R1, 48(SP)	
51			56	28	AE	C1 004D9	ADDL3	REC_OVHD, REC_ADDR, R1	2754
18		1C	A7	06	E1	004DE	BBC	#6, 28(IDX_DFN), 64\$	2750
		50	50	00B4	CA	3C 004E3	MOVZWL	180(IFAB), R0	2753
		14	AE	60	B940	DE 004E8	MOVAL	@96(IRAB)[R0], CURR_KEY	
				14	AE	DD 004EE	PUSHL	CURR_KEY	2754
				51	DD	004F1	PUSHL	R1	
				F5A4	30	004F3	BSBW	RM\$BUILD_KEY	
			5E	08	C0	004F6	ADDL2	#8, SP	
				04	11	004F9	BRB	65\$	2750
		14	AE	51	D0	004FB	MOVL	R1, CURR_KEY	2759
		5C	5C	03	D0	004FF	MOVL	#3, AP	2760
		0C	AE	00B4	CA	3C 00502	MOVZWL	180(IFAB), 12(SP)	2762
			50	0C	AE	D0 00508	MOVL	12(SP), R0	
			53	60	B940	3E 0050C	MOVAW	@96(IRAB)[R0], R3	
			50	20	A7	9A 00511	MOVZBL	32(IDX_DFN), R0	2761
			51	14	AE	D0 00515	MOVL	CURR_KEY, R1	
				0000G	30	00519	BSBW	RM\$COMPARE_KEY	
		2C	AE	50	D0	0051C	MOVL	R0, STATUS	
		42	AE	2C	AE	E9 00520	BLBC	STATUS, 70\$	2765
		08	AE	56	D1	00524	CMPL	REC_ADDR, POS_INSERT	2769
				15	1E	00528	BGEQU	66\$	
			50	0C	AE	D0 0052A	MOVL	12(SP), R0	2771
				60	B940	9F 0052E	PUSHAB	@96(IRAB)[R0]	
				18	AE	DD 00532	PUSHL	CURR_KEY	
			7E	20	A7	9A 00535	MOVZBL	32(IDX_DFN), -(SP)	
				0000G	30	00539	BSBW	RM\$MOVE	
			5E	0C	C0	0053C	ADDL2	#12, SP	
		04	AE	56	D1	0053F	CMPL	REC_ADDR, EOB	2773
				05	13	00543	BEQL	67\$	
				2C	AE	D5 00545	TSTL	STATUS	2775
				0F	18	00548	BGEQ	69\$	
			50	48	A9	3C 0054A	MOVZWL	72(IRAB), R0	2783
		4C	A9	50	B0	0054E	MOVW	R0, 76(IRAB)	
		4A	A9	50	B0	00552	MOVW	R0, 74(IRAB)	
				01E2	31	00556	BRW	96\$	2784
50		56	56	28	AE	C1 00559	ADDL3	REC_OVHD, REC_ADDR, R0	2790

			63	1F	00642	BLSSU	86\$		
48	A9	4A	A9	B1	00644	CMPL	74(IRAB), 72(IRAB)	2918	
			5C	13	00649	BEQL	86\$		
44	A9		08	88	0064B	BISB2	#8, 68(IRAB)	2920	
			56	11	0064F	BRB	86\$	2923	
		44	A9	95	00651	TSTB	68(IRAB)	2935	
			03	19	00654	BLSS	83\$		
			0086	31	00656	BRW	89\$		
08	AE	10	AE	D1	00659	CMPL	RHS, POS_INSERT	2937	
			7F	1B	0065E	BLEQU	89\$		
50	52	08	AE	C3	00660	SUBL3	POS_INSERT, LHS, R0	2941	
	50	18	AE	C0	00665	ADDL2	RRV, R0		
1C	AE		50	D1	00669	CMPL	R0, BKTSIZE		
			3A	1E	0066D	BGEQU	87\$		
44	A9		10	88	0066F	BISB2	#16, 68(IRAB)	2959	
4A	A9	48	A9	B0	00673	MOVW	72(IRAB), 74(IRAB)	2960	
04	AE	10	AE	D1	00678	CMPL	RHS, EOB	2962	
			08	12	0067D	BNEQ	84\$		
44	A9	02	01	F0	0067F	INSV	#1, #1, #2, 68(IRAB)	2964	
			06	11	00685	BRB	85\$		
	4C	A9	10	A3	00687	SUBW3	BKT_ADDR, RHS, 76(IRAB)	2966	
			56	D0	0068D	MOVL	LHS, REC_ADDR	2968	
			50	3C	00690	MOVZWL	180(IFABT), R0	2969	
		00B4	60	B940	9F	PUSHAB	296(IRAB)[R0]		
			60	B940	3F	PUSHAW	296(IRAB)[R0]		
			7E	A7	9A	MOVZBL	32(IDX DFN), -(SP)		
				0000G	30	BSBW	RMSMOVE		
			5E	C0	006A4	ADDL2	#12, SP		
				72	11	006A7	BRB	91\$	2970
50	08	AE	10	AE	C3	006A9	SUBL3	RHS, POS_INSERT, R0	2977
		50	18	AE	C0	006AF	ADDL2	RRV, R0	
1C	AE		50	D1	006B3	CMPL	R0, BKTSIZE		
			26	1E	006B7	BGEQU	89\$		
	4A	A9	52	A3	006B9	SUBW3	BKT_ADDR, LHS, 74(IRAB)	2998	
		4C	A9	B0	006BE	MOVW	72(IRAB), 76(IRAB)	2999	
			0E	A9	B1	006C3	CMPL	74(IRAB), #14	3001
				05	12	006C7	BNEQ	88\$	
44	A9	40	8F	88	006C9	BISB2	#64, 68(IRAB)	3003	
44	A9		08	88	006CE	BISB2	#8, 68(IRAB)	3005	
04	AE	10	AE	D1	006D2	CMPL	RHS, EOB	3007	
			52	1F	006D7	BLSSU	93\$		
44	A9		20	88	006D9	BISB2	#32, 68(IRAB)	3014	
			5C	11	006DD	BRB	96\$	3016	
	50	48	A9	3C	006DF	MOVZWL	72(IRAB), R0	3033	
4C	A9		50	B0	006E3	MOVW	R0, 76(IRAB)		
4A	A9		50	B0	006E7	MOVW	R0, 74(IRAB)		
		44	A9	95	006EB	TSTB	68(IRAB)	3035	
			1E	18	006EE	BGEQ	90\$		
44	A9		10	88	006F0	BISB2	#16, 68(IRAB)	3038	
	56		52	D0	006F4	MOVL	LHS, REC_ADDR	3039	
	50	00B4	60	CA	3C	006F7	MOVZWL	180(IFABT), R0	3040
			60	B940	9F	006FC	PUSHAB	296(IRAB)[R0]	
			60	B940	3F	00700	PUSHAW	296(IRAB)[R0]	
		7E	20	A7	9A	00704	MOVZBL	32(IDX DFN), -(SP)	
				0000G	30	00708	BSBW	RMSMOVE	
	5E		0C	C0	0070B	ADDL2	#12, SP		
04	AE	08	AE	D1	0070E	CMPL	POS_INSERT, EOB	3043	

44	A9	02	01		08	12	00713		BNEQ	92\$		
					01	F0	00715		INSV	#1, #2, 68(IRAB)		3045
					1E	11	0071B	91\$:	BRB	96\$		
		10	AE	08	AE	D1	0071D	92\$:	CMPL	POS_INSERT, RHS		3048
					17	1E	00722		BGEQU	96\$		
		04	AE	10	AE	D1	00724		CMPL	RHS, EOB		3052
					06	1E	00729		BGEQU	94\$		
		44	A9		06	88	0072B	93\$:	BISB2	#6, 68(IRAB)		3054
					04	11	0072F		BRB	95\$		
		44	A9		20	88	00731	94\$:	BISB2	#32, 68(IRAB)		3056
4E	A9				55	A3	00735	95\$:	SUBW3	BKT_ADDR, RHS, 78(IRAB)		3058
		10	AE	4A	A9	B1	0073B	96\$:	CMPL	74(IRAB), #14		3070
			0E		0A	12	0073F		BNEQ	97\$		
		05			03	E0	00741		BBS	#3, 68(IRAB), 97\$		3072
					44	A9	00746		BISB2	#64, 68(IRAB)		3074
			5E	40	8F	88	0074B	97\$:	ADDL2	#52, SP		3078
					34	C0	0074E		POPR	#^M<R2,R3>		
					0C	BA	0074E		RSB			
						05	00750					

; Routine Size: 1873 bytes, Routine Base: RM\$RMS3 + 05C0

; 3020	3079	1
; 3021	3080	1 END
; 3022	3081	1
; 3023	3082	0 ELUDOM

PSECT SUMMARY

Name	Bytes	Attributes
RM\$RMS3	3345	NOVEC, NOWRT, RD, EXE, NOSHR, GBL, REL, CON, PIC, ALIGN(2)

Library Statistics

File	Total	Symbols Loaded	Percent	Pages Mapped	Processing Time
_S255\$DUA28:[RMS.OBJ]RMS.L32;1	3109	63	2	154	00:00.4

COMMAND QUALIFIERS

; BLISS/CHECK=(FIELD, INITIAL, OPTIMIZE)/LIS=LIS\$:RM3SPLUDR/DBJ=OBJ\$:RM3SPLUDR MSRC\$:RM3SPLUDR/UPDATE=(ENH\$:RM3SPLUDR)

RM3SPLUDR
V04-000

RMSSPLIT_UDR_3

E 5
16-Sep-1984 02:03:28
14-Sep-1984 13:01:40

VAX-11 Bliss-32 V4.0-742
[RMS.SRC]RM3SPLUDR.B32;1

Page 75
(6)

RM
V0

: Size: 3345 code + 0 data bytes
: Run Time: 01:25.9
: Elapsed Time: 02:43.7
: Lines/CPU Min: 2152
: Lexemes/CPU-Min: 13733
: Memory Used: 634 pages
: Compilation Complete

0327 AH-BT13A-SE
VAX/VMS V4.0

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RM3PROBE
LIS

RM3IOXSP
LIS

RM3PUTERR
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RM3PUTUPD
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RM3SPLUDR
LIS

RM3RRU
LIS

RM3ROOT
LIS

RM3PUT
LIS

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